

Supporting Efficient Streaming and Insertion of XML Data in RDBMS

- **Timo Böhme**

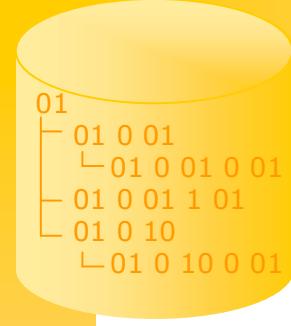
boehme@informatik.uni-leipzig.de

- **Erhard Rahm**

rahm@informatik.uni-leipzig.de

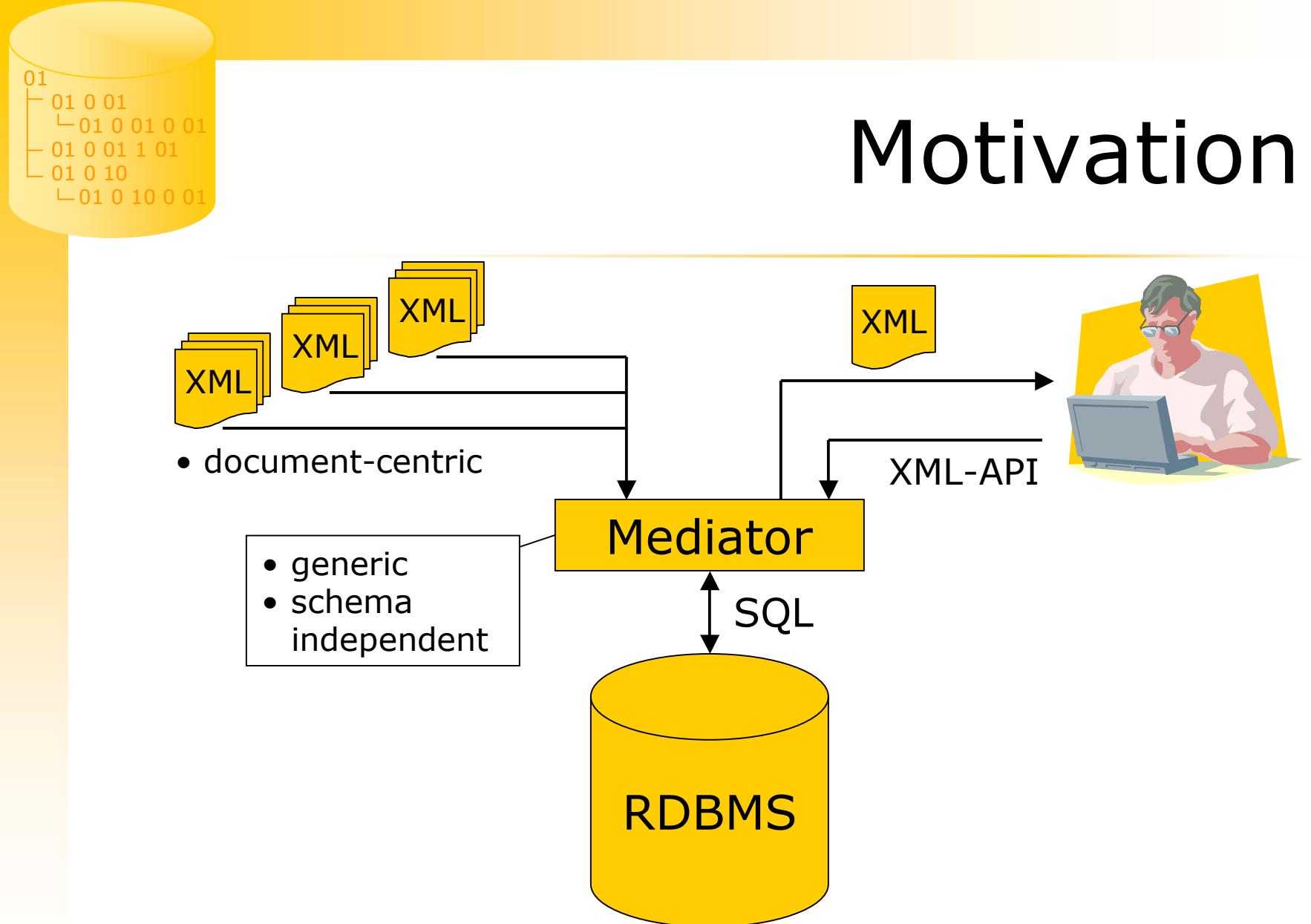
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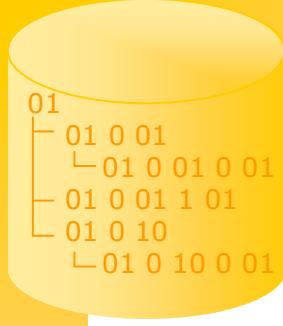
Overview



- Motivation
- Numbering Schemes
- Dynamic Level Numbering (DLN)
- Evaluation

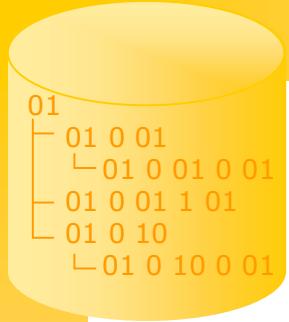
Motivation





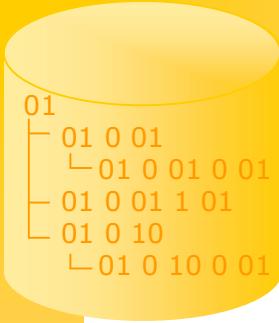
Motivation

- demand to manage XML data with RDBMS
 - mature, scalable
 - query across XML and SQL data
 - single (existing) DBMS
- mapping XML \leftrightarrow relational
 - document-centered
 - data/schema-centered
 - structure-centered



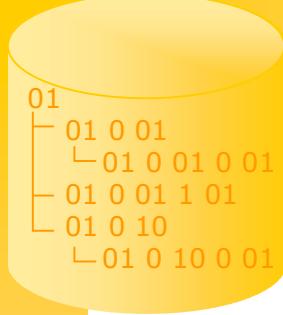
Motivation Mappings

- document-centered
 - save document as a whole (CLOB)
 - + fast storage/retrieval
 - + maintain original structure
 - updates costly
 - query evaluation may need parsing of XML data
- data/schema-centered
 - shredding XML into schema specific tables
 - + enable data-centered SQL operations
 - problems with structure-oriented queries
 - element sibling order, schema needed



Motivation Mappings (2)

- structure-centered
 - mapping graph structure into predefined relations
 - + no schema needed
 - + structure-oriented/navigational queries
 - + efficient transformation of XML interfaces
 - no direct SQL operations by user
 - mapping to parent-child relationships insufficient
 - determine ancestors/descendants
 - reconstruction of document fragments
 - solution: numbering schemes (NS)

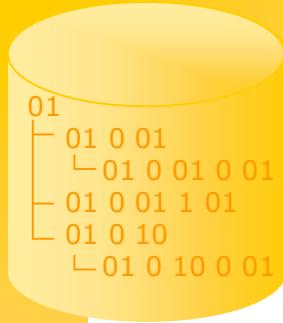


Motivation Goal

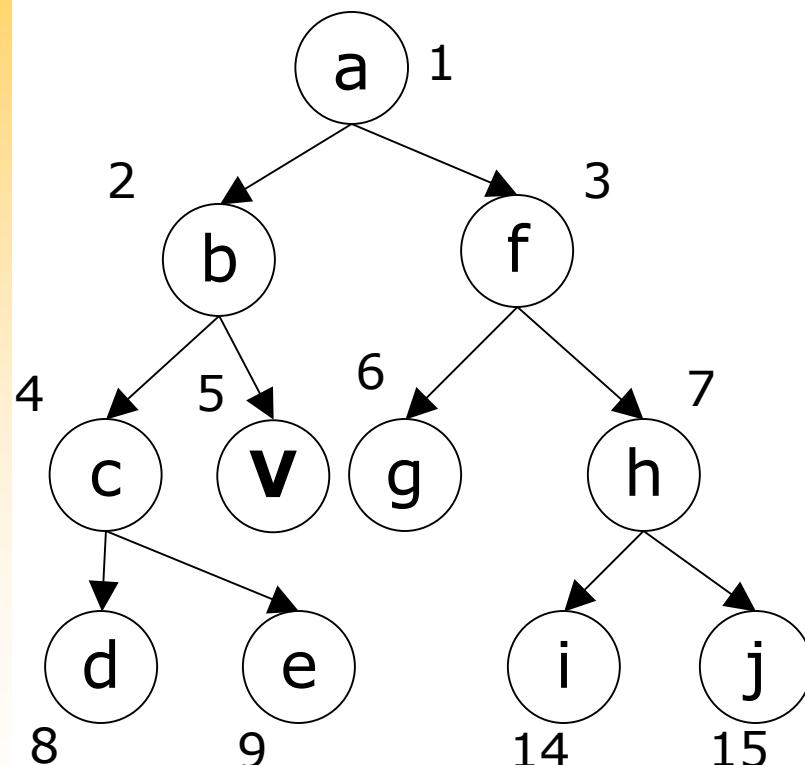
- efficiently manage document-centric XML data in RDBMS
 - support sequential processing (streaming)
 - efficient updates (insertion of complex subtrees)
 - fast reconstruction
 - no manual interaction
- best suited mapping strategy: structure-centered with semantically rich NS
 - existing NS fail on streaming or updating

Related Work

Numbering Schemes



k-ary tree

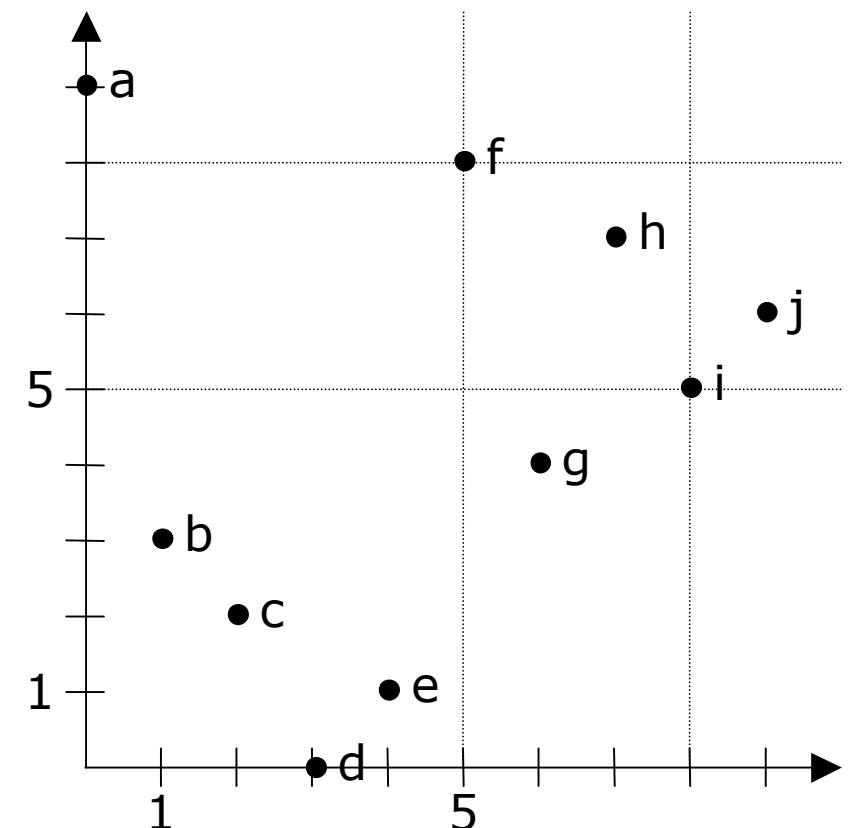
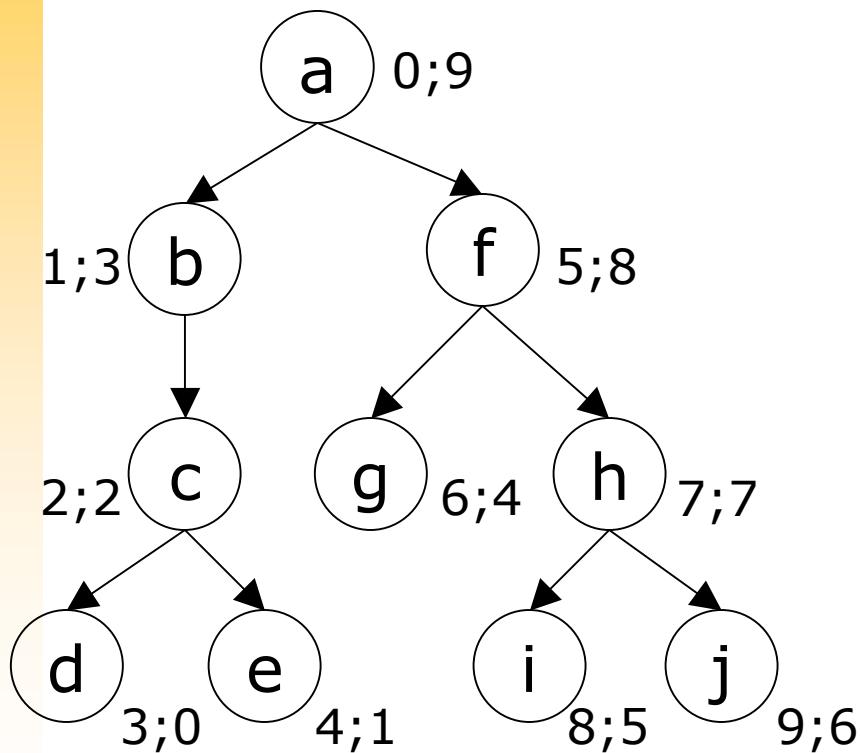


- supports
 - calculation of parent ID
$$\text{parent}(i) = \left\lfloor \frac{(i-2)}{k} + 1 \right\rfloor$$
 - calculation of j^{th} child ID
$$\text{child}(i,j) = k(i-1) + j + 1$$
- negative
 - fixed, known fan-out
 - insertion costs
 - missing support for other XPath axis

Related Work

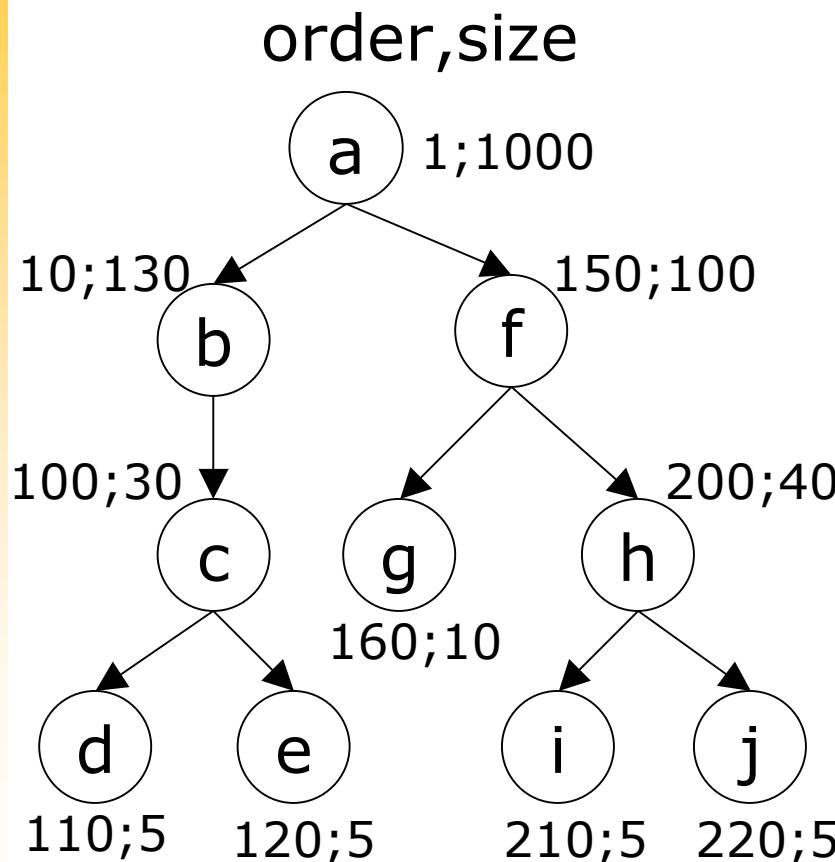
Numbering Schemes (2)

preorder, postorder

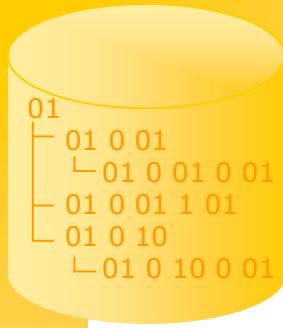


Related Work

Numbering Schemes (3)

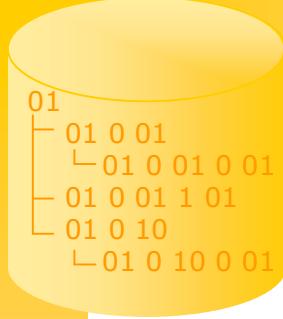


- supports
 - containment (descendant)
 - ancestor
 - restricted insert operations
- negative
 - parent, child not supported
 - assigning size values has to be delayed



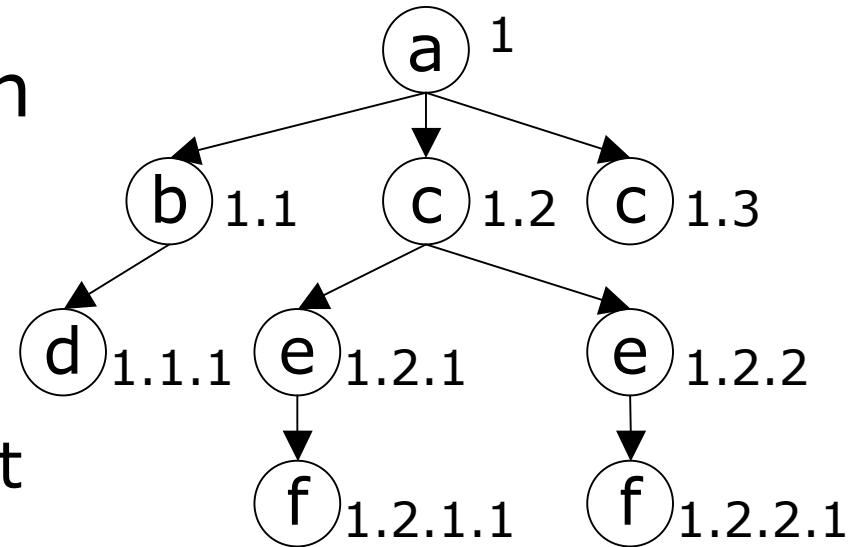
Objectives for a new NS

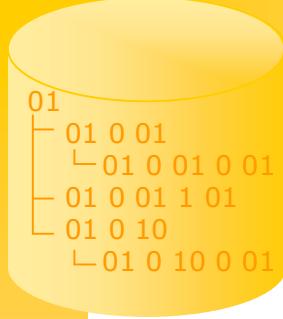
- support wide range of documents
 - irregular documents (low and high fan-out)
 - streaming of large documents
- support efficient operations
 - explicitly express total node order
 - maintain global node order
 - fast reconstruction of document fragments
 - minimize necessity for renumbering after update operations
 - efficient processing of XPath expressions
 - exploitable by query optimizer of RDBMS



Dynamic Level Numbering

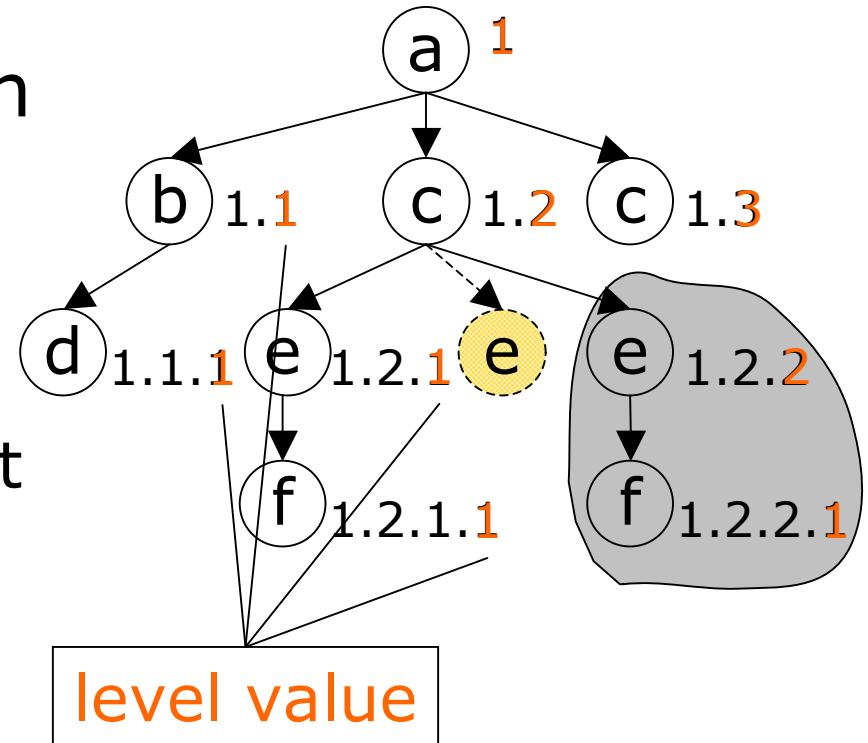
- based on Dewey Decimal Classification
- pro
 - ancestor/descendant deducible
 - sequential assignment of IDs





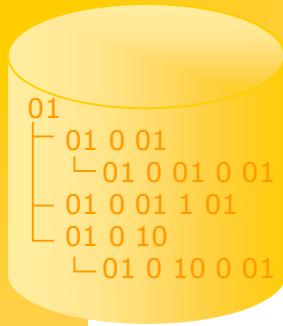
Dynamic Level Numbering

- based on Dewey Decimal Classification
- pro
 - ancestor/descendant deducible
 - sequential assignment of IDs
 - restricted update domain

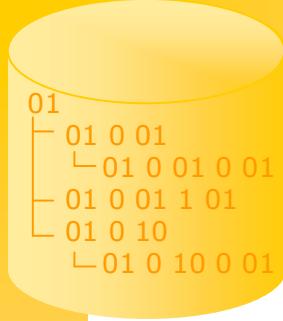


Decimal Classification

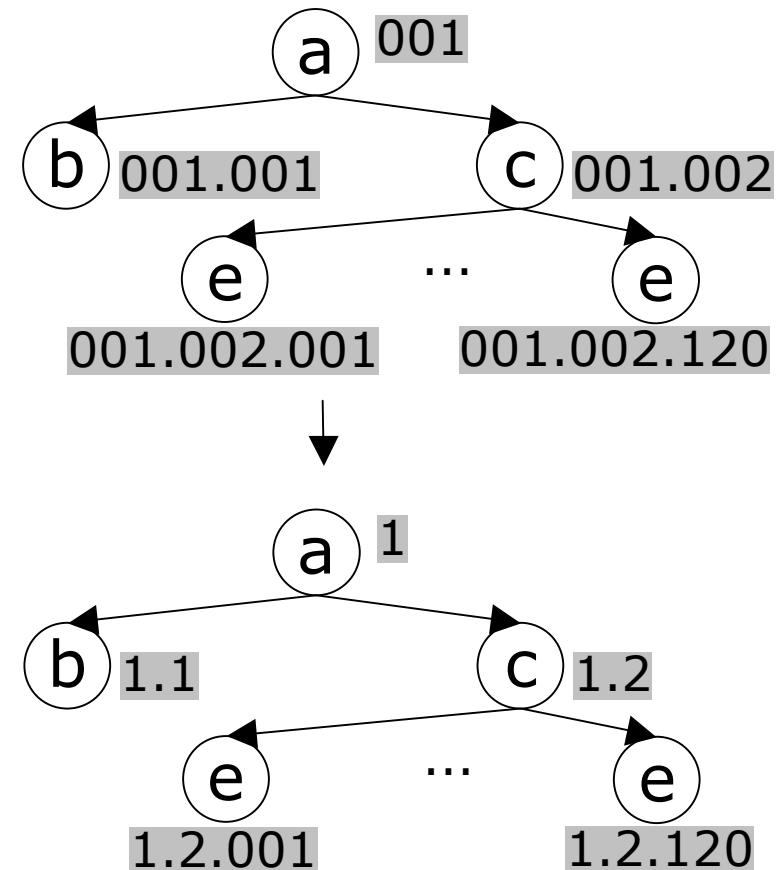
- cons of decimal classification
 - ID length depends on path length
 - updates can be costly
 - string comparison may return wrong result (e.g. 1.9 and 1.10)
- alternative:
 - level value with fix number of characters (001.009, 001.010)
 - level value coded in UTF-8 (0-127 1 Byte, 128-2047 2 Byte, 2.048-65.553 3 Byte etc., comparable at binary level)

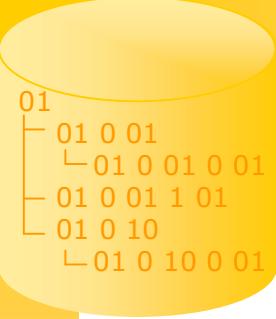


Enhancements



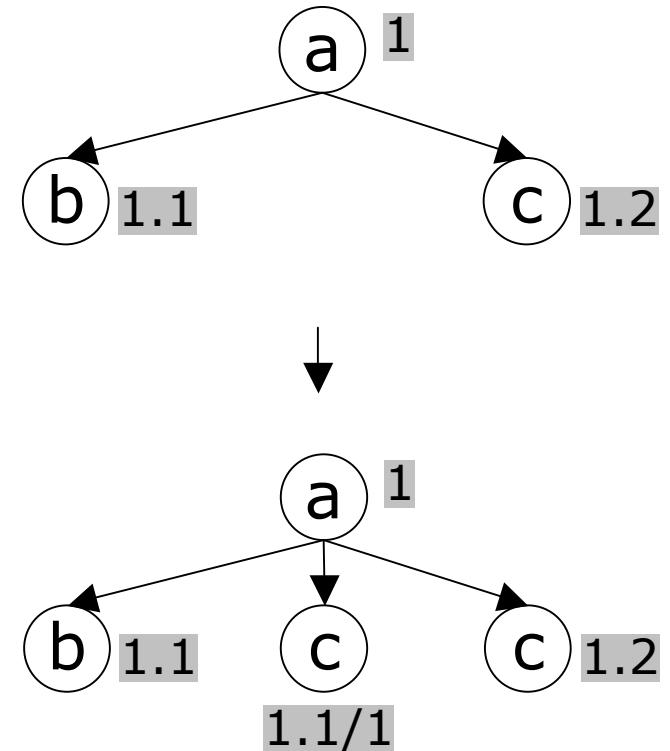
- optimization of needed characters
 - level values have variable number of digits
 - but: all sibling nodes have same number of digits

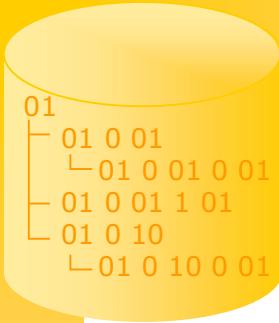




DLN Enhancements (2)

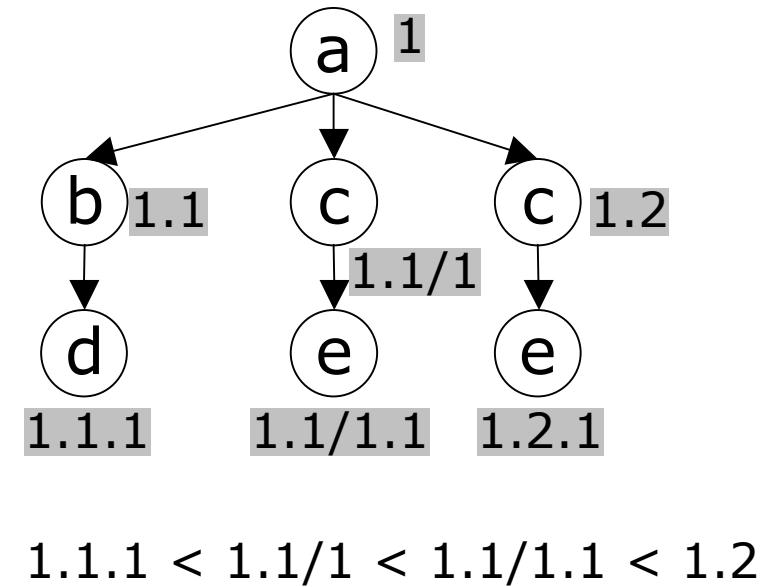
- prevention of renumbering
 - supplement level values with subvalues (recursive)
 - separate subvalues and level values by character
 - greater than level separator
 - last subvalue >0

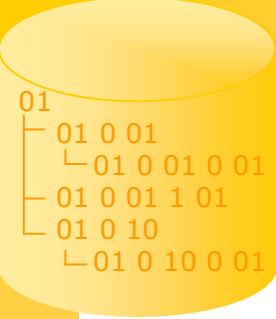




DLN Enhancements (2)

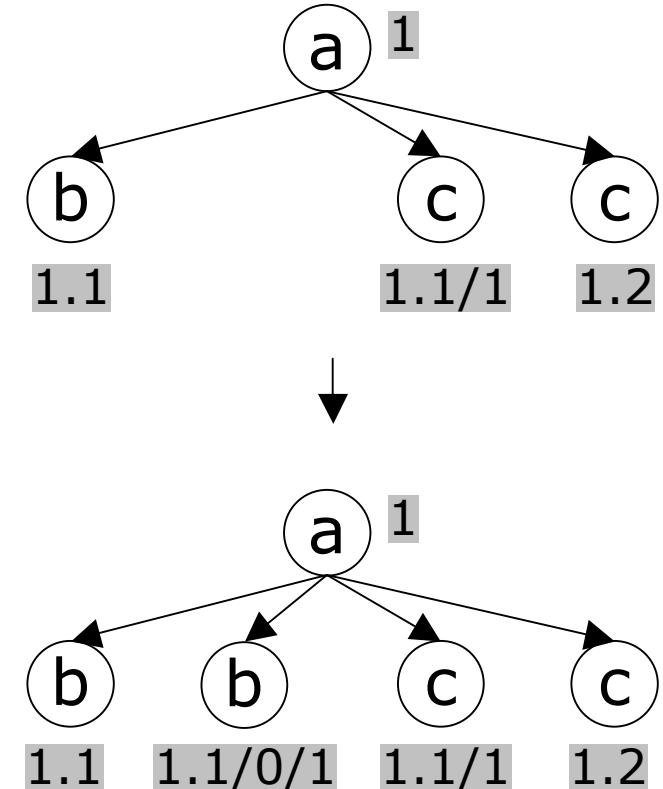
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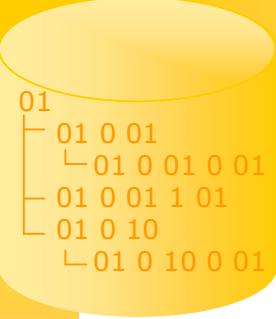




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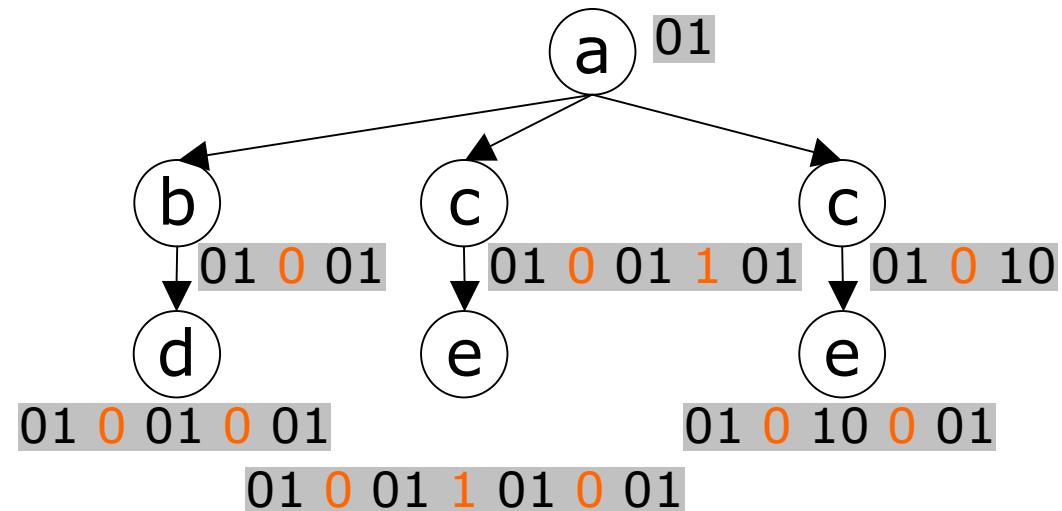
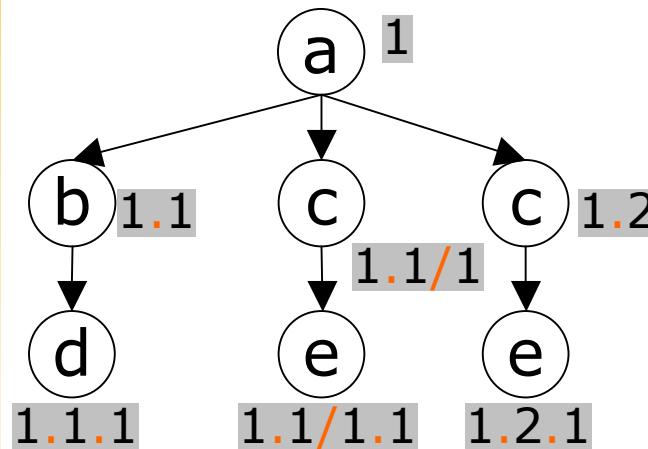
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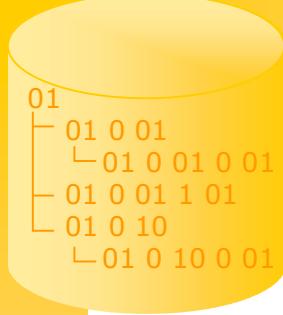
DLN Enhancements (3)

- binary coding to reduce ID length



$1.1 <$
 $1.1.1 <$
 $1.1/1 <$
 $1.1/1.1 <$
 1.2

$01 0 01 <$
 $01 0 01 0 01 <$
 $01 0 01 \textcolor{red}{1} 01 <$
 $01 0 01 1 01 0 01 <$
 $01 0 \textcolor{red}{1} 0$



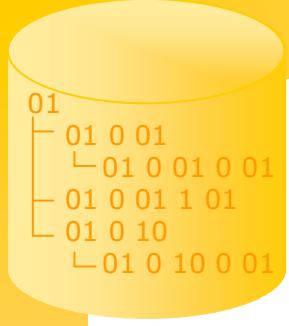
DLN Properties

- preorder numbering
- update domain
 - worst case: identical to decimal classification
 - need for renumbering highly reduced by the use of subvalues
- supports parent, ancestor/descendant, preceding, following relationships
- ID length reduced due to variable digit number and binary coding

Streamed Data

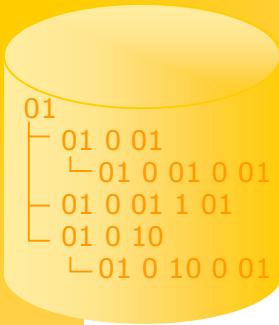
- variable digit number not applicable
- solution (algorithm)
 - start with small bit number n
 - after $n-1$ bits are exhausted add subvalue and use combined bit range

bit pattern	# of ids	enhanced # of ids
0XXX	1..7	1..7
10XX 1 XXXX	8..71	8..86
110X 1 XXXX 1 XXXX	72..583	87..613
1110 1 XXXX 1 XXXX 1 XXXX	584..4679	614..4724



DLN Variants

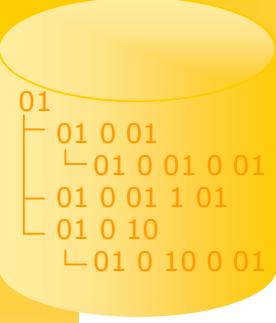
- streaming DLN
 - same number of bits for all values
 - different but fixed number of bits for level values and subvalues
- simple DOM DLN
 - variable number of bits per value
- fixed DOM DLN
 - fixed number of bits per value
- restricted DOM DLN
 - variable number of bits per value but restricted to a fixed upper bound
- no metadata
- suitable for all kinds of data
- big metadata
- minimal bit usage
- no metadata
- efficient use of combined bit range
- small metadata
- minimal overall bit usage



DLN – Evaluation

Maximum ID length

	Streaming DLN		DOM DLN			ORDPATH	
	4 bit	3/4 bit	simple	fixed	restr.	A	B
Nasa	64	2	-8	0	-9	0	-1
Cities	39	1	-1	-5	-5	5	7
Dictionary	69	-1	-13	-5	-12	9	-6
Novel	24	6	6	0	-2	3	-1
Pop. Places	44	2	-5	-10	-9	7	-1
Religion	54	-1	-8	-10	-10	-7	0
Shakespeare	44	5	-2	-5	-4	1	5
Sigmod	44	9	0	0	-2	2	7
Treebank	199	-21	-64	-5	-71	-3	-57
WFB	49	0	-3	0	-4	3	5
Courses	44	0	-3	-5	-7	-2	0



DLN – Evaluation XMLRDB Prototype

NODE

```
document : int
dlnId    : long
parent    : long
rightSibling: long
name      : int
value     : string
valType   : char
nodeType  : char
```

ATTR

```
document : int
dlnId    : long
name     : int
value    : string
```

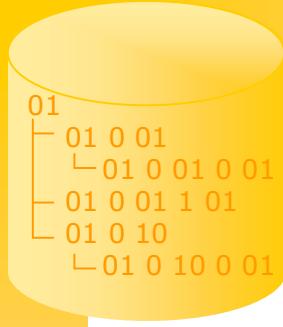
TEXT

```
document : int
dlnId    : long
value    : string
```

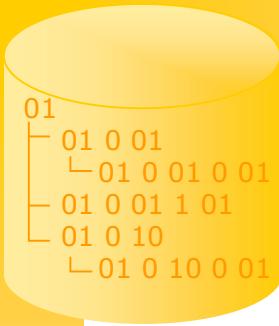
NAMEMAP

```
nameId : int
name   : string
```

DLN – Evaluation Performance



doc	#nodes	n/kB	Insertion (n/sec)			Reconstruction (n/sec)
			w/o idx	w/o txtidx	with txtidx	
XMach-1	$2,2 \cdot 10^6$	11	2.500	-	-	-
Cities	$3,5 \cdot 10^4$	44	2.380	1.729	1.699	10.310
Sigmod	$3,8 \cdot 10^4$	38	2.180	1.579	1.522	9.860
Bible	$2,5 \cdot 10^4$	8	1.899	1.405	1.152	8.525



DLN – Evaluation Performance

- queries with descendant axis (a//b) optimized by NS
 - range query with node id of a and maximum child id (calculated from a)
 - example
 - `//issue[./author='Michael Stonebraker']]//volume` – 40 ms
 - `/SigmodRecord/issue[articles/article/authors[author='Michael Stonebraker']]//volume` – 451 ms
(both queries needed 190 ms to build XML result)

Summary

- new numbering scheme (DLN) for generic structure-centered storage of XML data in RDBMS
- insert operations without renumbering of existing nodes
- usable with streamed data
- benefits from structure information
- inserting, retrieving and querying efficiently supported