

ASAP: Towards Accurate, Stable and Accelerative Penetrating-Rank Estimation on Large Graphs

Xuefei Li¹, Weiren Yu², Bo Yang³, Jiajin Le³

¹ Fudan University

² University of New South Wales

³ Donghua University

Presented by Weiren Yu

Roadmap

1

P-Rank Overview

2

Accuracy Estimate

3

Stability Analysis

4

Algorithm on Undirected Graphs

5

Empirical Evaluation

P-Rank Overview

- Information Network (IN)
 - Physical / Conceptual entities → vertices
 - Interconnected relationships → edges
- INs form a critical component of modern information infrastructure
 - highway or urban transportation networks
 - research collaboration and publication networks
 - Biological networks
 - social networks

P-Rank Overview (cont.)

- P(enetrating)-Rank similarity
 - A new promising structural measure (CIKM'09)
 - An extension of SimRank metrics
- Basic Philosophy
 - Two entities are similar, if
 - they are referenced by similar entities
 - they reference similar entities
- Mathematical Formula

$$s(u, u) = 1;$$

$$s(u, v) = \underbrace{\frac{\lambda \cdot C_{in}}{|\mathcal{I}(u)| |\mathcal{I}(v)|} \sum_{i=1}^{|\mathcal{I}(u)|} \sum_{j=1}^{|\mathcal{I}(v)|} s(\mathcal{I}_i(u), \mathcal{I}_j(v))}_{\text{in-link part}} + \underbrace{\frac{(1 - \lambda) \cdot C_{out}}{|\mathcal{O}(u)| |\mathcal{O}(v)|} \sum_{i=1}^{|\mathcal{O}(u)|} \sum_{j=1}^{|\mathcal{O}(v)|} s(\mathcal{O}_i(u), \mathcal{O}_j(v))}_{\text{out-link part}}.$$

P-Rank Overview (cont.)

- P-Rank Computation

- Naïve way: a fixed-point iterative paradigm

$$\begin{aligned} s^{(k+1)}(u, u) &= 1. \\ s^{(k+1)}(u, v) &= \frac{\lambda \cdot C_{in}}{|\mathcal{I}(u)| |\mathcal{I}(v)|} \sum_{i=1}^{|\mathcal{I}(u)|} \sum_{j=1}^{|\mathcal{I}(v)|} s^{(k)}(\mathcal{I}_i(u), \mathcal{I}_j(v)) \\ &\quad + \frac{(1-\lambda) \cdot C_{out}}{|\mathcal{O}(u)| |\mathcal{O}(v)|} \sum_{i=1}^{|\mathcal{O}(u)|} \sum_{j=1}^{|\mathcal{O}(v)|} s^{(k)}(\mathcal{O}_i(u), \mathcal{O}_j(v)). \end{aligned}$$

- Iterative P-Rank Properties

- Symmetry: $s^{(k)}(a, b) = s^{(k)}(b, a)$
- Monotonicity: $0 \leq s^{(k)}(a, b) \leq s^{(k+1)}(a, b) \leq 1$
- Existence & Uniqueness ($0 < c < 1$)

$$\lim_{k \rightarrow \infty} s^{(k)}(u, v) = \sup_{k \geq 0} \{s^{(k)}(u, v)\} = s(u, v)$$

Motivations

- Despite the convergence of P-Rank iteration, a precise P-Rank accuracy estimation is not provided.
- P-Rank condition number is not studied, which can measure how much networks may change in proportion to small perturbation in P-Rank scoring results.
- No efficient algorithm is designed specially for computing P-Rank on undirected graphs.

Contributions

- We provide an accuracy estimation of the P-Rank convergence rate with a prescribed iterative error in the fixed number of iterations.
- We show that P-Rank is well-conditioned for small choices of the damping factors, by providing a tight stability bound for κ_∞ .
- We propose a novel non-iterative $O(n^3)$ -time algorithm (ASAP) for efficiently computing similarities over undirected graphs.

Roadmap

1

P-Rank Overview

2

Accuracy Estimate

3

Stability Analysis

4

Algorithm on Undirected Graphs

5

Empirical Evaluation

P-Rank accuracy estimation

- P-Rank iterative paradigm:

$$\begin{aligned} s^{(k+1)}(u, u) &= 1. \\ s^{(k+1)}(u, v) &= \frac{\lambda \cdot C_{in}}{|\mathcal{I}(u)| |\mathcal{I}(v)|} \sum_{i=1}^{|\mathcal{I}(u)|} \sum_{j=1}^{|\mathcal{I}(v)|} s^{(k)}(\mathcal{I}_i(u), \mathcal{I}_j(v)) \\ &\quad + \frac{(1-\lambda) \cdot C_{out}}{|\mathcal{O}(u)| |\mathcal{O}(v)|} \sum_{i=1}^{|\mathcal{O}(u)|} \sum_{j=1}^{|\mathcal{O}(v)|} s^{(k)}(\mathcal{O}_i(u), \mathcal{O}_j(v)). \end{aligned}$$

$$\lim_{k \rightarrow \infty} s^{(k)}(u, v) = \sup_{k \geq 0} \{s^{(k)}(u, v)\} = s(u, v)$$

- P-Rank accuracy estimate problem:

Given a network G , for each iteration $k = 1, 2, \dots$, it is to find an upper bound ϵ_k s.t.

$$|s^{(k)}(u, v) - s(u, v)| \leq \epsilon_k$$

for any vertices u and v in G .

P-Rank accuracy estimation

- Theorem 1. The P-Rank accuracy estimate problem has a tight upper bound

$$\epsilon_k = (\lambda C_{in} + (1 - \lambda) C_{out})^{k+1}$$

such that $\forall k=0,1,\dots, \forall u, v \in V$

$$|s^{(k)}(u, v) - s(u, v)| \leq \epsilon_k.$$

- Theorem 1 provides an a-priori estimate for the gap between iterative and exact P-Rank similarity:

$$k = \lceil \log \epsilon / \log (\lambda \cdot C_{in} + (1 - \lambda) \cdot C_{out}) \rceil$$

P-Rank accuracy estimation

- Example:

Setting $C_{in} = 0.6, C_{out} = 0.4, \lambda = 0.3, k = 5$ produces the high accuracy :

$$\epsilon_k = (0.3 \times 0.6 + (1 - 0.3) \times 0.4)^{5+1} = 0.0095.$$

- The “=” in Theorem 1 can be attainable :

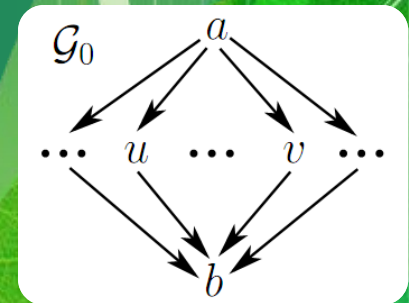
$$s^{(0)}(u, v) = 0,$$

$$\forall k=1, 2, \dots$$

$$s^{(k)}(u, v) = \lambda C_{in} + (1 - \lambda) C_{out}.$$

Hence, for $k=0$,

$$|s(u, v) - s^{(k)}(u, v)| = (\lambda C_{in} + (1 - \lambda) C_{out})^{0+1}$$



Roadmap

1

P-Rank Overview

2

Accuracy Estimate

3

Stability Analysis

4

Algorithm on Undirected Graphs

5

Empirical Evaluation

Stability Analysis of P-Rank

- P-Rank stability:
 - how the slight perturbation of the network affects P-Rank similarity scores $s(\cdot, \cdot)$.
- P-Rank Matrix Representation

$$q_{i,j} \triangleq \begin{cases} a_{j,i} / \sum_{j=1}^n a_{j,i}, & \text{if } \mathcal{I}(i) \neq \emptyset; \\ 0, & \text{if } \mathcal{I}(i) = \emptyset. \end{cases}$$

$$p_{i,j} \triangleq \begin{cases} a_{i,j} / \sum_{j=1}^n a_{i,j}, & \text{if } \mathcal{O}(i) \neq \emptyset; \\ 0, & \text{if } \mathcal{O}(i) = \emptyset. \end{cases}$$

$$\mathbf{S} = \lambda C_{\text{in}} \cdot \mathbf{Q} \cdot \mathbf{S} \cdot \mathbf{Q}^T + (1 - \lambda) C_{\text{out}} \cdot \mathbf{P} \cdot \mathbf{S} \cdot \mathbf{P}^T + (1 - \lambda C_{\text{in}} - (1 - \lambda) C_{\text{out}}) \cdot \mathbf{I}_n,$$

$$\underbrace{(\mathbf{I}_{n^2} - \lambda C_{\text{in}}(\mathbf{Q} \otimes \mathbf{Q}) - (1 - \lambda) C_{\text{out}}(\mathbf{P} \otimes \mathbf{P}))}_{\triangleq \mathbf{M}} \cdot \underbrace{\text{vec}(\mathbf{S})}_{\triangleq \mathbf{s}} = \underbrace{\text{vec}(\mathbf{I}_n)}_{\triangleq \mathbf{b}}.$$

Stability Analysis of P-Rank

- P-Rank conditional number :

Let

$$\mathbf{M} \triangleq \mathbf{I}_{n^2} - \lambda C_{in}(\mathbf{Q} \otimes \mathbf{Q}) - (1 - \lambda) C_{out}(\mathbf{P} \otimes \mathbf{P}).$$

P-Rank conditional number of G is defined as

$$\kappa_{\infty}(\mathcal{G}) \triangleq \|\mathbf{M}\|_{\infty} \cdot \|\mathbf{M}^{-1}\|_{\infty}$$

- $\kappa_{\infty}(\mathcal{G})$ measures how stable the P-Rank similarity score is to the changes in the link structure of the network G.
(e.g., inserting or deleting vertices or edges)

Stability Analysis of P-Rank

- Theorem 2. Given a network G , $\forall \lambda \in [0,1]$ and $\forall C_{in}, C_{out} \in (0,1)$, P-Rank conditional number has the following tight bound:

$$\kappa_{\infty}(\mathcal{G}) \leq \frac{1 + \lambda \cdot C_{in} + (1 - \lambda) \cdot C_{out}}{1 - \lambda \cdot C_{in} - (1 - \lambda) \cdot C_{out}}.$$

- Small choices of $\kappa_{\infty}(\mathcal{G})$ would make P-Rank stable (well-conditioned).

(i.e., a small change ΔM in link structure to M may not cause a large change Δs in P-Rank scores).

$$\frac{\|\Delta s\|_{\infty}}{\|s\|_{\infty}} \leq \kappa_{\infty}(\mathcal{G}) \cdot \frac{\|\Delta M\|_{\infty}}{\|M\|_{\infty}}$$

Stability Analysis of P-Rank

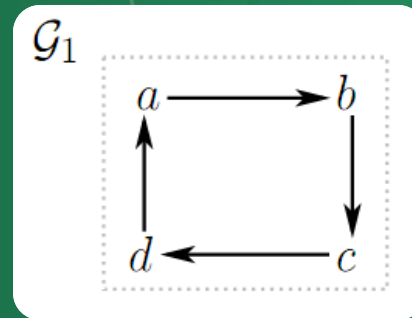
- The weighting factor λ affects $\kappa_{\infty}(G)$ as follows:

$$\frac{\partial}{\partial \lambda} \left(\frac{1 + \lambda \cdot C_{in} + (1 - \lambda) \cdot C_{out}}{1 - \lambda \cdot C_{in} - (1 - \lambda) \cdot C_{out}} \right) = \frac{2(C_{in} - C_{out})}{(1 - \lambda \cdot C_{in} - (1 - \lambda) \cdot C_{out})^2},$$

- when $C_{in} > C_{out}$ and $\lambda \nearrow$,
a small change in G produces a large change in P-Rank, which makes P-Rank ill-conditioned.
- when $C_{in} < C_{out}$ and $\lambda \nearrow$,
a small change in G produces a small change in P-Rank, which makes P-Rank well-conditioned.
- when $C_{in} = C_{out}$, $\kappa_{\infty}(G)$ is independent of λ .

Stability Analysis of P-Rank

- The upper bound of $\kappa_\infty(G)$ is attainable iff each vertex in G has at least one in-degree and one out-degree.



Example:

$$\kappa_\infty(G) = \| M \|_\infty \cdot \| M^{-1} \|_\infty = 1.7 \times 3.333 = 5.667;$$

$$\frac{1 + \lambda \cdot C_{\text{in}} + (1 - \lambda) \cdot C_{\text{out}}}{1 - \lambda \cdot C_{\text{in}} - (1 - \lambda) \cdot C_{\text{out}}} = \frac{1 + 0.5 \times 0.8 + (1 - 0.5) \times 0.6}{1 - 0.5 \times 0.8 - (1 - 0.5) \times 0.6} \doteq 5.667.$$

Roadmap

1

P-Rank Overview

2

Accuracy Estimate

3

Stability Analysis

4

Algorithm on Undirected Graphs

5

Empirical Evaluation

Estimating P-Rank On Undirected Graphs

- Theorem 3. For undirected networks, the P-Rank similarity problem

$$\mathbf{S} = \lambda C_{\text{in}} \cdot \mathbf{Q} \cdot \mathbf{S} \cdot \mathbf{Q}^T + (1 - \lambda) C_{\text{out}} \cdot \mathbf{P} \cdot \mathbf{S} \cdot \mathbf{P}^T + (1 - \lambda C_{\text{in}} - (1 - \lambda) C_{\text{out}}) \cdot \mathbf{I}_n,$$

can be solvable in $O(n^3)$ worst-case time.

Comparison:

- $O(Kn^4)$ time [CIKM 09'] via naive iterative fashion
- $O(Kn^3)$ time [EDBT 10'] via matrix iteration
- $O(n^3)$ time [this work] via non-iterative paradigm

Estimating P-Rank On Undirected Graphs

- The key idea in our optimization is to maximally use the adjacency matrix A :
 - characterizing S as a power series form

$$S = \sum_{k=0}^{+\infty} f(A^k)$$

$A = A^T$ for undirected graphs, implying $\exists D$ s.t.

$$Q = P = D \cdot A$$

- diagonalizing A into Λ to compute A^k

Hence, calculating $f(A^k)$ reduces to computing the function on each eigenvalue for A .

Estimating P-Rank On Undirected Graphs

- Proposition. For the undirected network G with n vertices, let

$$\mathbf{D} = \text{diag}\left(\left(\sum_{j=1}^n a_{1,j}\right)^{-1}, \dots, \left(\sum_{j=1}^n a_{n,j}\right)^{-1}\right),$$

and

$$[\mathbf{U}, \Lambda] = \text{eig}(\mathbf{D}^{1/2} \mathbf{A} \mathbf{D}^{1/2})$$

Then, \mathbf{S}' can be computed as

$$\mathbf{S}' = \mathbf{D}^{1/2} \mathbf{U} \cdot \Psi \cdot \mathbf{U}^T \mathbf{D}^{1/2},$$

where

$$\Psi = (\Psi_{i,j})_{n \times n} = \left(\frac{[\mathbf{U}^T \mathbf{D}^{-1} \mathbf{U}]_{i,j}}{1 - (\lambda \cdot C_{in} + (1 - \lambda) \cdot C_{out}) \Lambda_{i,i} \Lambda_{j,j}} \right)_{n \times n}$$

Estimating P-Rank On Undirected Graphs

Algorithm 1: ASAP ($\mathcal{G}, \lambda, C_{in}, C_{out}$)

Input : a labeled undirected network $\mathcal{G} = (\mathcal{V}, \mathcal{E}, \mathcal{A}; l)$, the weighting factor λ , and in- and out-link damping factors C_{in} and C_{out} .

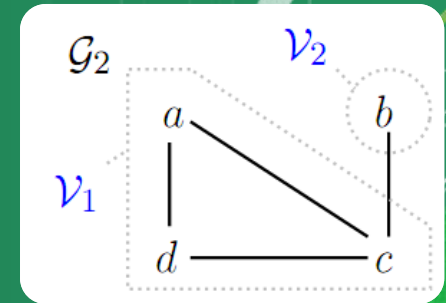
Output: similarity matrix $\mathbf{S} = (s_{i,j})_{n \times n}$ with $s_{i,j}$ denoting P-Rank score between vertices i and j .

- 1 initialize the adjacency matrix \mathbf{A} of \mathcal{G} ; $\mathcal{O}(n^2)$
- 2 compute the diagonal matrix $\mathbf{D} = \text{diag}(d_{1,1}, d_{2,2}, \dots, d_{n,n})$ $\mathcal{O}(m)$
with its entry $d_{i,i} = (\sum_{j=1}^n a_{i,j})^{-1}$, if $\sum_{j=1}^n a_{i,j} \neq 0$; and $d_{i,i} = 0$, otherwise;
- 3 compute the auxiliary matrix $\mathbf{T} = \mathbf{D}^{1/2} \cdot \mathbf{A} \cdot \mathbf{D}^{1/2}$ $\mathcal{O}(n^2)$
- 4 decompose \mathbf{T} into the diagonal matrix $\mathbf{\Lambda} = \text{diag}(\Lambda_{1,1}, \Lambda_{2,2}, \dots, \Lambda_{n,n})$ and the orthogonal \mathbf{U} via QR factorization *s.t.* $\mathbf{T} = \mathbf{U} \cdot \mathbf{\Lambda} \cdot \mathbf{U}^T$; $\mathcal{O}(n^3)$
- 5 compute the auxiliary matrix $\mathbf{\Gamma} = (\Gamma_{i,j})_{n \times n} = \mathbf{U}^T \cdot \mathbf{D}^{-1} \cdot \mathbf{U}$ and $\mathbf{V} = \mathbf{D}^{1/2} \cdot \mathbf{U}$ $\mathcal{O}(n^3+n^2)$
and the constant $C = \lambda C_{in} + (1 - \lambda) C_{out}$;
- 6 compute the matrix $\mathbf{\Psi} = (\psi_{i,j})_{n \times n}$ whose entry $\psi_{i,j} = \Gamma_{i,j} / (1 - C \cdot \Lambda_{i,i} \cdot \Lambda_{j,j})$; $\mathcal{O}(n^2)$
- 7 compute the P-Rank similarity matrix $\mathbf{S} = (1 - C) \cdot \mathbf{V} \cdot \mathbf{\Psi} \cdot \mathbf{V}^T$; $\mathcal{O}(n^3)$
- 8 **return** \mathbf{S} ;

The total time complexity of ASAP is bounded by $\mathcal{O}(n^3)$.

Estimating P-Rank On Undirected Graphs

- Running Example for ASAP:
Consider an undirected G_2 with
vertex set $V = V_1 \cup V_2 = \{a, c, d\} \cup \{b\}$
edge set $E = \{(a, c), (a, d), (c, d), (b, c)\}$.



$$\begin{aligned} \textcircled{1} \quad \mathbf{A} = \begin{pmatrix} 0 & 0 & 1 & 1 \\ 0 & 0 & 1 & 0 \\ 1 & 1 & 0 & 1 \\ 1 & 0 & 1 & 0 \end{pmatrix} &\xrightarrow{\textcircled{2}} \mathbf{D} = \{\text{Using (16)}\} = \begin{pmatrix} .5 & & & \\ & 1 & & \\ & & .333 & \\ & & & .5 \end{pmatrix} &\xrightarrow{\textcircled{3}} \mathbf{Q} = \mathbf{P} = \mathbf{D}\mathbf{A} = \begin{pmatrix} 0 & 0 & .5 & .5 \\ 0 & 0 & 1 & 0 \\ .333 & .333 & 0 & .333 \\ .5 & 0 & .5 & 0 \end{pmatrix} &\xrightarrow{\textcircled{4}} \mathbf{\Lambda} = \text{eigval}(\mathbf{D}^{1/2}\mathbf{A}\mathbf{D}^{1/2}) = \begin{pmatrix} -.729 & & 0 \\ & -.5 & \\ & & .229 \\ 0 & & & 1 \end{pmatrix} &\mathbf{U} = \text{eigvec}(\mathbf{D}^{1/2}\mathbf{A}\mathbf{D}^{1/2}) = \begin{pmatrix} -.244 & .707 & .436 & .5 \\ -.583 & 0 & -.732 & .354 \\ .736 & 0 & -.290 & .612 \\ -.244 & -.707 & .436 & .5 \end{pmatrix} \end{aligned}$$

$$\begin{aligned} \textcircled{5} \quad \mathbf{\Gamma} \triangleq \mathbf{U}^T \mathbf{D}^{-1} \mathbf{U} = \begin{pmatrix} 2.201 & 0 & -.640 & .656 \\ 0 & 2 & 0 & 0 \\ -.640 & 0 & 1.549 & .081 \\ .656 & 0 & .081 & 2.25 \end{pmatrix} &\xrightarrow{\textcircled{6}} \mathbf{\Psi} = \{\text{Using Eq.(18)}\} = \begin{pmatrix} 3.231 & 0 & -.582 & .457 \\ 0 & 2.353 & 0 & 0 \\ -.582 & 0 & 1.599 & .094 \\ .457 & 0 & .094 & 5.625 \end{pmatrix} &\xrightarrow{\textcircled{7}} \mathbf{S} = \{\text{Using Eq.(17)}\} = \begin{pmatrix} .627 & .225 & .134 & .156 \\ .225 & .770 & .067 & .225 \\ .134 & .067 & .615 & .134 \\ .156 & .225 & .134 & .627 \end{pmatrix} \end{aligned}$$

Roadmap

1

P-Rank Overview

2

Accuracy Estimate

3

Stability Analysis

4

Algorithm on Undirected Graphs

5

Empirical Evaluation

Experimental Evaluation

- Dataset

DBLP Data	1998-1999	1998-2001	1998-2003	1998-2005	1998-2007
<i>n</i>	1,525	3,208	5,307	7,984	10,682
<i>m</i>	5,929	13,441	24,762	39,399	54,844

Real-life.

DBLP (co-authorships among scientists from 1998 to 2007)

*The papers published on 6 conferences are picked up
("ICDE", "VLDB", "SIGMOD", "WWW", "SIGIR", "KDD").*

Synthetic.

*Using a C++ boost generator to produce graphs
with vertices ranging from 100K to 1M
and edges being randomly chosen*

Algorithms.

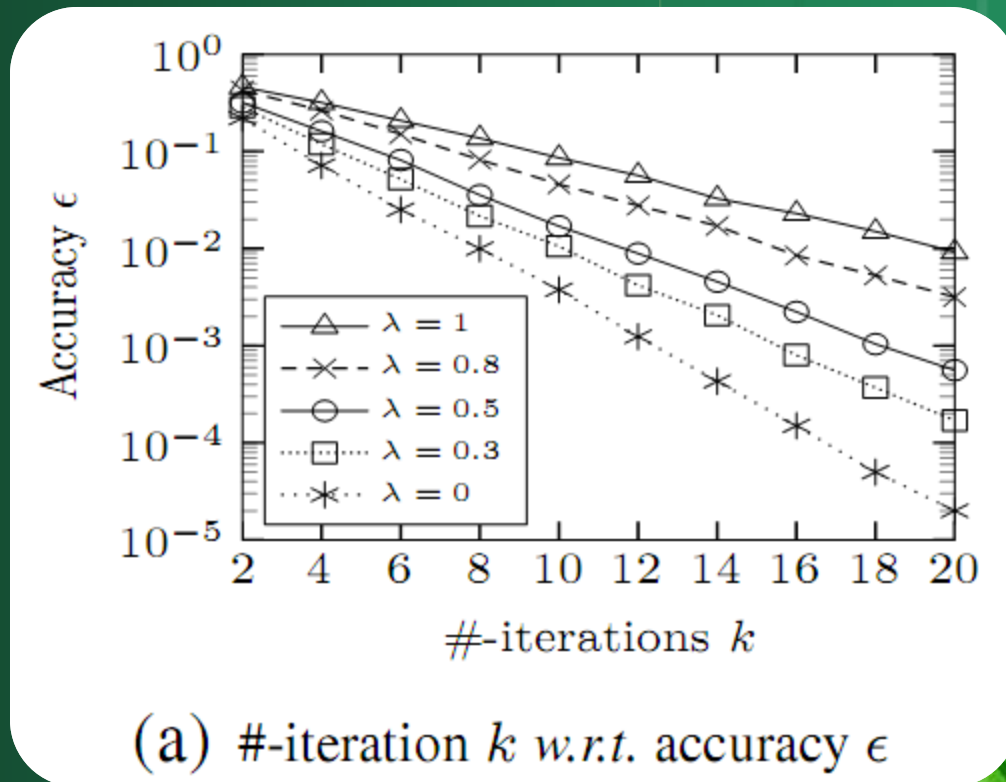
(i) Iter: conventional P-Rank algorithm [CIKM '09]
with the radius-based pruning technique

(ii) Memo: the memoization-based algorithm [VLDB J. '10]

(iii) AUG : SimRank algorithm [WAIM '10] on undirected graphs.

Experimental Results

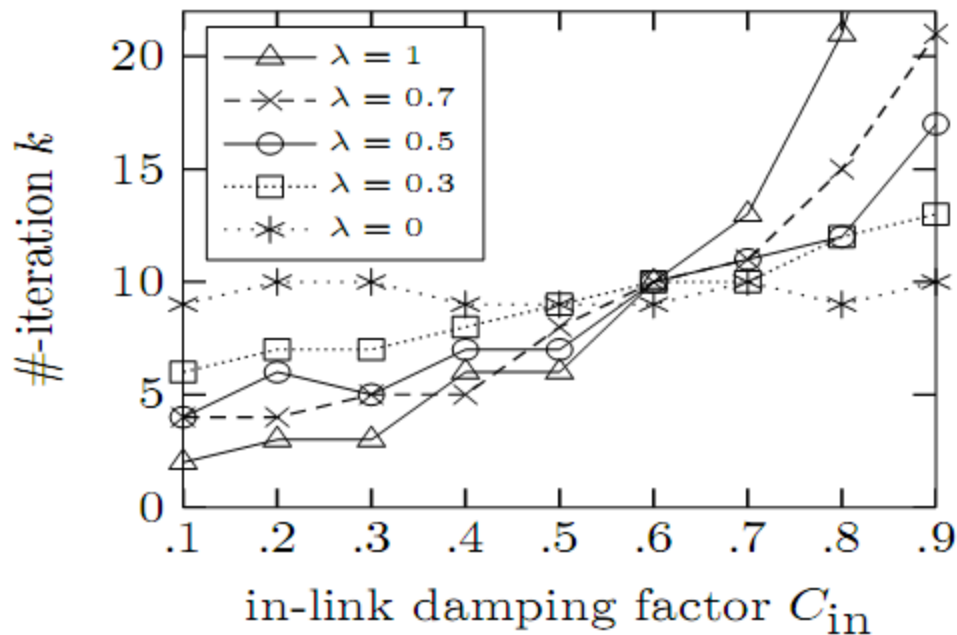
- P-Rank Accuracy



For each fixed λ , the downward lines for P-Rank iterations reveal an exponential accuracy as k increases, as expected in Theorem 1.

Experimental Results

- P-Rank Accuracy

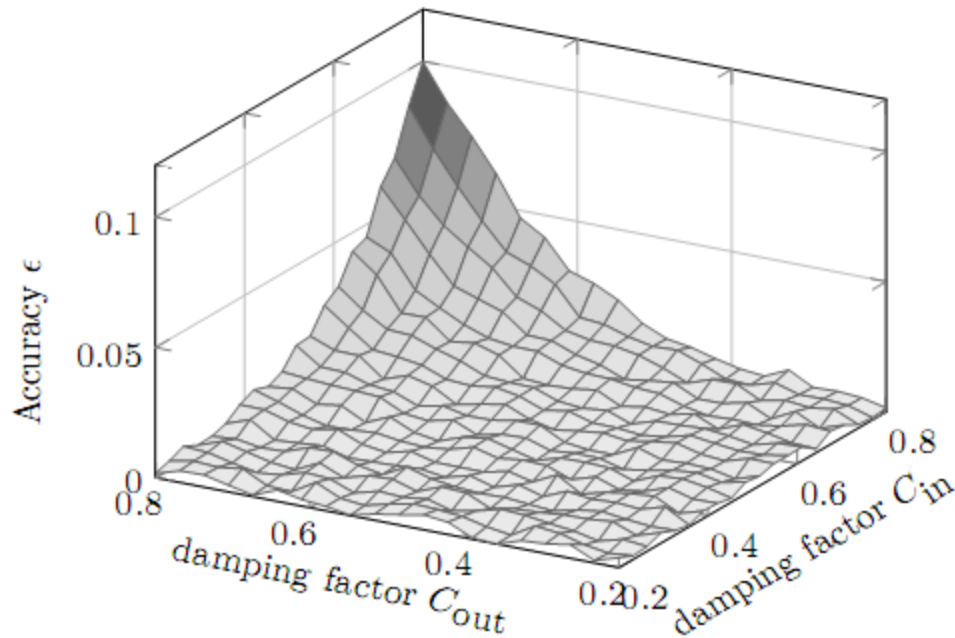


(b) damping factor C_{in} w.r.t. k

When $0 < \lambda \leq 1$, k shows a general increased tendency as C_{in} is growing. This tells us that small choices of damping factors may reduce the number of iterations required for a fixed accuracy.

Experimental Results

- P-Rank Accuracy

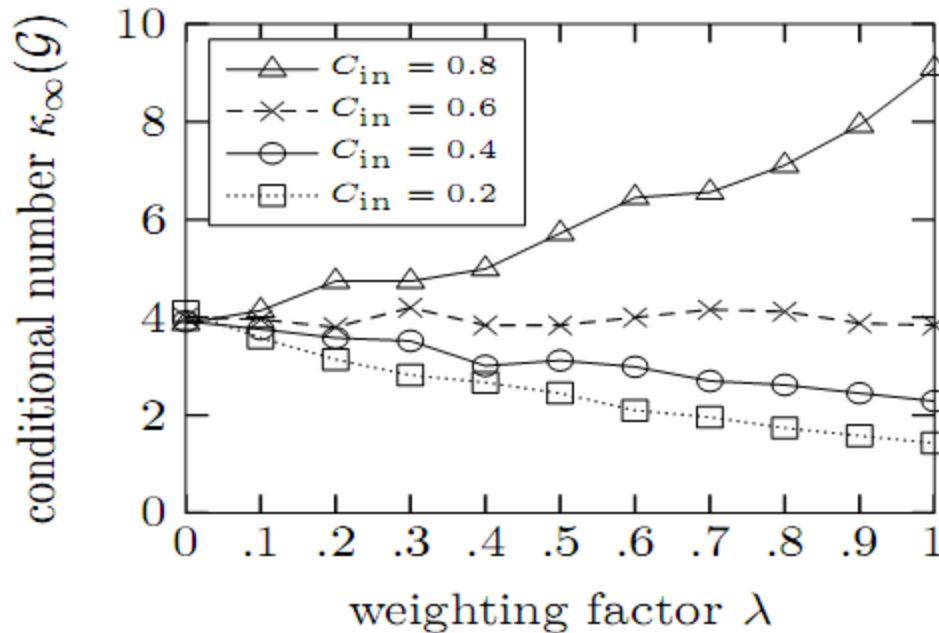


(c) damping factors C_{in} , C_{out} w.r.t. ϵ

The residual becomes huge only when C_{in} and C_{out} are both increasing to 1; and the iterative P-Rank is accurate when C_{in} and C_{out} are less than 0.6. This explains why small choices of damping factors are suggested in P-Rank iteration.

Experimental Results

- P-Rank Stability



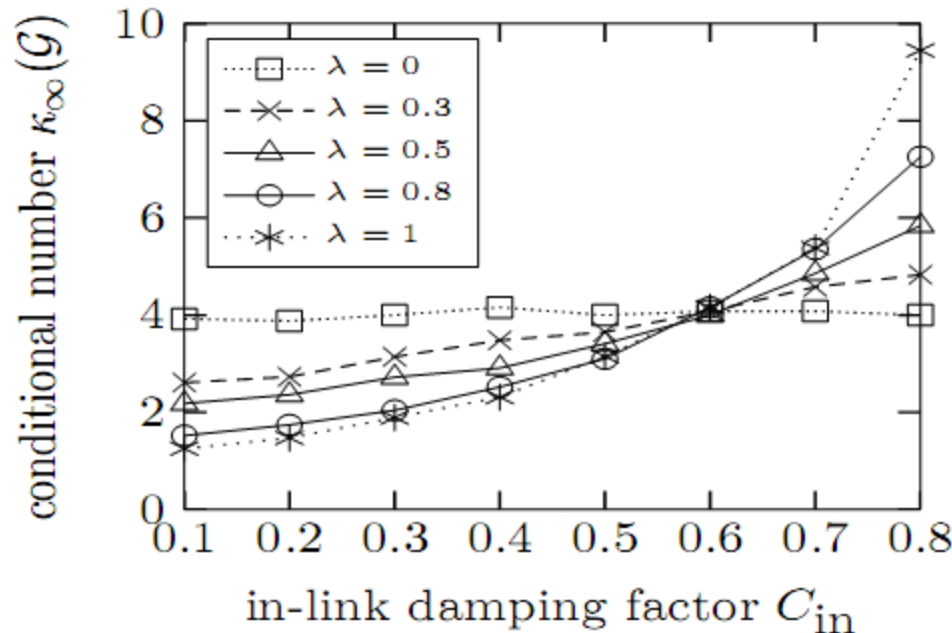
(a) weighting factor λ w.r.t. κ_{∞}

(g) weighting factor γ w.r.t. κ_{∞}

Increasing λ induces a large P-Rank conditional number when $C_{in} > 0.6$.
When $C_{in} < 0.6$, $\kappa_{\infty}(G)$ is decreased as λ grows.

Experimental Results

- P-Rank Stability

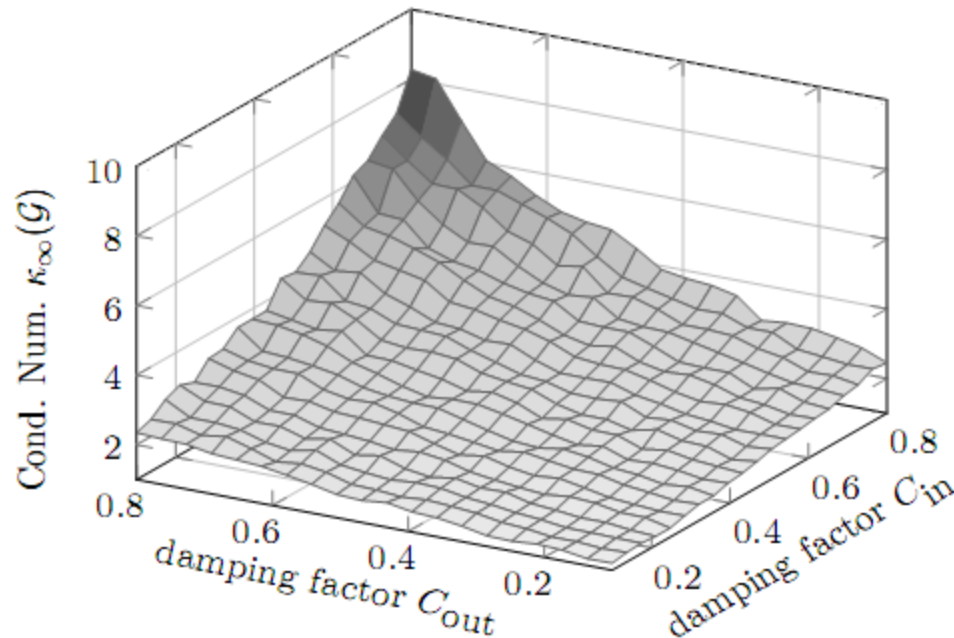


(b) damping factor C_{in} w.r.t. κ_{∞}

When $\lambda = 0$, the curve approaches to a horizontal line. These indicate that varying C_{in} as $\lambda = 0$ has no effect on the stability κ_{∞} of P-Rank, for in this case only the contribution of out-links is considered for computing P-Rank similarity.

Experimental Results

- P-Rank Stability

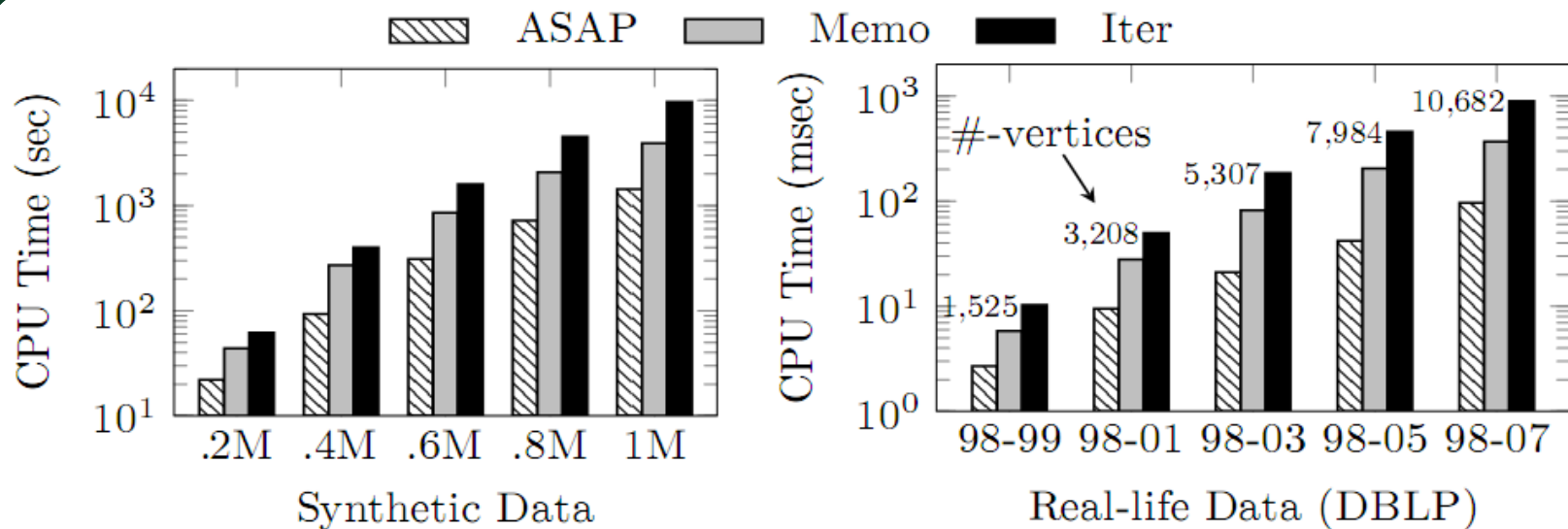


(c) damping factors C_{in} , C_{out} w.r.t. κ_∞

The result demonstrates that P-Rank is comparatively stable when both C_{in} and C_{out} are small (less than 0.6). When C_{in} and $C_{out} \rightarrow 1$, P-Rank is ill-conditioned since small perturbations in similarity computation may cause P-Rank scores drastically altered.

Experimental Results

- P-Rank Time Efficiency



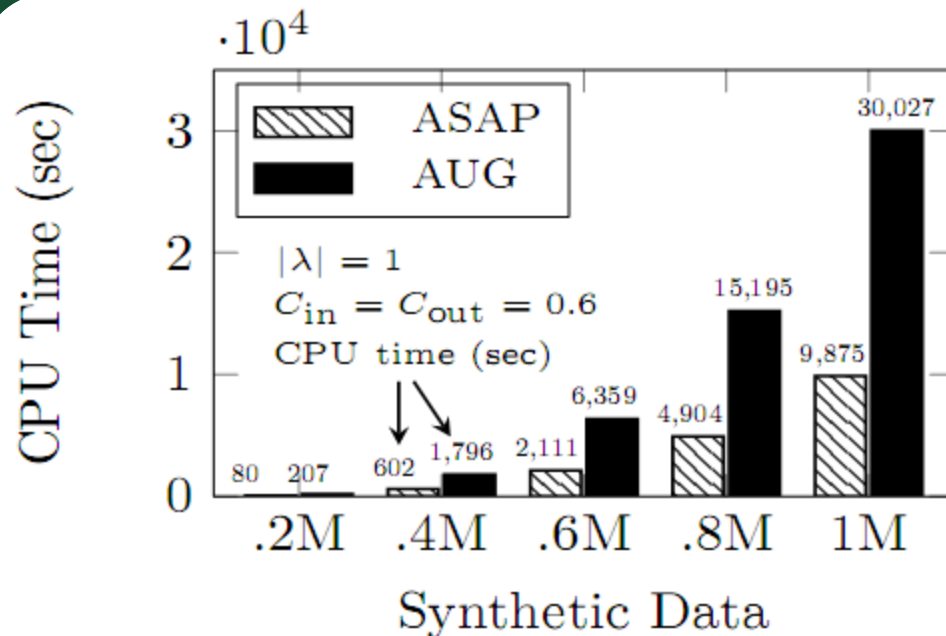
(d) #-vertices n w.r.t. CPU time over synthetic and real-life data

(e) #-vertices n w.r.t. CPU time over synthetic and real-life data

In all cases, ASAP performed the best, by taking advantage of its non-iterative paradigm.

Experimental Results

- P-Rank Time Efficiency



(e) ASAP v.s. AUG on synthetic data

ASAP runs approx. 3x faster than AUG because after eigen-decomposition, AUG still requires extra iterations to be performed in the small eigen-subspace, which takes a significant amount of time, whereas ASAP can straightforwardly compute similarities in terms of eigenvectors with no need for iterations, and therefore takes less time.³³

Conclusions

- An accuracy estimate has been proposed for the P-Rank iterative paradigm, by finding out the exact number of iterations needed to attain a given accuracy.
- The notion of P-Rank conditional number was introduced based on P-Rank matrix representation. A tight bound of P-Rank conditional number was provided to show how the weighting factor and the damping factors affect the P-Rank stability.
- An $O(n^3)$ -time algorithm has been devised to deal with the P-Rank optimization problem over undirected networks.

Thank You!

Q / A ?