

Actuarial Mathematics

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Abstract

Actuarial Mathematics is a theory in applied mathematics, which is mainly used for determining the prices of insurance products and evaluating the liability of a company associating with insurance contracts. It is related to calculus, probability theory and financial theory, etc.

In this entry, I formalize the very basic part of Actuarial Mathematics in Isabelle/HOL. It includes the theory of interest, survival model, and life table. The theory of interest deals with interest rates, present value factors, an annuity certain, etc. The survival model is a probabilistic model that represents the human mortality. The life table is based on the survival model and used for practical calculations.

I have already formalized the basic part of Actuarial Mathematics in Coq (<https://github.com/Yosuke-Ito-345/Actuary>) in a purely axiomatic manner. In contrast, Isabelle formalization is based on the probability theory and the survival model is developed as generally as possible. Such rigorous and general formulation seems very rare; at least I cannot find any similar documentation on the web.

This formalization in Isabelle is still at an early stage, and I cannot guarantee the backward compatibility in the future development. If you heavily depend on the “Actuarial Mathematics” library, please let me know.

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theory *Preliminaries*

imports *HOL-Library.Rewrite HOL-Library.Extended-Nonnegative-Real HOL-Library.Extended-Real HOL-Probability.Probability*

begin

declare *[[show-types]]*

notation *powr (infixr .^ 80)*

1 Preliminary Lemmas

lemma *Collect-conj-eq2*: $\{x \in A. P x \wedge Q x\} = \{x \in A. P x\} \cap \{x \in A. Q x\}$
by *blast*

lemma *vimage-compl-atMost*:

fixes $f :: 'a \Rightarrow 'b::linorder$

shows $-(f - \cdot \{..y\}) = f - \cdot \{y<..\}$

by *fastforce*

context *linorder*

begin

lemma *Icc-minus-Ico*:

fixes $a b$

assumes $a \leq b$

shows $\{a..b\} - \{a..<b\} = \{b\}$

proof

{ fix x **assume** $x \in \{a..b\} - \{a..<b\}$

hence $x \in \{b\}$ **by** *force* **}**

thus $\{a..b\} - \{a..<b\} \subseteq \{b\}$ **by** *blast*

next

show $\{b\} \subseteq \{a..b\} - \{a..<b\}$ **using** *assms* **by** *simp*

qed

lemma *Icc-minus-Ioc*:

fixes $a b$

assumes $a \leq b$

shows $\{a..b\} - \{a<..b\} = \{a\}$
proof
 { **fix** x **assume** $x \in \{a..b\} - \{a<..b\}$
 hence $x \in \{a\}$ **by** *force* }
thus $\{a..b\} - \{a<..b\} \subseteq \{a\}$ **by** *blast*
next
show $\{a\} \subseteq \{a..b\} - \{a<..b\}$ **using** *assms* **by** *simp*
qed

lemma *Int-atLeastAtMost-Unbounded*[*simp*]: $\{a..\} \text{ Int } \{..b\} = \{a..b\}$
by *auto*

lemma *Int-greaterThanAtMost-Unbounded*[*simp*]: $\{a<..\} \text{ Int } \{..b\} = \{a<..b\}$
by *auto*

lemma *Int-atLeastLessThan-Unbounded*[*simp*]: $\{a..\} \text{ Int } \{..<b\} = \{a..<b\}$
by *auto*

lemma *Int-greaterThanLessThan-Unbounded*[*simp*]: $\{a<..\} \text{ Int } \{..<b\} = \{a<..<b\}$
by *auto*

end

lemma *Ico-real-nat-disjoint*:
disjoint-family $(\lambda n::\text{nat}. \{a + \text{real } n ..< a + \text{real } n + 1\})$ **for** $a::\text{real}$
proof –
 { **fix** $m\ n :: \text{nat}$
 assume $\{a + \text{real } m ..< a + \text{real } m + 1\} \cap \{a + \text{real } n ..< a + \text{real } n + 1\}$
 $\neq \{\}$
 then obtain $x::\text{real}$
 where $x \in \{a + \text{real } m ..< a + \text{real } m + 1\} \cap \{a + \text{real } n ..< a + \text{real } n + 1\}$ **by** *force*
 hence $m = n$ **by** *simp* }
thus *?thesis* **unfolding** *disjoint-family-on-def* **by** *blast*
qed

corollary *Ico-nat-disjoint*: *disjoint-family* $(\lambda n::\text{nat}. \{\text{real } n ..< \text{real } n + 1\})$
using *Ico-real-nat-disjoint*[*of 0*] **by** *simp*

lemma *Ico-real-nat-union*: $(\bigcup n::\text{nat}. \{a + \text{real } n ..< a + \text{real } n + 1\}) = \{a..\}$
for $a::\text{real}$
proof
show $(\bigcup n::\text{nat}. \{a + \text{real } n ..< a + \text{real } n + 1\}) \subseteq \{a..\}$ **by** *auto*
next
show $\{a..\} \subseteq (\bigcup n::\text{nat}. \{a + \text{real } n ..< a + \text{real } n + 1\})$
proof
fix x **assume** $x \in \{a..\}$
hence $a \leq x$ **by** *simp*

hence $\text{nat } \lfloor x-a \rfloor \leq x-a \wedge x-a < \text{nat } \lfloor x-a \rfloor + 1$ **by** *linarith*
 hence $a + \text{nat } \lfloor x-a \rfloor \leq x \wedge x < a + \text{nat } \lfloor x-a \rfloor + 1$ **by** *auto*
 thus $x \in (\bigcup n::\text{nat}. \{a + \text{real } n ..< a + \text{real } n + 1\})$ **by** *auto*
qed
qed

corollary *Ico-nat-union*: $(\bigcup n::\text{nat}. \{\text{real } n ..< \text{real } n + 1\}) = \{0..\}$
using *Ico-real-nat-union[of 0]* **by** *simp*

lemma *Ico-nat-union-finite*: $(\bigcup (n::\text{nat})<m. \{\text{real } n ..< \text{real } n + 1\}) = \{0..<m\}$

proof

show $(\bigcup (n::\text{nat})<m. \{\text{real } n ..< \text{real } n + 1\}) \subseteq \{0..<m\}$ **by** *auto*

next

show $\{0..<m\} \subseteq (\bigcup (n::\text{nat})<m. \{\text{real } n ..< \text{real } n + 1\})$

proof

fix $x::\text{real}$

assume $\star: x \in \{0..<m\}$

hence $\text{nat } \lfloor x \rfloor < m$ **using** *of-nat-floor* **by** *fastforce*

moreover with \star **have** $\text{nat } \lfloor x \rfloor \leq x \wedge x < \text{nat } \lfloor x \rfloor + 1$

by (*metis* *Nat.add-0-right* *atLeastLessThan-iff* *le-nat-floor*
less-one *linorder-not-le* *nat-add-left-cancel-le* *of-nat-floor*)

ultimately show $x \in (\bigcup (n::\text{nat})<m. \{\text{real } n ..< \text{real } n + 1\})$

unfolding *atLeastLessThan-def* **by** *force*

qed

qed

lemma *seq-part-multiple*: **fixes** $m n :: \text{nat}$ **assumes** $m \neq 0$ **defines** $A \equiv \lambda i::\text{nat}. \{i*m ..< (i+1)*m\}$

shows $\forall i j. i \neq j \longrightarrow A i \cap A j = \{\}$ **and** $(\bigcup i<n. A i) = \{..< n*m\}$

proof –

{ fix $i j :: \text{nat}$

have $i \neq j \implies A i \cap A j = \{\}$

proof (*erule* *contrapos-np*)

assume $A i \cap A j \neq \{\}$

then obtain k **where** $k \in A i \cap A j$ **by** *blast*

hence $i*m < (j+1)*m \wedge j*m < (i+1)*m$ **unfolding** *A-def* **by** *force*

hence $i < j+1 \wedge j < i+1$ **using** *mult-less-cancel2* **by** *blast*

thus $i = j$ **by** *force*

qed }

thus $\forall i j. i \neq j \longrightarrow A i \cap A j = \{\}$ **by** *blast*

next

show $(\bigcup i<n. A i) = \{..< n*m\}$

proof

show $(\bigcup i<n. A i) \subseteq \{..< n*m\}$

proof

fix $x::\text{nat}$

assume $x \in (\bigcup i<n. A i)$

then obtain i **where** $i:n: i < n$ **and** $i-x: x < (i+1)*m$ **unfolding** *A-def* **by**

force

hence $i+1 \leq n$ **by** *linarith*
hence $x < n*m$ **by** (*meson less-le-trans mult-le-cancel2 i-x*)
thus $x \in \{.. < n*m\}$
using *diff-mult-distrib mult-1 i-n* **by** *auto*
qed
next
show $\{.. < n*m\} \subseteq (\bigcup i < n. A i)$
proof
fix $x::nat$
let $?i = x \text{ div } m$
assume $x \in \{.. < n*m\}$
hence $?i < n$ **by** (*simp add: less-mult-imp-div-less*)
moreover have $?i*m \leq x \wedge x < (?i+1)*m$
using *assms div-times-less-eq-dividend dividend-less-div-times* **by** *auto*
ultimately show $x \in (\bigcup i < n. A i)$ **unfolding** *A-def* **by** *force*
qed
qed
qed

lemma *frontier-Icc-real*: *frontier {a..b} = {a, b} if $a \leq b$ for $a b :: real$*
unfolding *frontier-def* **using** *that* **by** *force*

lemma(*in field*) *divide-mult-cancel[simp]*: **fixes** $a b$ **assumes** $b \neq 0$
shows $a / b * b = a$
by (*simp add: assms*)

lemma *inverse-powr*: $(1/a).\hat{\ }b = a.\hat{\ }-b$ **if** $a > 0$ **for** $a b :: real$
by (*smt that powr-divide powr-minus-divide powr-one-eq-one*)

lemma *powr-eq-one-iff-gen[simp]*: $a.\hat{\ }x = 1 \iff x = 0$ **if** $a > 0$ $a \neq 1$ **for** $a x :: real$
by (*metis powr-eq-0-iff powr-inj powr-zero-eq-one that*)

lemma *powr-less-cancel2*: $0 < a \implies 0 < x \implies 0 < y \implies x.\hat{\ }a < y.\hat{\ }a \implies x < y$

for $a x y :: real$

proof –

assume *a-pos*: $0 < a$ **and** *x-pos*: $0 < x$ **and** *y-pos*: $0 < y$

show $x.\hat{\ }a < y.\hat{\ }a \implies x < y$

proof (*erule contrapos-pp*)

assume $\neg x < y$

hence $x \geq y$ **by** *fastforce*

hence $x.\hat{\ }a \geq y.\hat{\ }a$

proof (*cases x = y*)

case *True*

thus *?thesis* **by** *simp*

next

case *False*

hence $x.\hat{\ }a > y.\hat{\ }a$

using $\langle x \geq y \rangle$ *powr-less-mono2 a-pos y-pos* **by** *auto*
thus *?thesis* **by** *auto*
qed
thus $\neg x.\hat{a} < y.\hat{a}$ **by** *fastforce*
qed
qed

lemma *geometric-increasing-sum-aux*: $(1-r)^{\wedge 2} * (\sum k < n. (k+1)*r^{\wedge k}) = 1 - (n+1)*r^{\wedge n} + n*r^{\wedge (n+1)}$
for $n::\text{nat}$ **and** $r::\text{real}$
proof (*induct n*)
case 0
thus *?case* **by** *simp*
next
case (*Suc n*)
thus *?case*
apply (*simp only: sum.lessThan-Suc*)
apply (*subst distrib-left*)
apply (*subst Suc.hyps*)
by (*subst power2-diff, simp add: field-simps power2-eq-square*)
qed

lemma *geometric-increasing-sum*: $(\sum k < n. (k+1)*r^{\wedge k}) = (1 - (n+1)*r^{\wedge n} + n*r^{\wedge (n+1)}) / (1-r)^{\wedge 2}$
if $r \neq 1$ **for** $n::\text{nat}$ **and** $r::\text{real}$
by (*subst geometric-increasing-sum-aux[THEN sym], simp add: that*)

lemma *Reals-UNIV[simp]*: $\mathbb{R} = \{x::\text{real}. \text{True}\}$
unfolding *Reals-def* **by** *auto*

lemma *Lim-cong*:
assumes $\forall_F x \text{ in } F. f x = g x$
shows $\text{Lim } F f = \text{Lim } F g$
unfolding *t2-space-class.Lim-def* **using** *tendsto-cong assms* **by** *fastforce*

lemma *LIM-zero-iff'*: $((\lambda x. l - f x) \longrightarrow 0) F = (f \longrightarrow l) F$
for $f :: 'a \Rightarrow 'b::\text{real-normed-vector}$
unfolding *tendsto-iff dist-norm*
by (*rewrite minus-diff-eq[THEN sym], rewrite norm-minus-cancel*) *simp*

lemma *antimono-onI*:
 $(\bigwedge r s. r \in A \implies s \in A \implies r \leq s \implies f r \geq f s) \implies \text{antimono-on } A f$
by (*rule monotone-onI*)

lemma *antimono-onD*:
 $\llbracket \text{antimono-on } A f; r \in A; s \in A; r \leq s \rrbracket \implies f r \geq f s$
by (*rule monotone-onD*)

lemma *antimono-imp-mono-on*: $\text{antimono } f \implies \text{antimono-on } A f$

by (simp add: antimononD antimonon-onI)

lemma *antimonon-on-subset*: $\text{antimonon-on } A \ f \implies B \subseteq A \implies \text{antimonon-on } B \ f$
 by (rule monotone-on-subset)

lemma *monon-on-antimonon-on*:

fixes $f :: 'a::\text{order} \Rightarrow 'b::\text{ordered-ab-group-add}$
 shows $\text{monon-on } A \ f \longleftrightarrow \text{antimonon-on } A \ (\lambda r. - f r)$
 by (simp add: monotone-on-def)

corollary *monon-antimonon*:

fixes $f :: 'a::\text{order} \Rightarrow 'b::\text{ordered-ab-group-add}$
 shows $\text{monon } f \longleftrightarrow \text{antimonon } (\lambda r. - f r)$
 by (rule monon-on-antimonon-on)

lemma *monon-on-at-top-le*:

fixes $a :: 'a::\text{linorder}$ and $b :: 'b::\{\text{order-topology, linordered-ab-group-add}\}$
 and $f :: 'a \Rightarrow 'b$
 assumes $f\text{-monon}$: $\text{monon-on } \{a..\} \ f$ and $f\text{-to-l}$: $(f \longrightarrow l) \text{ at-top}$
 shows $\bigwedge x. x \in \{a..\} \implies f x \leq l$
proof (unfold atomize-ball)
 { fix b assume $b\text{-a}$: $b \geq a$ and $fb\text{-l}$: $\neg f b \leq l$
 let $?eps = f b - l$
 have lim-top : $\bigwedge S. \text{open } S \implies l \in S \implies \text{eventually } (\lambda x. f x \in S) \text{ at-top}$
 using *assms tendsto-def* by *auto*
 have eventually $(\lambda x. f x \in \{l - ?eps <..< l + ?eps\}) \text{ at-top}$
 using $fb\text{-l}$ by (intro lim-top ; *force*)
 then obtain N where $fn\text{-in}$: $\bigwedge n. n \geq N \implies f n \in \{l - ?eps <..< l + ?eps\}$
 using $\text{eventually-at-top-linorder}$ by *metis*
 let $?n = \max b N$
 have $f ?n < f b$ using $fn\text{-in}$ by *force*
 moreover have $f ?n \geq f b$ using $f\text{-monon } b\text{-a}$ by (simp add: *le-max-iff-disj*
monon-on-def)
 ultimately have *False* by *simp* }
 thus $\forall x \in \{a..\}. f x \leq l$
 apply -
 by (rule *notnotD*, rewrite *Set.ball-simps*) *auto*
qed

corollary *monon-at-top-le*:

fixes $b :: 'b::\{\text{order-topology, linordered-ab-group-add}\}$ and $f :: 'a::\text{linorder} \Rightarrow 'b$
 assumes $\text{monon } f$ and $(f \longrightarrow b) \text{ at-top}$
 shows $\bigwedge x. f x \leq b$
 using *monon-on-at-top-le assms* by (*metis atLeast-iff monon-imp-monon-on nle-le*)

lemma *monon-on-at-bot-ge*:

fixes $a :: 'a::\text{linorder}$ and $b :: 'b::\{\text{order-topology, linordered-ab-group-add}\}$
 and $f :: 'a \Rightarrow 'b$
 assumes $f\text{-monon}$: $\text{monon-on } \{..a\} \ f$ and $f\text{-to-l}$: $(f \longrightarrow l) \text{ at-bot}$

shows $\bigwedge x. x \in \{..a\} \implies f x \geq l$
proof (*unfold atomize-ball*)
{ **fix** b **assume** $b-a: b \leq a$ **and** $fb-l: \neg f b \geq l$
 let $?eps = l - f b$
 have $lim-bot: \bigwedge S. open S \implies l \in S \implies eventually (\lambda x. f x \in S)$ *at-bot*
 using *assms tendsto-def* **by** *auto*
 have $eventually (\lambda x. f x \in \{l - ?eps <..< l + ?eps\})$ *at-bot*
 using $fb-l$ **by** (*intro lim-bot; force*)
 then obtain N **where** $fn-in: \bigwedge n. n \leq N \implies f n \in \{l - ?eps <..< l + ?eps\}$
 using *eventually-at-bot-linorder* **by** *metis*
 let $?n = \min b N$
 have $f ?n > f b$ **using** $fn-in$ **by** *force*
 moreover have $f ?n \leq f b$ **using** $f-mono\ b-a$ **by** (*simp add: min.coboundedII*
mono-onD)
 ultimately have *False* **by** *simp* }
thus $\forall x \in \{..a\}. f x \geq l$
 apply $-$
 by (*rule notnotD, rewrite Set.ball-simps*) *auto*
qed

corollary *mono-at-bot-ge*:

fixes $b :: 'b::\{order-topology, linordered-ab-group-add\}$ **and** $f :: 'a::linorder \Rightarrow 'b$
assumes *mono f* **and** ($f \longrightarrow b$) *at-bot*
shows $\bigwedge x. f x \geq b$
using *mono-on-at-bot-ge assms* **by** (*metis atMost-iff mono-imp-mono-on nle-le*)

lemma *antimono-on-at-top-ge*:

fixes $a :: 'a::linorder$ **and** $b :: 'b::\{order-topology, linordered-ab-group-add\}$
 and $f :: 'a \Rightarrow 'b$
assumes $f-antimono: antimono-on \{a..\} f$ **and** $f-to-l: (f \longrightarrow l)$ *at-top*
shows $\bigwedge x. x \in \{a..\} \implies f x \geq l$
proof (*unfold atomize-ball*)
{ **fix** b **assume** $b-a: b \geq a$ **and** $fb-l: \neg f b \geq l$
 let $?eps = l - f b$
 have $lim-top: \bigwedge S. open S \implies l \in S \implies eventually (\lambda x. f x \in S)$ *at-top*
 using *assms tendsto-def* **by** *auto*
 have $eventually (\lambda x. f x \in \{l - ?eps <..< l + ?eps\})$ *at-top*
 using $fb-l$ **by** (*intro lim-top; force*)
 then obtain N **where** $fn-in: \bigwedge n. n \geq N \implies f n \in \{l - ?eps <..< l + ?eps\}$
 using *eventually-at-top-linorder* **by** *metis*
 let $?n = \max b N$
 have $f ?n > f b$ **using** $fn-in$ **by** *force*
 moreover have $f ?n \leq f b$ **using** $f-antimono\ b-a$
 by (*simp add: antimono-onD le-max-iff-disj*)
 ultimately have *False* **by** *simp* }
thus $\forall x \in \{a..\}. f x \geq l$
 apply $-$
 by (*rule notnotD, rewrite Set.ball-simps*) *auto*
qed

corollary *antimono-at-top-le*:

fixes $b :: 'b::\{\text{order-topology, linordered-ab-group-add}\}$ **and** $f :: 'a::\text{linorder} \Rightarrow 'b$
assumes *antimono* f **and** $(f \longrightarrow b)$ *at-top*
shows $\bigwedge x. f\ x \geq b$
using *antimono-on-at-top-ge* *assms* *antimono-imp-mono-on* **by** *blast*

lemma *antimono-on-at-bot-ge*:

fixes $a :: 'a::\text{linorder}$ **and** $b :: 'b::\{\text{order-topology, linordered-ab-group-add}\}$
and $f :: 'a \Rightarrow 'b$
assumes $f\text{-antimono}$: *antimono-on* $\{..a\}$ f **and** $f\text{-to-l}$: $(f \longrightarrow l)$ *at-bot*
shows $\bigwedge x. x \in \{..a\} \implies f\ x \leq l$
proof (*unfold* *atomize-ball*)
{ **fix** b **assume** $b\text{-a}$: $b \leq a$ **and** $fb\text{-l}$: $\neg f\ b \leq l$
 let $?eps = f\ b - l$
 have $lim\text{-bot}$: $\bigwedge S. \text{open } S \implies l \in S \implies \text{eventually } (\lambda x. f\ x \in S)$ *at-bot*
 using *assms* *tendsto-def* **by** *auto*
 have $\text{eventually } (\lambda x. f\ x \in \{l - ?eps <..< l + ?eps\})$ *at-bot*
 using $fb\text{-l}$ **by** (*intro* $lim\text{-bot}$; *force*)
 then obtain N **where** $fn\text{-in}$: $\bigwedge n. n \leq N \implies f\ n \in \{l - ?eps <..< l + ?eps\}$
 using *eventually-at-bot-linorder* **by** *metis*
 let $?n = \min\ b\ N$
 have $f\ ?n < f\ b$ **using** $fn\text{-in}$ **by** *force*
 moreover have $f\ ?n \geq f\ b$ **using** $f\text{-antimono}$ $b\text{-a}$ **by** (*simp* *add*: *min.coboundedI1*
antimono-onD)
 ultimately have *False* **by** *simp* }
thus $\forall x \in \{..a\}. f\ x \leq l$
apply –
by (*rule* *notnotD*, *rewrite* *Set.ball-simps*) *auto*

qed

corollary *antimono-at-bot-ge*:

fixes $b :: 'b::\{\text{order-topology, linordered-ab-group-add}\}$ **and** $f :: 'a::\text{linorder} \Rightarrow 'b$
assumes *antimono* f **and** $(f \longrightarrow b)$ *at-bot*
shows $\bigwedge x. f\ x \leq b$
using *antimono-on-at-bot-ge* *assms* *antimono-imp-mono-on* **by** *blast*

lemma *continuous-cdivide*:

fixes $c :: 'a::\text{real-normed-field}$
assumes $c \neq 0$ *continuous* $F\ f$
shows *continuous* F $(\lambda x. f\ x / c)$
using *assms* *continuous-mult-right* **by** (*rewrite* *field-class.field-divide-inverse*)
auto

lemma *continuous-mult-left-iff*:

fixes $c :: 'a::\text{real-normed-field}$
assumes $c \neq 0$
shows *continuous* $F\ f \iff \text{continuous } F$ $(\lambda x. c * f\ x)$
using *continuous-mult-left* *continuous-cdivide* *assms* **by** *force*

lemma *continuous-mult-right-iff*:
fixes $c :: 'a :: \text{real-normed-field}$
assumes $c \neq 0$
shows $\text{continuous } F f \longleftrightarrow \text{continuous } F (\lambda x. f x * c)$
using *continuous-mult-right continuous-cdivide assms* **by force**

lemma *continuous-cdivide-iff*:
fixes $c :: 'a :: \text{real-normed-field}$
assumes $c \neq 0$
shows $\text{continuous } F f \longleftrightarrow \text{continuous } F (\lambda x. f x / c)$
proof
assume $\text{continuous } F f$
thus $\text{continuous } F (\lambda x. f x / c)$
by (*intro continuous-cdivide*) (*simp add: assms*)
next
assume $\text{continuous } F (\lambda x. f x / c)$
hence $\text{continuous } F (\lambda x. f x / c * c)$ **using** *continuous-mult-right* **by fastforce**
thus $\text{continuous } F f$ **using** *assms* **by force**
qed

lemma *continuous-cong*:
assumes $\text{eventually } (\lambda x. f x = g x) F f (\text{Lim } F (\lambda x. x)) = g (\text{Lim } F (\lambda x. x))$
shows $\text{continuous } F f \longleftrightarrow \text{continuous } F g$
unfolding *continuous-def* **using** *assms filterlim-cong* **by force**

lemma *continuous-at-within-cong*:
assumes $f x = g x \text{ eventually } (\lambda x. f x = g x) (\text{at } x \text{ within } s)$
shows $\text{continuous } (\text{at } x \text{ within } s) f \longleftrightarrow \text{continuous } (\text{at } x \text{ within } s) g$
proof (*cases* $\langle x \in \text{closure } (s - \{x\}) \rangle$)
case *True*
thus *?thesis*
using *assms apply* (*intro continuous-cong, simp*)
by (*rewrite Lim-ident-at, simp add: at-within-eq-bot-iff*) **+** *simp*
next
case *False*
hence $\text{at } x \text{ within } s = \text{bot}$ **using** *not-trivial-limit-within* **by blast**
thus *?thesis* **by simp**
qed

lemma *continuous-within-shift*:
fixes $a x :: 'a :: \{\text{topological-ab-group-add, } t2\text{-space}\}$
and $s :: 'a \text{ set}$
and $f :: 'a \Rightarrow 'b :: \text{topological-space}$
shows $\text{continuous } (\text{at } x \text{ within } s) (\lambda x. f (x+a)) \longleftrightarrow \text{continuous } (\text{at } (x+a) \text{ within } \text{plus } a ' s) f$
proof
assume $\text{continuous } (\text{at } x \text{ within } s) (\lambda x. f (x+a))$
moreover **have** $\text{continuous } (\text{at } (x+a) \text{ within } \text{plus } a ' s) (\text{plus } (-a))$

by (simp add: continuous-at-imp-continuous-at-within)
 moreover have plus $(-a)$ ' plus a ' $s = s$ by force
 ultimately show continuous (at $(x+a)$ within plus a ' s) f
 using continuous-within-compose unfolding comp-def by force
 next
 assume continuous (at $(x+a)$ within plus a ' s) f
 moreover have continuous (at x within s) (plus a)
 by (simp add: continuous-at-imp-continuous-at-within)
 ultimately show continuous (at x within s) $(\lambda x. f (x+a))$
 using continuous-within-compose unfolding comp-def by (force simp add:
 add.commute)
 qed

lemma *isCont-shift*:
 fixes $a x :: 'a :: \{\text{topological-ab-group-add, } t2\text{-space}\}$
 and $f :: 'a \Rightarrow 'b::\text{topological-space}$
 shows *isCont* $(\lambda x. f (x+a)) x \longleftrightarrow \text{isCont } f (x+a)$
 using continuous-within-shift by force

lemma *has-real-derivative-at-split*:
 $(f \text{ has-real-derivative } D) (at x) \longleftrightarrow$
 $(f \text{ has-real-derivative } D) (at-left x) \wedge (f \text{ has-real-derivative } D) (at-right x)$
proof –
 have $at x = at x \text{ within } (\{..<x\} \cup \{x<..\})$ by (simp add: at-eq-sup-left-right
 at-within-union)
 thus $(f \text{ has-real-derivative } D) (at x) \longleftrightarrow$
 $(f \text{ has-real-derivative } D) (at-left x) \wedge (f \text{ has-real-derivative } D) (at-right x)$
 using *Lim-within-Un has-field-derivative-iff* by fastforce
 qed

lemma *DERIV-cmult-iff*:
 assumes $c \neq 0$
 shows $(f \text{ has-field-derivative } D) (at x \text{ within } s) \longleftrightarrow$
 $((\lambda x. c * f x) \text{ has-field-derivative } c * D) (at x \text{ within } s)$
proof
 assume $(f \text{ has-field-derivative } D) (at x \text{ within } s)$
 thus $((\lambda x. c * f x) \text{ has-field-derivative } c * D) (at x \text{ within } s)$ using *DERIV-cmult*
 by force
 next
 assume $((\lambda x. c * f x) \text{ has-field-derivative } c * D) (at x \text{ within } s)$
 hence $((\lambda x. c * f x / c) \text{ has-field-derivative } c * D / c) (at x \text{ within } s)$
 using *DERIV-cdivide assms* by blast
 thus $(f \text{ has-field-derivative } D) (at x \text{ within } s)$ by (simp add: assms field-simps)
 qed

lemma *DERIV-cmult-right-iff*:
 assumes $c \neq 0$
 shows $(f \text{ has-field-derivative } D) (at x \text{ within } s) \longleftrightarrow$
 $((\lambda x. f x * c) \text{ has-field-derivative } D * c) (at x \text{ within } s)$

by (rewrite DERIV-cmult-iff[of c], simp-all add: assms mult-ac)

lemma DERIV-cdivide-iff:

assumes $c \neq 0$

shows (f has-field-derivative D) (at x within s) \longleftrightarrow

(($\lambda x. f x / c$) has-field-derivative D / c) (at x within s)

apply (rewrite field-class.field-divide-inverse)+

using DERIV-cmult-right-iff assms inverse-nonzero-iff-nonzero by blast

lemma DERIV-ln-divide-chain:

fixes $f :: \text{real} \Rightarrow \text{real}$

assumes $f x > 0$ and (f has-real-derivative D) (at x within s)

shows (($\lambda x. \ln (f x)$) has-real-derivative (D / f x)) (at x within s)

proof –

have DERIV $\ln (f x) :> 1 / f x$ using assms by (intro DERIV-ln-divide) blast

thus ?thesis using DERIV-chain2 assms by fastforce

qed

lemma inverse-fun-has-integral-ln:

fixes $f :: \text{real} \Rightarrow \text{real}$ and $f' :: \text{real} \Rightarrow \text{real}$

assumes $a \leq b$ and

$\bigwedge x. x \in \{a..b\} \implies f x > 0$ and

continuous-on $\{a..b\}$ f and

$\bigwedge x. x \in \{a <..<b\} \implies (f \text{ has-real-derivative } f' x) \text{ (at } x)$

shows (($\lambda x. f' x / f x$) has-integral ($\ln (f b) - \ln (f a)$)) $\{a..b\}$

proof –

have continuous-on $\{a..b\}$ ($\lambda x. \ln (f x)$) using assms by (intro continuous-intros; force)

moreover have $\bigwedge x. x \in \{a <..<b\} \implies ((\lambda x. \ln (f x)) \text{ has-vector-derivative } f' x / f x) \text{ (at } x)$

apply (rewrite has-real-derivative-iff-has-vector-derivative[THEN sym])

using assms by (intro DERIV-ln-divide-chain; simp)

ultimately show ?thesis using assms by (intro fundamental-theorem-of-calculus-interior; simp)

qed

lemma DERIV-fun-powr2:

fixes $a :: \text{real}$

assumes a-pos: $a > 0$

and f: DERIV $f x :> r$

shows DERIV ($\lambda x. a. \wedge (f x)$) $x :> a. \wedge (f x) * r * \ln a$

proof –

let ?g = ($\lambda x. a$)

have g: DERIV ?g $x :> 0$ by simp

have pos: ?g $x > 0$ by (simp add: a-pos)

show ?thesis

using DERIV-powr[OF g pos f] a-pos by (auto simp add: field-simps)

qed

lemma *has-real-derivative-powr2*:
assumes $a\text{-pos}$: $a > 0$
shows $((\lambda x. a. \hat{x}) \text{ has-real-derivative } a. \hat{x} * \ln a)$ $(at\ x)$
proof –
let $?f = (\lambda x. x::real)$
have f : $DERIV\ ?f\ x\ :>\ 1$ **by** *simp*
thus $?thesis$ **using** $DERIV\text{-fun-powr2}[OF\ a\text{-pos}\ f]$ **by** *simp*
qed

lemma *field-differentiable-shift*:
 $(f \text{ field-differentiable } (at\ (x + z))) = ((\lambda x. f\ (x + z)) \text{ field-differentiable } (at\ x))$
unfolding *field-differentiable-def* **using** $DERIV\text{-shift}$ **by** *force*

1.1 Lemmas on *indicator* for a Linearly Ordered Type

lemma *indicator-Icc-shift*:
fixes $a\ b\ t\ x :: 'a::linordered-ab-group-add$
shows $indicator\ \{a..b\}\ x = indicator\ \{t+a..t+b\}\ (t+x)$
by $(simp\ add:\ indicator\text{-def})$

lemma *indicator-Icc-shift-inverse*:
fixes $a\ b\ t\ x :: 'a::linordered-ab-group-add$
shows $indicator\ \{a-t..b-t\}\ x = indicator\ \{a..b\}\ (t+x)$
by $(metis\ add.\ commute\ diff\text{-add}\text{-cancel}\ indicator\text{-Icc}\text{-shift})$

lemma *indicator-Ici-shift*:
fixes $a\ t\ x :: 'a::linordered-ab-group-add$
shows $indicator\ \{a..\}\ x = indicator\ \{t+a..\}\ (t+x)$
by $(simp\ add:\ indicator\text{-def})$

lemma *indicator-Ici-shift-inverse*:
fixes $a\ t\ x :: 'a::linordered-ab-group-add$
shows $indicator\ \{a-t..\}\ x = indicator\ \{a..\}\ (t+x)$
by $(metis\ add.\ commute\ diff\text{-add}\text{-cancel}\ indicator\text{-Ici}\text{-shift})$

lemma *indicator-Iic-shift*:
fixes $b\ t\ x :: 'a::linordered-ab-group-add$
shows $indicator\ \{..b\}\ x = indicator\ \{..t+b\}\ (t+x)$
by $(simp\ add:\ indicator\text{-def})$

lemma *indicator-Iic-shift-inverse*:
fixes $b\ t\ x :: 'a::linordered-ab-group-add$
shows $indicator\ \{..b-t\}\ x = indicator\ \{..b\}\ (t+x)$
by $(metis\ add.\ commute\ diff\text{-add}\text{-cancel}\ indicator\text{-Iic}\text{-shift})$

lemma *indicator-Icc-reverse*:
fixes $a\ b\ t\ x :: 'a::linordered-ab-group-add$
shows $indicator\ \{a..b\}\ x = indicator\ \{t-b..t-a\}\ (t-x)$

by (metis add-le-cancel-left atLeastAtMost-iff diff-add-cancel indicator-simps le-diff-eq)

lemma *indicator-Icc-reverse-inverse*:

fixes $a\ b\ t\ x :: 'a::\text{linordered-ab-group-add}$

shows $\text{indicator } \{t-b..t-a\} x = \text{indicator } \{a..b\} (t-x)$

by (metis add-diff-cancel-left' diff-add-cancel indicator-Icc-reverse)

lemma *indicator-Ici-reverse*:

fixes $a\ t\ x :: 'a::\text{linordered-ab-group-add}$

shows $\text{indicator } \{a..x\} = \text{indicator } \{..t-a\} (t-x)$

by (simp add: indicator-def)

lemma *indicator-Ici-reverse-inverse*:

fixes $b\ t\ x :: 'a::\text{linordered-ab-group-add}$

shows $\text{indicator } \{t-b..x\} = \text{indicator } \{..b\} (t-x)$

by (metis add-diff-cancel-left' diff-add-cancel indicator-Ici-reverse)

lemma *indicator-Iic-reverse*:

fixes $b\ t\ x :: 'a::\text{linordered-ab-group-add}$

shows $\text{indicator } \{..b\} x = \text{indicator } \{t-b..x\} (t-x)$

by (simp add: indicator-def)

lemma *indicator-Iic-reverse-inverse*:

fixes $a\ t\ x :: 'a::\text{linordered-field}$

shows $\text{indicator } \{..t-a\} x = \text{indicator } \{a..x\} (t-x)$

by (simp add: indicator-def)

lemma *indicator-Icc-affine-pos*:

fixes $a\ b\ c\ t\ x :: 'a::\text{linordered-field}$

assumes $c > 0$

shows $\text{indicator } \{a..b\} x = \text{indicator } \{t+c*a..t+c*b\} (t + c*x)$

unfolding indicator-def **using** assms **by** simp

lemma *indicator-Icc-affine-pos-inverse*:

fixes $a\ b\ c\ t\ x :: 'a::\text{linordered-field}$

assumes $c > 0$

shows $\text{indicator } \{(a-t)/c..(b-t)/c\} x = \text{indicator } \{a..b\} (t + c*x)$

using indicator-Icc-affine-pos[**where** $a=(a-t)/c$ **and** $b=(b-t)/c$ **and** $c=c$ **and** $t=t$] **assms** **by** simp

lemma *indicator-Ici-affine-pos*:

fixes $a\ c\ t\ x :: 'a::\text{linordered-field}$

assumes $c > 0$

shows $\text{indicator } \{a..x\} = \text{indicator } \{t+c*a..x\} (t + c*x)$

unfolding indicator-def **using** assms **by** simp

lemma *indicator-Ici-affine-pos-inverse*:

fixes $a\ c\ t\ x :: 'a::\text{linordered-field}$

assumes $c > 0$

shows $\text{indicator } \{(a-t)/c..\} x = \text{indicator } \{a..\} (t + c*x)$
using $\text{indicator-Ici-affine-pos}$ [**where** $a=(a-t)/c$ **and** $c=c$ **and** $t=t$] **assms by simp**

lemma $\text{indicator-Iic-affine-pos}$:
fixes $b\ c\ t\ x :: 'a::\text{linordered-field}$
assumes $c > 0$
shows $\text{indicator } \{..b\} x = \text{indicator } \{..t+c*b\} (t + c*x)$
unfolding indicator-def **using** assms **by simp**

lemma $\text{indicator-Iic-affine-pos-inverse}$:
fixes $b\ c\ t\ x :: 'a::\text{linordered-field}$
assumes $c > 0$
shows $\text{indicator } \{..(b-t)/c\} x = \text{indicator } \{..b\} (t + c*x)$
using $\text{indicator-Iic-affine-pos}$ [**where** $b=(b-t)/c$ **and** $c=c$ **and** $t=t$] **assms by simp**

lemma $\text{indicator-Icc-affine-neg}$:
fixes $a\ b\ c\ t\ x :: 'a::\text{linordered-field}$
assumes $c < 0$
shows $\text{indicator } \{a..b\} x = \text{indicator } \{t+c*b..t+c*a\} (t + c*x)$
unfolding indicator-def **using** assms **by auto**

lemma $\text{indicator-Icc-affine-neg-inverse}$:
fixes $a\ b\ c\ t\ x :: 'a::\text{linordered-field}$
assumes $c < 0$
shows $\text{indicator } \{(b-t)/c..(a-t)/c\} x = \text{indicator } \{a..b\} (t + c*x)$
using $\text{indicator-Icc-affine-neg}$ [**where** $a=(b-t)/c$ **and** $b=(a-t)/c$ **and** $c=c$ **and** $t=t$] **assms by simp**

lemma $\text{indicator-Ici-affine-neg}$:
fixes $a\ c\ t\ x :: 'a::\text{linordered-field}$
assumes $c < 0$
shows $\text{indicator } \{a..\} x = \text{indicator } \{..t+c*a\} (t + c*x)$
unfolding indicator-def **using** assms **by simp**

lemma $\text{indicator-Ici-affine-neg-inverse}$:
fixes $b\ c\ t\ x :: 'a::\text{linordered-field}$
assumes $c < 0$
shows $\text{indicator } \{(b-t)/c..\} x = \text{indicator } \{..b\} (t + c*x)$
using $\text{indicator-Ici-affine-neg}$ [**where** $a=(b-t)/c$ **and** $c=c$ **and** $t=t$] **assms by simp**

lemma $\text{indicator-Iic-affine-neg}$:
fixes $b\ c\ t\ x :: 'a::\text{linordered-field}$
assumes $c < 0$
shows $\text{indicator } \{..b\} x = \text{indicator } \{t+c*b..\} (t + c*x)$
unfolding indicator-def **using** assms **by simp**

lemma *indicator-Iic-affine-neg-inverse*:
fixes $a\ c\ t\ x :: 'a::\text{linordered-field}$
assumes $c < 0$
shows $\text{indicator } \{..(a-t)/c\} x = \text{indicator } \{a..\} (t + c*x)$
using *indicator-Iic-affine-neg*[**where** $b=(a-t)/c$ **and** $c=c$ **and** $t=t$] *assms* **by**
simp

2 Additional Lemmas for the *HOL-Analysis* Library

lemma *differentiable-eq-field-differentiable-real*:
fixes $f :: \text{real} \Rightarrow \text{real}$
shows $f \text{ differentiable } F \longleftrightarrow f \text{ field-differentiable } F$
unfolding *field-differentiable-def differentiable-def has-real-derivative*
using *has-real-derivative-iff* **by** *presburger*

lemma *differentiable-on-eq-field-differentiable-real*:
fixes $f :: \text{real} \Rightarrow \text{real}$
shows $f \text{ differentiable-on } s \longleftrightarrow (\forall x \in s. f \text{ field-differentiable (at } x \text{ within } s))$
unfolding *differentiable-on-def* **using** *differentiable-eq-field-differentiable-real* **by**
simp

lemma *differentiable-on-cong* :
assumes $\bigwedge x. x \in s \implies f\ x = g\ x$ **and** $f \text{ differentiable-on } s$
shows $g \text{ differentiable-on } s$
proof –
{ **fix** x **assume** $x \in s$
hence $f \text{ differentiable at } x \text{ within } s$ **using** *assms* **unfolding** *differentiable-on-def*
by *simp*
from this obtain D **where** ($f \text{ has-derivative } D$) (at x within s)
unfolding *differentiable-def* **by** *blast*
hence ($g \text{ has-derivative } D$) (at x within s)
using *has-derivative-transform assms* $\langle x \in s \rangle$ **by** *metis*
hence $g \text{ differentiable at } x \text{ within } s$ **unfolding** *differentiable-def* **by** *blast* }
hence $\forall x \in s. g \text{ differentiable at } x \text{ within } s$ **by** *simp*
thus *?thesis* **unfolding** *differentiable-on-def* **by** *simp*
qed

lemma *C1-differentiable-imp-deriv-continuous-on*:
 $f \text{ C1-differentiable-on } S \implies \text{continuous-on } S (deriv\ f)$
using *C1-differentiable-on-eq field-derivative-eq-vector-derivative* **by** *auto*

lemma *deriv-shift*:
assumes $f \text{ field-differentiable at } (x+a)$
shows $deriv\ f\ (x+a) = deriv\ (\lambda s. f\ (x+s))\ a$
proof –
have ($f \text{ has-field-derivative } deriv\ f\ (x+a)$) (at $(x+a)$)
using *DERIV-deriv-iff-field-differentiable* *assms*
by *force*
hence $((\lambda s. f\ (x+s)) \text{ has-field-derivative } deriv\ f\ (x+a))$ (at a)

using *DERIV-at-within-shift has-field-derivative-at-within* **by** *blast*
moreover have $((\lambda s. f (x+s)) \text{ has-field-derivative } \text{deriv } (\lambda s. f (x+s)) a) (at a)$
using *DERIV-imp-deriv calculation* **by** *fastforce*
ultimately show *?thesis* **using** *DERIV-unique* **by** *force*
qed

lemma *piecewise-differentiable-on-cong*:
assumes *f piecewise-differentiable-on i*
and $\bigwedge x. x \in i \implies f x = g x$
shows *g piecewise-differentiable-on i*
proof –
have *continuous-on i g*
using *continuous-on-cong-simp assms piecewise-differentiable-on-imp-continuous-on*
by *force*
moreover have $\exists S. \text{finite } S \wedge (\forall x \in i - S. g \text{ differentiable } (at x \text{ within } i))$
proof –
from *assms piecewise-differentiable-on-def*
obtain *S where fin: finite S and $\forall x \in i - S. f \text{ differentiable } (at x \text{ within } i)$*
by *metis*
hence $\bigwedge x. x \in i - S \implies f \text{ differentiable } (at x \text{ within } i)$ **by** *simp*
hence $\bigwedge x. x \in i - S \implies g \text{ differentiable } (at x \text{ within } i)$
using *has-derivative-transform assms* **by** *(metis DiffD1 differentiable-def)*
with fin show *?thesis* **by** *blast*
qed
ultimately show *?thesis* **unfolding** *piecewise-differentiable-on-def* **by** *simp*
qed

lemma *differentiable-piecewise*:
assumes *continuous-on i f*
and *f differentiable-on i*
shows *f piecewise-differentiable-on i*
unfolding *piecewise-differentiable-on-def* **using** *assms differentiable-onD* **by** *auto*

lemma *piecewise-differentiable-scaleR*:
assumes *f piecewise-differentiable-on S*
shows $(\lambda x. a *_{\mathbb{R}} f x) \text{ piecewise-differentiable-on } S$
proof –
from *assms* **obtain** *T where fin: finite T $\bigwedge x. x \in S - T \implies f \text{ differentiable at } x \text{ within } S$*
unfolding *piecewise-differentiable-on-def* **by** *blast*
hence $\bigwedge x. x \in S - T \implies (\lambda x. a *_{\mathbb{R}} f x) \text{ differentiable at } x \text{ within } S$
using *differentiable-scaleR* **by** *simp*
moreover have *continuous-on S $(\lambda x. a *_{\mathbb{R}} f x)$*
using *assms continuous-on-scaleR continuous-on-const piecewise-differentiable-on-def*
by *blast*
ultimately show $(\lambda x. a *_{\mathbb{R}} f x) \text{ piecewise-differentiable-on } S$
using *fin piecewise-differentiable-on-def* **by** *blast*
qed

lemma *differentiable-on-piecewise-compose*:
assumes f *piecewise-differentiable-on* S
and g *differentiable-on* $f \text{ ' } S$
shows $g \circ f$ *piecewise-differentiable-on* S
proof –
from *assms* **obtain** T **where** $\text{fin: finite } T \wedge x. x \in S - T \implies f$ *differentiable at } x \text{ within } S
unfolding *piecewise-differentiable-on-def* **by** *blast*
hence $\wedge x. x \in S - T \implies g \circ f$ *differentiable at } x \text{ within } S
by (*meson DiffD1 assms differentiable-chain-within differentiable-onD image-eqI*)
hence $\exists T. \text{finite } T \wedge (\forall x \in S - T. g \circ f$ *differentiable at } x \text{ within } S) **using** *fin*
by *blast*
moreover **have** *continuous-on* S $(g \circ f)$
using *assms continuous-on-compose differentiable-imp-continuous-on*
unfolding *piecewise-differentiable-on-def* **by** *blast*
ultimately show *?thesis*
unfolding *piecewise-differentiable-on-def* **by** *force*
qed***

lemma *MVT-order-free*:
fixes $r \ h :: \text{real}$
defines $I \equiv \{r..r+h\} \cup \{r+h..r\}$
assumes *continuous-on* I f **and** f *differentiable-on interior* I
obtains t **where** $t \in \{0 <..<1\}$ **and** $f(r+h) - f\ r = h * \text{deriv } f(r + t*h)$
proof –
consider (*h-pos*) $h > 0$ | (*h-0*) $h = 0$ | (*h-neg*) $h < 0$ **by** *force*
thus *?thesis*
proof *cases*
case *h-pos*
hence $r < r+h$ $I = \{r..r+h\}$ **unfolding** *I-def* **by** *simp-all*
moreover **hence** *interior* $I = \{r <..<r+h\}$ **by** *simp*
moreover **hence** $\wedge x. \llbracket r < x; x < r+h \rrbracket \implies f$ *differentiable (at } x)*
using *assms by (simp add: differentiable-on-eq-differentiable-at)*
ultimately obtain z **where** $r < z \wedge z < r+h \wedge f(r+h) - f\ r = h * \text{deriv } f\ z$
using *MVT assms by (smt (verit) DERIV-imp-deriv)*
moreover **hence** $(z-r) / h \in \{0 <..<1\}$ **by** *simp*
moreover **have** $z = r + (z-r)/h * h$ **using** *h-pos* **by** *simp*
ultimately show *?thesis* **using** *that* **by** *presburger*
next
case *h-0*
have $1/2 \in \{0::\text{real} <..<1\}$ **by** *simp*
moreover **have** $f(r+h) - f\ r = 0$ **using** *h-0* **by** *simp*
moreover **have** $h * \text{deriv } f(r + (1/2)*h) = 0$ **using** *h-0* **by** *simp*
ultimately show *?thesis* **using** *that* **by** *presburger*
next **case** *h-neg*
hence $r+h < r$ $I = \{r+h..r\}$ **unfolding** *I-def* **by** *simp-all*
moreover **hence** *interior* $I = \{r+h <..<r\}$ **by** *simp*
moreover **hence** $\wedge x. \llbracket r+h < x; x < r \rrbracket \implies f$ *differentiable (at } x)*

using *assms* **by** (*simp add: differentiable-on-eq-differentiable-at*)
ultimately obtain z **where** $r+h < z \wedge z < r \wedge f r - f (r+h) = -h * deriv$
 $f z$
using *MVT assms* **by** (*smt (verit) DERIV-imp-deriv*)
moreover hence $(z-r) / h \in \{0 < .. < 1\}$ **by** (*simp add: neg-less-divide-eq*)
moreover have $z = r + (z-r)/h * h$ **using** *h-neg* **by** *simp*
ultimately show *?thesis* **using** *that mult-minus-left* **by** *fastforce*
qed
qed

lemma *integral-combine2*:
fixes $f :: real \Rightarrow 'a::banach$
assumes $a < c \leq b$
and f *integrable-on* $\{a..c\}$ f *integrable-on* $\{c..b\}$
shows $integral \{a..c\} f + integral \{c..b\} f = integral \{a..b\} f$
apply (*rule integral-unique[THEN sym]*)
apply (*rule has-integral-combine[of a c b], simp-all add: assms*)
using *has-integral-integral assms* **by** *auto*

lemma *has-integral-null-interval*: **fixes** $a b :: real$ **and** $f :: real \Rightarrow real$ **assumes** $a \geq b$
shows $(f \text{ has-integral } 0) \{a..b\}$
using *assms content-real-eq-0* **by** *blast*

lemma *has-integral-interval-reverse*: **fixes** $f :: real \Rightarrow real$ **and** $a b :: real$
assumes $a \leq b$
and *continuous-on* $\{a..b\}$ f
shows $((\lambda x. f (a+b-x)) \text{ has-integral } (integral \{a..b\} f)) \{a..b\}$
proof –
let $?g = \lambda x. a + b - x$
let $?g' = \lambda x. -1$
have g -*C0*: *continuous-on* $\{a..b\}$ $?g$ **using** *continuous-on-op-minus* **by** *simp*
have Dg - g' : $\bigwedge x. x \in \{a..b\} \implies (?g \text{ has-field-derivative } ?g' x) \text{ (at } x \text{ within } \{a..b\})$
by (*auto intro!: derivative-eq-intros*)
show *?thesis*
using *has-integral-substitution-general*
 $[of \{ \} a b ?g a b f, simplified, OF \text{ assms } g\text{-}C0 \text{ } Dg\text{-}g', simplified]$
apply (*simp add: has-integral-null-interval[OF assms(1), THEN integral-unique]*)
by (*simp add: has-integral-neg-iff*)
qed

lemma *FTC-real-deriv-has-integral*:
fixes $F :: real \Rightarrow real$
assumes $a \leq b$
and F *piecewise-differentiable-on* $\{a < .. < b\}$
and *continuous-on* $\{a..b\}$ F
shows $(deriv F \text{ has-integral } F b - F a) \{a..b\}$
proof –
obtain S **where** *fin*: *finite* S **and**

diff: $\bigwedge x. x \in \{a <..<b\} - S \implies F$ differentiable at x within $\{a <..<b\} - S$
using *assms* **unfolding** *piecewise-differentiable-on-def*
by (*meson Diff-subset differentiable-within-subset*)
hence $\bigwedge x. x \in \{a <..<b\} - S \implies (F \text{ has-real-derivative } \text{deriv } F x)$ (at x)
proof –
fix x **assume** $x\text{-in}$: $x \in \{a <..<b\} - S$
have *open* ($\{a <..<b\} - S$)
using *fin finite-imp-closed* **by** (*metis open-Diff open-greaterThanLessThan*)
hence at x within $\{a <..<b\} - S =$ at x **by** (*meson x-in at-within-open*)
hence F differentiable at x **using** *diff x-in* **by** *smt*
thus ($F \text{ has-real-derivative } \text{deriv } F x$) (at x)
using *DERIV-deriv-iff-real-differentiable* **by** *simp*
qed
thus *?thesis*
by (*intro fundamental-theorem-of-calculus-interior-strong*[**where** $S=S$];
simp add: assms fin has-real-derivative-iff-has-vector-derivative)
qed

lemma *integrable-spike-cong*:
assumes *negligible* $S \bigwedge x. x \in T - S \implies g x = f x$
shows f integrable-on $T \longleftrightarrow g$ integrable-on T
using *integrable-spike assms* **by** *force*

lemma *has-integral-powr2-from-0*:
fixes $a c :: \text{real}$
assumes *a-pos*: $a > 0$ **and** *a-neq-1*: $a \neq 1$ **and** *c-nneg*: $c \geq 0$
shows $((\lambda x. a. \hat{x}) \text{ has-integral } ((a. \hat{c} - 1) / (\ln a))) \{0..c\}$
proof –
have $((\lambda x. a. \hat{x}) \text{ has-integral } ((a. \hat{c}) / (\ln a) - (a. \hat{0}) / (\ln a))) \{0..c\}$
proof (*rule fundamental-theorem-of-calculus*[*OF c-nneg*])
fix $x :: \text{real}$
assume $x \in \{0..c\}$
show $((\lambda y. a. \hat{y} / \ln a) \text{ has-vector-derivative } a. \hat{x})$ (at x within $\{0..c\}$)
apply (*insert has-real-derivative-powr2*[*OF a-pos, of x*])
apply (*drule DERIV-cdivide*[**where** $c = \ln a$], *simp add: assms*)
apply (*rule has-vector-derivative-within-subset*[**where** $S=UNIV$ **and** $T=\{0..c\}$],
auto)
by (*rule iffD1*[*OF has-real-derivative-iff-has-vector-derivative*])
qed
thus *?thesis*
using *assms powr-zero-eq-one* **by** (*simp add: field-simps*)
qed

lemma *integrable-on-powr2-from-0*:
fixes $a c :: \text{real}$
assumes *a-pos*: $a > 0$ **and** *a-neq-1*: $a \neq 1$ **and** *c-nneg*: $c \geq 0$
shows $(\lambda x. a. \hat{x})$ integrable-on $\{0..c\}$
using *has-integral-powr2-from-0*[*OF assms*] **unfolding** *integrable-on-def* **by** *blast*

lemma *integrable-on-powr2-from-0-general*:
fixes $a\ c :: \text{real}$
assumes $a\text{-pos}: a > 0$ **and** $c\text{-nneg}: c \geq 0$
shows $(\lambda x. a \cdot \hat{x})$ *integrable-on* $\{0..c\}$
proof (*cases* $a = 1$)
case *True*
thus *?thesis*
using *has-integral-const-real* **by** *auto*
next
case *False*
thus *?thesis*
using *has-integral-powr2-from-0* *False assms* **by** *auto*
qed

lemma *has-bochner-integral-power*:
fixes $a\ b :: \text{real}$ **and** $k :: \text{nat}$
assumes $a \leq b$
shows *has-bochner-integral* *lborel* $(\lambda x. x^k * \text{indicator } \{a..b\} x)$ $((b^{k+1}) - a^{k+1}) / (k+1)$
proof $-$
have $\bigwedge x. ((\lambda x. x^{k+1}) / (k+1))$ *has-real-derivative* x^k (*at* x)
using *DERIV-pow* **by** (*intro derivative-eq-intros*) *auto*
hence *has-bochner-integral* *lborel* $(\lambda x. x^k * \text{indicator } \{a..b\} x)$ $(b^{k+1}) / (k+1)$
 $- a^{k+1} / (k+1)$
by (*intro has-bochner-integral-FTC-Icc-real; simp add: assms*)
thus *?thesis* **by** (*simp add: diff-divide-distrib*)
qed

corollary *integrable-power*: $(a :: \text{real}) \leq b \implies$ *integrable* *lborel* $(\lambda x. x^k * \text{indicator } \{a..b\} x)$
using *has-bochner-integral-power integrable.intros* **by** *blast*

lemma *has-integral-set-integral-real*:
fixes $f :: 'a :: \text{euclidean-space} \Rightarrow \text{real}$ **and** $A :: 'a \text{ set}$
assumes f : *set-integrable* *lborel* A f
shows $(f$ *has-integral* (*set-lebesgue-integral* *lborel* A f)) A
using *assms has-integral-integral-real* [**where** $f = \lambda x. \text{indicat-real } A\ x * f\ x$]
unfolding *set-integrable-def set-lebesgue-integral-def*
by *simp (smt (verit, ccfv-SIG) has-integral-cong has-integral-restrict-UNIV indicator-times-eq-if)*

lemma *set-borel-measurable-lborel*:
set-borel-measurable *lborel* A $f \iff$ *set-borel-measurable* *borel* A f
unfolding *set-borel-measurable-def* **by** *simp*

lemma *restrict-space-whole*[*simp*]: *restrict-space* M (*space* M) = M
unfolding *restrict-space-def* **by** (*simp add: measure-of-of-measure*)

lemma *deriv-measurable-real*:
fixes $f :: \text{real} \Rightarrow \text{real}$
assumes f *differentiable-on* S *open* S $f \in \text{borel-measurable borel}$
shows *set-borel-measurable* $\text{borel } S$ $(\text{deriv } f)$
proof –
have $\bigwedge x. x \in S \implies \text{deriv } f x = \lim (\lambda i. (f (x + 1 / \text{Suc } i) - f x) / (1 / \text{Suc } i))$
proof –
fix x **assume** $x \in S$: $x \in S$
hence f *field-differentiable* (at x within S)
using *differentiable-on-eq-field-differentiable-real* **assms** **by** *simp*
hence f *has-field-derivative* $\text{deriv } f x$ (at x)
using *assms DERIV-deriv-iff-field-differentiable* $x \in S$ *at-within-open* **by** *force*
hence $(\lambda h. (f (x+h) - f x) / h) \xrightarrow{0} \text{deriv } f x$ **using** *DERIV-def* **by** *auto*
hence $\forall d. (\forall i. d i \in \text{UNIV} - \{0 :: \text{real}\}) \longrightarrow d \longrightarrow 0 \longrightarrow$
 $((\lambda h. (f (x+h) - f x) / h) \circ d) \longrightarrow \text{deriv } f x$
using *tendsto-at-iff-sequentially* **by** *blast*
moreover **have** $\forall i. 1 / \text{Suc } i \in \text{UNIV} - \{0 :: \text{real}\}$ **by** *simp*
moreover **have** $(\lambda i. 1 / \text{Suc } i) \longrightarrow 0$ **using** *LIMSEQ-Suc* *lim-const-over-n*
by *blast*
ultimately **have** $((\lambda h. (f (x+h) - f x) / h) \circ (\lambda i. 1 / \text{Suc } i)) \longrightarrow \text{deriv } f x$ **by** *auto*
thus $\text{deriv } f x = \lim (\lambda i. (f (x + 1 / \text{Suc } i) - f x) / (1 / \text{Suc } i))$
unfolding *comp-def* **by** (*simp add: limI*)
qed
moreover **have** $(\lambda x. \text{indicator } S x * \lim (\lambda i. (f (x + 1 / \text{Suc } i) - f x) / (1 / \text{Suc } i)))$
 $\in \text{borel-measurable borel}$
using *assms* **by** (*measurable, simp, measurable*)
ultimately **show** *?thesis*
unfolding *set-borel-measurable-def* *measurable-cong*
by *simp* (*smt* (*verit*) *indicator-simps(2)* *measurable-cong* *mult-eq-0-iff*)
qed

lemma *piecewise-differentiable-on-deriv-measurable-real*:
fixes $f :: \text{real} \Rightarrow \text{real}$
assumes f *piecewise-differentiable-on* S *open* S $f \in \text{borel-measurable borel}$
shows *set-borel-measurable* $\text{borel } S$ $(\text{deriv } f)$
proof –
from *assms* **obtain** T **where** *fin*: *finite* T **and**
 $\text{diff}: \bigwedge x. x \in S - T \implies f$ *differentiable* at x within S
unfolding *piecewise-differentiable-on-def* **by** *blast*
with *assms* **have** *open* $(S - T)$ **using** *finite-imp-closed* **by** *blast*
moreover **hence** f *differentiable-on* $(S - T)$
unfolding *differentiable-on-def* **using** *assms* **by** (*metis* *Diff-iff* *at-within-open* *diff*)
ultimately **have** *set-borel-measurable* $\text{borel } (S - T)$ $(\text{deriv } f)$
by (*intro* *deriv-measurable-real*; *simp add: assms*)

thus *?thesis*
unfolding *set-borel-measurable-def* **apply** *simp*
apply (*rule measurable-discrete-difference*
[**where** $X=T$ **and** $f=\lambda x. \text{indicat-real } (S - T) x * \text{deriv } f x$], *simp-all*)
using *fin uncountable-infinite* **apply** *blast*
by (*simp add: indicator-diff*)
qed

lemma *borel-measurable-antimono*:
fixes $f :: \text{real} \Rightarrow \text{real}$
shows $\text{antimono } f \Longrightarrow f \in \text{borel-measurable borel}$
using *borel-measurable-mono* **by** (*smt (verit, del-insts) borel-measurable-uminus-eq monotone-on-def*)

lemma *set-borel-measurable-restrict-space-iff*:
fixes $f :: 'a \Rightarrow 'b::\text{real-normed-vector}$
assumes $\Omega[\text{measurable, simp}]: \Omega \cap \text{space } M \in \text{sets } M$
shows $f \in \text{borel-measurable } (\text{restrict-space } M \ \Omega) \longleftrightarrow \text{set-borel-measurable } M \ \Omega$
using *assms borel-measurable-restrict-space-iff set-borel-measurable-def* **by** *blast*

lemma *set-integrable-restrict-space-iff*:
fixes $f :: 'a \Rightarrow 'b::\{\text{banach, second-countable-topology}\}$
assumes $A \in \text{sets } M$
shows $\text{set-integrable } M \ A \ f \longleftrightarrow \text{integrable } (\text{restrict-space } M \ A) \ f$
unfolding *set-integrable-def* **using** *assms*
by (*rewrite integrable-restrict-space; simp*)

lemma *set-lebesgue-integral-restrict-space*:
fixes $f :: 'a \Rightarrow 'b::\{\text{banach, second-countable-topology}\}$
assumes $A \in \text{sets } M$
shows $\text{set-lebesgue-integral } M \ A \ f = \text{integral}^L (\text{restrict-space } M \ A) \ f$
unfolding *set-lebesgue-integral-def* **using** *assms integral-restrict-space*
by (*metis (mono-tags) sets.Int-space-eq2*)

lemma *distr-borel-lborel*: $\text{distr } M \ \text{borel } f = \text{distr } M \ \text{lborel } f$
by (*metis distr-cong sets-lborel*)

lemma *AE-translation*:
assumes $\text{AE } x \ \text{in } \text{lborel}. P \ x$ **shows** $\text{AE } x \ \text{in } \text{lborel}. P \ (a+x)$
proof –
from *assms* **obtain** N **where** $P: \bigwedge x. x \in \text{space } \text{lborel} - N \Longrightarrow P \ x$ **and** *null*:
 $N \in \text{null-sets } \text{lborel}$
using *AE-E3* **by** *blast*
hence $\{y. a+y \in N\} \in \text{null-sets } \text{lborel}$
using *null-sets-translation[of N -a, simplified]* **by** (*simp add: add.commute*)
moreover **have** $\bigwedge y. y \in \text{space } \text{lborel} - \{y. a+y \in N\} \Longrightarrow P \ (a+y)$ **using** P
by *force*
ultimately **show** $\text{AE } y \ \text{in } \text{lborel}. P \ (a+y)$

by (smt (verit, del-insts) Diff-iff eventually-ae-filter mem-Collect-eq subsetI)
qed

lemma *set-AE-translation*:

assumes $AE\ x \in S$ in lborel. $P\ x$ shows $AE\ x \in plus\ (-a)\ 'S$ in lborel. $P\ (a+x)$

proof –

have $AE\ x$ in lborel. $a+x \in S \longrightarrow P\ (a+x)$ **using** *assms* **by** (rule *AE-translation*)

moreover have $\bigwedge x. a+x \in S \longleftrightarrow x \in plus\ (-a)\ 'S$ **by** *force*

ultimately show *?thesis* **by** *simp*

qed

lemma *AE-scale-measure-iff*:

assumes $r > 0$

shows $(AE\ x$ in $(scale-measure\ r\ M)$. $P\ x) \longleftrightarrow (AE\ x$ in M . $P\ x)$

unfolding *ae-filter-def null-sets-def*

apply (*rewrite space-scale-measure, simp*)

using *assms* **by** (*smt Collect-cong not-gr-zero*)

lemma *nn-set-integral-cong2*:

assumes $AE\ x \in A$ in M . $f\ x = g\ x$

shows $(\int^{+x \in A}. f\ x\ \partial M) = (\int^{+x \in A}. g\ x\ \partial M)$

proof –

{ **fix** x

assume $x \in space\ M$

have $(x \in A \longrightarrow f\ x = g\ x) \longrightarrow f\ x * indicator\ A\ x = g\ x * indicator\ A\ x$ **by**

force }

hence $AE\ x$ in M . $(x \in A \longrightarrow f\ x = g\ x) \longrightarrow f\ x * indicator\ A\ x = g\ x * indicator\ A\ x$

by (*rule AE-I2*)

hence $AE\ x$ in M . $f\ x * indicator\ A\ x = g\ x * indicator\ A\ x$ **using** *assms* *AE-mp*

by *auto*

thus *?thesis* **by** (*rule nn-integral-cong-AE*)

qed

lemma *set-lebesgue-integral-cong-AE2*:

assumes [*measurable*]: $A \in sets\ M$ *set-borel-measurable* $M\ A$ *set-borel-measurable* $M\ A\ g$

assumes $AE\ x \in A$ in M . $f\ x = g\ x$

shows $(LINT\ x:A|M. f\ x) = (LINT\ x:A|M. g\ x)$

proof –

let $?fA = \lambda x. indicator\ A\ x *_{R}\ f\ x$ **and** $?gA = \lambda x. indicator\ A\ x *_{R}\ g\ x$

have $?fA \in borel-measurable\ M$ $?gA \in borel-measurable\ M$

using *assms* **unfolding** *set-borel-measurable-def* **by** *simp-all*

moreover have $AE\ x \in A$ in M . $?fA\ x = ?gA\ x$ **using** *assms* **by** *simp*

ultimately have $(LINT\ x:A|M. ?fA\ x) = (LINT\ x:A|M. ?gA\ x)$

by (*intro set-lebesgue-integral-cong-AE; simp*)

moreover have $(LINT\ x:A|M. f\ x) = (LINT\ x:A|M. ?fA\ x)$ $(LINT\ x:A|M. g\ x) = (LINT\ x:A|M. ?gA\ x)$

unfolding *set-lebesgue-integral-def*

by (metis indicator-scaleR-eq-if)+
ultimately show ?thesis by simp
qed

proposition *set-nn-integral-eq-set-integral*:

assumes $\forall x \in A$ in M . $0 \leq f x$ set-integrable $M A f$

shows $(\int^{+x \in A}. f x \partial M) = (\int x \in A. f x \partial M)$

proof –

have $(\int^{+x \in A}. f x \partial M) = \int^{+x}. \text{ennreal } (f x * \text{indicator } A x) \partial M$

using *nn-integral-set-ennreal* by blast

also have $\dots = \int x. f x * \text{indicator } A x \partial M$

using *assms unfolding set-integrable-def*

by (rewrite *nn-integral-eq-integral*; force *simp add: mult.commute*)

also have $\dots = (\int x \in A. f x \partial M)$ **unfolding** *set-lebesgue-integral-def* by (*simp add: mult.commute*)

finally show ?thesis .

qed

proposition *nn-integral-disjoint-family-on-finite*:

assumes [*measurable*]: $f \in \text{borel-measurable } M \wedge (n :: \text{nat}). n \in S \implies B n \in \text{sets } M$

and *disjoint-family-on* $B S$ *finite* S

shows $(\int^{+x \in (\bigcup n \in S. B n)}. f x \partial M) = (\sum n \in S. (\int^{+x \in B n}. f x \partial M))$

proof –

let $?A = \lambda n :: \text{nat}. \text{if } n \in S \text{ then } B n \text{ else } \{\}$

have $\bigwedge n :: \text{nat}. ?A n \in \text{sets } M$ by *simp*

moreover have *disjoint-family* $?A$

unfolding *disjoint-family-on-def*

proof –

{ fix $m n :: \text{nat}$

assume $m \neq n$

hence $(\text{if } m \in S \text{ then } B m \text{ else } \{\}) \cap (\text{if } n \in S \text{ then } B n \text{ else } \{\}) = \{\}$

apply *simp*

using *assms unfolding disjoint-family-on-def* by blast }

thus $\forall m :: \text{nat} \in UNIV. \forall n :: \text{nat} \in UNIV. m \neq n \longrightarrow$

$(\text{if } m \in S \text{ then } B m \text{ else } \{\}) \cap (\text{if } n \in S \text{ then } B n \text{ else } \{\}) = \{\}$

by *blast*

qed

ultimately have *set-nn-integral* $M (\bigcup (\text{range } ?A)) f = (\sum n. \text{set-nn-integral } M (?A n) f)$

by (rewrite *nn-integral-disjoint-family*; *simp*)

moreover have *set-nn-integral* $M (\bigcup (\text{range } ?A)) f = (\int^{+x \in (\bigcup n \in S. B n)}. f x \partial M)$

proof –

have $\bigcup (\text{range } ?A) = (\bigcup n \in S. B n)$ by *simp*

thus ?thesis by *simp*

qed

moreover have $(\sum n. \text{set-nn-integral } M (?A n) f) = (\sum n \in S. \text{set-nn-integral } M (B n) f)$

by (rewrite suminf-finite[of S]; simp add: assms)
ultimately show ?thesis by simp
qed

lemma nn-integral-distr-set:

assumes $T \in \text{measurable } M \ M'$ and $f \in \text{borel-measurable } (\text{distr } M \ M' \ T)$
and $A \in \text{sets } M'$ and $\bigwedge x. x \in \text{space } M \implies T \ x \in A$
shows $\text{integral}^N (\text{distr } M \ M' \ T) \ f = \text{set-nn-integral } (\text{distr } M \ M' \ T) \ A \ f$
proof –
have $\text{integral}^N (\text{distr } M \ M' \ T) \ f = (\int^{+x \in (\text{space } M')} . f \ x \ \partial(\text{distr } M \ M' \ T))$
by (rewrite nn-set-integral-space[THEN sym], simp)
also have $\dots = (\int^{+x \in A} . f \ x \ \partial(\text{distr } M \ M' \ T))$
proof –
have [simp]: $\text{sym-diff } (\text{space } M') \ A = \text{space } M' - A$
using assms by (metis Diff-mono sets.sets-into-space sup.orderE)
show ?thesis
apply (rule nn-integral-null-delta; simp add: assms)
unfolding null-sets-def **using** assms
apply (simp, rewrite emeasure-distr; simp)
unfolding vimage-def **using** emeasure-empty
by (smt (z3) Collect-empty-eq Diff-iff Int-def mem-Collect-eq)
qed
finally show ?thesis .
qed

lemma measure-eqI-Ioc:

fixes $M \ N :: \text{real measure}$
assumes sets: $\text{sets } M = \text{sets borel sets } N = \text{borel}$
assumes fin: $\bigwedge a \ b. a \leq b \implies \text{emeasure } M \ \{a <.. b\} < \infty$
assumes eq: $\bigwedge a \ b. a \leq b \implies \text{emeasure } M \ \{a <.. b\} = \text{emeasure } N \ \{a <.. b\}$
shows $M = N$
proof (rule measure-eqI-generator-eq-countable)
let ?Ioc = $\lambda(a::\text{real}, b::\text{real}). \{a <.. b\}$ **let** ?E = $\text{range } ?\text{Ioc}$
show Int-stable ?E **using** Int-stable-def Int-greaterThanAtMost **by** force
show ?E $\subseteq \text{Pow UNIV sets } M = \text{sigma-sets UNIV } ?E \text{ sets } N = \text{sigma-sets UNIV } ?E$
unfolding sets **by** (auto simp add: borel-sigma-sets-Ioc)
show $\bigwedge I. I \in ?E \implies \text{emeasure } M \ I = \text{emeasure } N \ I$
proof –
fix I **assume** $I \in ?E$
then obtain $a \ b$ **where** $I = \{a <.. b\}$ **by** auto
thus $\text{emeasure } M \ I = \text{emeasure } N \ I$ **by** (smt (verit, best) eq greaterThanAtMost-empty)
qed
show ?Ioc $' (Rats \times Rats) \subseteq ?E (\bigcup i \in (Rats \times Rats). ?\text{Ioc } i) = \text{UNIV}$
using Rats-no-bot-less Rats-no-top-le **by** auto
show countable (?Ioc $' (Rats \times Rats)$) **using** countable-rat **by** blast
show $\bigwedge I. I \in ?\text{Ioc } ' (Rats \times Rats) \implies \text{emeasure } M \ I \neq \infty$

proof –
fix I **assume** $I \in ?Ioc \text{ ' } (Rats \times Rats)$
then obtain $a \ b$ **where** $(a, b) \in (Rats \times Rats) \ I = \{a <.. b\}$ **by** *blast*
thus $emeasure \ M \ I \neq \infty$ **by** $(smt \ (verit, \ best) \ Ioc\text{-}inj \ fin \ order.\textit{strict}\text{-}implies\text{-}not\text{-}eq)$
qed
qed

lemma $(in \ finite\text{-}measure) \ distributed\text{-}measure:$

assumes $distributed \ M \ N \ X \ f$
and $\bigwedge x. x \in space \ N \implies f \ x \geq 0$
and $A \in sets \ N$
shows $measure \ M \ (X \ \text{' } A \cap space \ M) = (\int x. indicator \ A \ x \ * \ f \ x \ \partial N)$

proof –

have $[simp]: (\lambda x. indicat\text{-}real \ A \ x \ * \ f \ x) \in borel\text{-}measurable \ N$
using *assms* **apply** $(measurable; \ simp?)$
using $distributed\text{-}real\text{-}measurable \ assms$ **by** *force*
have $emeasure \ M \ (X \ \text{' } A \cap space \ M) = (\int^{+x \in A. ennreal \ (f \ x) \ \partial N)$
by $(rule \ distributed\text{-}emeasure; \ simp \ add: \ assms)$
moreover **have** $enn2real \ (\int^{+x \in A. ennreal \ (f \ x) \ \partial N) = \int x. indicator \ A \ x \ * \ f \ x \ \partial N$
apply $(rewrite \ enn2real\text{-}nn\text{-}integral\text{-}eq\text{-}integral$
 $\ [where \ f = \lambda x. ennreal \ (indicator \ A \ x \ * \ f \ x), \ THEN \ sym]; \ (simp \ add: \ assms)?)$
using $distributed\text{-}emeasure \ assms$
by $(smt \ (verit) \ emeasure\text{-}finite \ indicator\text{-}mult\text{-}ennreal \ mult.\textit{commute}$
 $\ nn\text{-}integral\text{-}cong \ top.\textit{not}\text{-}eq\text{-}extremum)$
ultimately show $?thesis$ **using** $measure\text{-}def$ **by** *metis*
qed

lemma $set\text{-}integrable\text{-}const[simp]:$

$A \in sets \ M \implies emeasure \ M \ A < \infty \implies set\text{-}integrable \ M \ A \ (\lambda\text{-}. \ c)$
using $has\text{-}bochner\text{-}integral\text{-}indicator$ **unfolding** $set\text{-}integrable\text{-}def$ **by** *simp*

lemma $set\text{-}integral\text{-}const[simp]:$

$A \in sets \ M \implies emeasure \ M \ A < \infty \implies set\text{-}lebesgue\text{-}integral \ M \ A \ (\lambda\text{-}. \ c) =$
 $measure \ M \ A \ *_{\mathbb{R}} \ c$
unfolding $set\text{-}lebesgue\text{-}integral\text{-}def$ **using** $has\text{-}bochner\text{-}integral\text{-}indicator$ **by** *force*

lemma $set\text{-}integral\text{-}empty\text{-}0[simp]: set\text{-}lebesgue\text{-}integral \ M \ \{\} \ f = 0$

unfolding $set\text{-}lebesgue\text{-}integral\text{-}def$ **by** *simp*

lemma $set\text{-}integral\text{-}nonneg[simp]:$

fixes $f :: 'a \Rightarrow real$ **and** $A :: 'a \ set$
shows $(\bigwedge x. x \in A \implies 0 \leq f \ x) \implies 0 \leq set\text{-}lebesgue\text{-}integral \ M \ A \ f$
unfolding $set\text{-}lebesgue\text{-}integral\text{-}def$ **by** $(simp \ add: \ indicator\text{-}times\text{-}eq\text{-}if(1))$

lemma

fixes $f :: 'a \Rightarrow 'b::\{banach, \ second\text{-}countable\text{-}topology\}$ **and** $w :: 'a \Rightarrow real$
assumes $A \in sets \ M \ set\text{-}borel\text{-}measurable \ M \ A \ f$

$\bigwedge i.$ *set-borel-measurable* $M A (s i)$ *set-integrable* $M A w$
assumes *lim*: $AE x \in A$ in $M.$ $(\lambda i. s i x) \longrightarrow f x$
assumes *bound*: $\bigwedge i::nat.$ $AE x \in A$ in $M.$ $norm (s i x) \leq w x$
shows *set-integrable-dominated-convergence*: *set-integrable* $M A f$
and *set-integrable-dominated-convergence2*: $\bigwedge i.$ *set-integrable* $M A (s i)$
and *set-integral-dominated-convergence*:
 $(\lambda i. set-lebesgue-integral M A (s i)) \longrightarrow set-lebesgue-integral M A f$
proof –
have $(\lambda x. indicator A x *_R f x) \in borel-measurable M$ **and**
 $\bigwedge i. (\lambda x. indicator A x *_R s i x) \in borel-measurable M$ **and**
integrable $M (\lambda x. indicator A x *_R w x)$
using *assms unfolding set-borel-measurable-def set-integrable-def* **by** *simp-all*
moreover **have** $AE x$ in $M.$ $(\lambda i. indicator A x *_R s i x) \longrightarrow indicator A x$
 $*_R f x$
proof –
obtain N **where** *N-null*: $N \in null-sets M$ **and**
 $si-f: \bigwedge x. x \in space M - N \implies x \in A \longrightarrow (\lambda i. s i x) \longrightarrow f x$
using *lim AE-E3* **by** *(smt (verit))*
hence $\bigwedge x. x \in space M - N \implies (\lambda i. indicator A x *_R s i x) \longrightarrow indicator$
 $A x *_R f x$
proof –
fix x **assume** *asm*: $x \in space M - N$
thus $(\lambda i. indicator A x *_R s i x) \longrightarrow indicator A x *_R f x$
proof *(cases ⟨x ∈ A⟩)*
case *True*
with *si-f asm* **show** *?thesis* **by** *simp*
next
case *False*
thus *?thesis* **by** *simp*
qed
qed
thus *?thesis* **by** *(smt (verit) AE-I' DiffI N-null mem-Collect-eq subsetI)*
qed
moreover **have** $\bigwedge i. AE x$ in $M.$ $norm (indicator A x *_R s i x) \leq indicator A x$
 $*_R w x$
proof –
fix i
from *bound* **obtain** N **where** *N-null*: $N \in null-sets M$ **and**
 $\bigwedge x. x \in space M - N \implies x \in A \longrightarrow norm (s i x) \leq w x$
using *AE-E3* **by** *(smt (verit))*
hence $\bigwedge x. x \in space M - N \implies norm (indicator A x *_R s i x) \leq indicator$
 $A x *_R w x$
by *(simp add: indicator-scaleR-eq-if)*
with *N-null* **show** $AE x$ in $M.$ $norm (indicator A x *_R s i x) \leq indicator A x$
 $*_R w x$
by *(smt (verit) DiffI eventually-ae-filter mem-Collect-eq subsetI)*
qed
ultimately **show** *set-integrable* $M A f \bigwedge i.$ *set-integrable* $M A (s i)$
 $(\lambda i. set-lebesgue-integral M A (s i)) \longrightarrow set-lebesgue-integral M A f$

unfolding *set-integrable-def set-lebesgue-integral-def*
by (*rule integrable-dominated-convergence, rule integrable-dominated-convergence2,*
rule integral-dominated-convergence)
qed

lemma *absolutely-integrable-on-iff-set-integrable:*
fixes $f :: 'a::\text{euclidean-space} \Rightarrow \text{real}$
assumes $f \in \text{borel-measurable lborel}$
and $S \in \text{sets lborel}$
shows $\text{set-integrable lborel } S f \longleftrightarrow f \text{ absolutely-integrable-on } S$
unfolding *set-integrable-def* **apply** (*simp, rewrite integrable-completion[THEN*
sym])
apply *measurable using assms by simp-all*

corollary *integrable-on-iff-set-integrable-nonneg:*
fixes $f :: 'a::\text{euclidean-space} \Rightarrow \text{real}$
assumes $\bigwedge x. x \in S \implies f x \geq 0$ $f \in \text{borel-measurable lborel}$
and $S \in \text{sets lborel}$
shows $\text{set-integrable lborel } S f \longleftrightarrow f \text{ integrable-on } S$
using *absolutely-integrable-on-iff-set-integrable assms*
by (*metis absolutely-integrable-on-iff-nonneg*)

lemma *integrable-on-iff-set-integrable-nonneg-AE:*
fixes $f :: 'a::\text{euclidean-space} \Rightarrow \text{real}$
assumes *AE* $x \in S$ *in lborel.* $f x \geq 0$ $f \in \text{borel-measurable lborel}$
and $S \in \text{sets lborel}$
shows $\text{set-integrable lborel } S f \longleftrightarrow f \text{ integrable-on } S$
proof –
from *assms* **obtain** N **where** $\bigwedge x. x \in S - N \implies f x \geq 0$ **and** *null:*
 $N \in \text{null-sets lborel}$
by (*smt (verit, ccfv-threshold) AE-E3 Diff-iff UNIV-I space-borel space-lborel*)
let $?g = \lambda x. \text{if } x \in N \text{ then } 0 \text{ else } f x$
have [*simp*]: *negligible* N **using** *null negligible-iff-null-sets null-sets-completionI*
by *blast*
have $N \in \text{sets lborel}$ **using** *null by auto*
hence [*simp*]: $?g \in \text{borel-measurable borel}$ **using** *assms by force*
have $\text{set-integrable lborel } S f \longleftrightarrow \text{set-integrable lborel } S ?g$
proof –
have *AE* $x \in S$ *in lborel.* $f x = ?g x$ **by** (*rule AE-I'[of N], simp-all add: null,*
blast)
thus *?thesis* **using** *assms by (intro set-integrable-cong-AE[of f - ?g S]; simp)*
qed
also **have** $\dots \longleftrightarrow ?g \text{ integrable-on } S$
using *assms by (intro integrable-on-iff-set-integrable-nonneg; simp add: nonneg)*
also **have** $\dots \longleftrightarrow f \text{ integrable-on } S$ **by** (*rule integrable-spike-cong[of N]; simp*)
finally **show** *?thesis* .
qed

lemma *set-borel-integral-eq-integral-nonneg:*

```

fixes  $f :: 'a::euclidean-space \Rightarrow real$ 
assumes  $\bigwedge x. x \in S \implies f x \geq 0$   $f \in \text{borel-measurable}$   $\text{borel } S \in \text{sets borel}$ 
shows  $(LINT x : S \mid \text{borel. } f x) = \text{integral } S f$ 
  — Note that  $0 = 0$  holds when the integral diverges.
proof (cases ‹set-integrable lborel S ‹›)
  case True
  thus ?thesis using set-borel-integral-eq-integral by force
next
  case False
  hence  $(LINT x : S \mid \text{borel. } f x) = 0$ 
  unfolding set-lebesgue-integral-def set-integrable-def
  by (rewrite not-integrable-integral-eq; simp)
  moreover have  $\text{integral } S f = 0$ 
  apply (rule not-integrable-integral)
  using False assms by (rewrite in asm integrable-on-iff-set-integrable-nonneg;
simp)
  ultimately show ?thesis ..
qed

```

```

lemma set-borel-integral-eq-integral-nonneg-AE:
  fixes  $f :: 'a::euclidean-space \Rightarrow real$ 
  assumes AE  $x \in S$  in lborel.  $f x \geq 0$   $f \in \text{borel-measurable}$   $\text{borel } S \in \text{sets borel}$ 
  shows  $(LINT x : S \mid \text{borel. } f x) = \text{integral } S f$ 
  — Note that  $0 = 0$  holds when the integral diverges.
proof (cases ‹set-integrable lborel S ‹›)
  case True
  thus ?thesis using set-borel-integral-eq-integral by force
next
  case False
  hence  $(LINT x : S \mid \text{borel. } f x) = 0$ 
  unfolding set-lebesgue-integral-def set-integrable-def
  by (rewrite not-integrable-integral-eq; simp)
  moreover have  $\text{integral } S f = 0$ 
  apply (rule not-integrable-integral)
  using False assms by (rewrite in asm integrable-on-iff-set-integrable-nonneg-AE;
simp)
  ultimately show ?thesis ..
qed

```

2.1 Set Lebesgue Integrability on Affine Transformation

```

lemma set-integrable-Icc-affine-pos-iff:
  fixes  $f :: real \Rightarrow 'a::\{\text{banach, second-countable-topology}\}$  and  $a b c t :: real$ 
  assumes  $c > 0$ 
  shows set-integrable lborel  $\{(a-t)/c..(b-t)/c\}$   $(\lambda x. f (t + c*x))$ 
   $\longleftrightarrow$  set-integrable lborel  $\{a..b\}$   $f$ 
  unfolding set-integrable-def using assms
  apply (rewrite indicator-Icc-affine-pos-inverse, simp)
  by (rule lborel-integrable-real-affine-iff) simp

```

corollary *set-integrable-Icc-shift*:

fixes $f :: \text{real} \Rightarrow 'a::\{\text{banach, second-countable-topology}\}$ **and** $a\ b\ t :: \text{real}$
shows *set-integrable lborel* $\{a-t..b-t\} (\lambda x. f (t+x)) \longleftrightarrow \text{set-integrable lborel}$
 $\{a..b\} f$
using *set-integrable-Icc-affine-pos-iff*[**where** $c=1$] **by** *simp*

lemma *set-integrable-Ici-affine-pos-iff*:

fixes $f :: \text{real} \Rightarrow 'a::\{\text{banach, second-countable-topology}\}$ **and** $a\ c\ t :: \text{real}$
assumes $c > 0$
shows *set-integrable lborel* $\{(a-t)/c..\}$ $(\lambda x. f (t + c*x))$
 $\longleftrightarrow \text{set-integrable lborel } \{a..\} f$
unfolding *set-integrable-def* **using** *assms*
apply (*rewrite indicator-Ici-affine-pos-inverse, simp*)
by (*rule lborel-integrable-real-affine-iff*) *simp*

corollary *set-integrable-Ici-shift*:

fixes $f :: \text{real} \Rightarrow 'a::\{\text{banach, second-countable-topology}\}$ **and** $a\ t :: \text{real}$
shows *set-integrable lborel* $\{a-t..\}$ $(\lambda x. f (t+x)) \longleftrightarrow \text{set-integrable lborel } \{a..\} f$
using *set-integrable-Ici-affine-pos-iff*[**where** $c=1$] **by** *simp*

lemma *set-integrable-Iic-affine-pos-iff*:

fixes $f :: \text{real} \Rightarrow 'a::\{\text{banach, second-countable-topology}\}$ **and** $b\ c\ t :: \text{real}$
assumes $c > 0$
shows *set-integrable lborel* $\{..(b-t)/c\}$ $(\lambda x. f (t + c*x))$
 $\longleftrightarrow \text{set-integrable lborel } \{..b\} f$
unfolding *set-integrable-def* **using** *assms*
apply (*rewrite indicator-Iic-affine-pos-inverse, simp*)
by (*rule lborel-integrable-real-affine-iff*) *simp*

corollary *set-integrable-Iic-shift*:

fixes $f :: \text{real} \Rightarrow 'a::\{\text{banach, second-countable-topology}\}$ **and** $b\ t :: \text{real}$
shows *set-integrable lborel* $\{..b-t\} (\lambda x. f (t+x)) \longleftrightarrow \text{set-integrable lborel } \{..b\} f$
using *set-integrable-Iic-affine-pos-iff*[**where** $c=1$] **by** *simp*

lemma *set-integrable-Icc-affine-neg-iff*:

fixes $f :: \text{real} \Rightarrow 'a::\{\text{banach, second-countable-topology}\}$ **and** $a\ b\ c\ t :: \text{real}$
assumes $c < 0$
shows *set-integrable lborel* $\{(b-t)/c..(a-t)/c\}$ $(\lambda x. f (t + c*x))$
 $\longleftrightarrow \text{set-integrable lborel } \{a..b\} f$
unfolding *set-integrable-def* **using** *assms*
apply (*rewrite indicator-Icc-affine-neg-inverse, simp*)
by (*rule lborel-integrable-real-affine-iff*) *simp*

corollary *set-integrable-Icc-reverse*:

fixes $f :: \text{real} \Rightarrow 'a::\{\text{banach, second-countable-topology}\}$ **and** $a\ b\ t :: \text{real}$
shows *set-integrable lborel* $\{t-b..t-a\} (\lambda x. f (t-x)) \longleftrightarrow \text{set-integrable lborel}$
 $\{a..b\} f$
using *set-integrable-Icc-affine-neg-iff*[**where** $c=-1$] **by** *simp*

lemma *set-integrable-Ici-affine-neg-iff*:

fixes $f :: \text{real} \Rightarrow 'a :: \{\text{banach, second-countable-topology}\}$ **and** $b \ c \ t :: \text{real}$
assumes $c < 0$
shows $\text{set-integrable lborel } \{(b-t)/c..\}$ $(\lambda x. f (t + c*x))$
 $\longleftrightarrow \text{set-integrable lborel } \{..b\}$ f
unfolding *set-integrable-def* **using** *assms*
apply (*rewrite indicator-Ici-affine-neg-inverse, simp*)
by (*rule lborel-integrable-real-affine-iff*) *simp*

corollary *set-integrable-Ici-reverse*:

fixes $f :: \text{real} \Rightarrow 'a :: \{\text{banach, second-countable-topology}\}$ **and** $b \ t :: \text{real}$
shows $\text{set-integrable lborel } \{t-b..\}$ $(\lambda x. f (t-x)) \longleftrightarrow \text{set-integrable lborel } \{..b\}$ f
using *set-integrable-Ici-affine-neg-iff* [**where** $c=-1$] **by** *simp*

lemma *set-integrable-Iic-affine-neg-iff*:

fixes $f :: \text{real} \Rightarrow 'a :: \{\text{banach, second-countable-topology}\}$ **and** $a \ c \ t :: \text{real}$
assumes $c < 0$
shows $\text{set-integrable lborel } \{..(a-t)/c\}$ $(\lambda x. f (t + c*x))$
 $\longleftrightarrow \text{set-integrable lborel } \{a..\}$ f
unfolding *set-integrable-def* **using** *assms*
apply (*rewrite indicator-Iic-affine-neg-inverse, simp*)
by (*rule lborel-integrable-real-affine-iff*) *simp*

corollary *set-integrable-Iic-reverse*:

fixes $f :: \text{real} \Rightarrow 'a :: \{\text{banach, second-countable-topology}\}$ **and** $a \ t :: \text{real}$
shows $\text{set-integrable lborel } \{..t-a\}$ $(\lambda x. f (t-x)) \longleftrightarrow \text{set-integrable lborel } \{a..\}$ f
using *set-integrable-Iic-affine-neg-iff* [**where** $c=-1$] **by** *simp*

2.2 Set Lebesgue Integral on Affine Transformation

lemma *lborel-set-integral-Icc-affine-pos*:

fixes $f :: \text{real} \Rightarrow 'a :: \{\text{banach, second-countable-topology}\}$ **and** $a \ b \ c :: \text{real}$
assumes $c > 0$
shows $(\int x \in \{a..b\}. f \ x \ \partial \text{lborel}) = c *_R (\int x \in \{(a-t)/c..(b-t)/c\}. f (t + c*x) \ \partial \text{lborel})$
unfolding *set-lebesgue-integral-def* **using** *assms*
apply (*rewrite indicator-Icc-affine-pos-inverse, simp*)
using *lborel-integral-real-affine* [**where** $c=c$] **by** *force*

corollary *lborel-set-integral-Icc-shift*:

fixes $f :: \text{real} \Rightarrow 'a :: \{\text{banach, second-countable-topology}\}$ **and** $a \ b :: \text{real}$
shows $(\int x \in \{a..b\}. f \ x \ \partial \text{lborel}) = (\int x \in \{a-t..b-t\}. f (t+x) \ \partial \text{lborel})$
using *lborel-set-integral-Icc-affine-pos* [**where** $c=1$] **by** *simp*

lemma *lborel-set-integral-Ici-affine-pos*:

fixes $f :: \text{real} \Rightarrow 'a :: \{\text{banach, second-countable-topology}\}$ **and** $a \ c :: \text{real}$
assumes $c > 0$
shows $(\int x \in \{a..\}. f \ x \ \partial \text{lborel}) = c *_R (\int x \in \{(a-t)/c..\}. f (t + c*x) \ \partial \text{lborel})$

unfolding *set-lebesgue-integral-def* **using** *assms*
apply (*rewrite indicator-Ici-affine-pos-inverse, simp*)
using *lborel-integral-real-affine*[**where** $c=c$] **by force**

corollary *lborel-set-integral-Ici-shift*:

fixes $f :: \text{real} \Rightarrow 'a :: \{\text{banach, second-countable-topology}\}$ **and** $a :: \text{real}$
shows $(\int x \in \{a..b\}. f x \partial \text{lborel}) = (\int x \in \{a-t..b-t\}. f (t+x) \partial \text{lborel})$
using *lborel-set-integral-Ici-affine-pos*[**where** $c=1$] **by simp**

lemma *lborel-set-integral-Ici-affine-pos*:

fixes $f :: \text{real} \Rightarrow 'a :: \{\text{banach, second-countable-topology}\}$ **and** $b c :: \text{real}$
assumes $c > 0$
shows $(\int x \in \{..b\}. f x \partial \text{lborel}) = c *_R (\int x \in \{..(b-t)/c\}. f (t + c*x) \partial \text{lborel})$
unfolding *set-lebesgue-integral-def* **using** *assms*
apply (*rewrite indicator-Ici-affine-pos-inverse, simp*)
using *lborel-integral-real-affine*[**where** $c=c$] **by force**

corollary *lborel-set-integral-Ici-shift*:

fixes $f :: \text{real} \Rightarrow 'a :: \{\text{banach, second-countable-topology}\}$ **and** $b :: \text{real}$
shows $(\int x \in \{..b\}. f x \partial \text{lborel}) = (\int x \in \{..b-t\}. f (t+x) \partial \text{lborel})$
using *lborel-set-integral-Ici-affine-pos*[**where** $c=1$] **by simp**

lemma *lborel-set-integral-Icc-affine-neg*:

fixes $f :: \text{real} \Rightarrow 'a :: \{\text{banach, second-countable-topology}\}$ **and** $a b c :: \text{real}$
assumes $c < 0$
shows $(\int x \in \{a..b\}. f x \partial \text{lborel}) = -c *_R (\int x \in \{(b-t)/c..(a-t)/c\}. f (t + c*x) \partial \text{lborel})$
unfolding *set-lebesgue-integral-def* **using** *assms*
apply (*rewrite indicator-Icc-affine-neg-inverse, simp*)
using *lborel-integral-real-affine*[**where** $c=c$] **by force**

corollary *lborel-set-integral-Icc-reverse*:

fixes $f :: \text{real} \Rightarrow 'a :: \{\text{banach, second-countable-topology}\}$ **and** $a b :: \text{real}$
shows $(\int x \in \{a..b\}. f x \partial \text{lborel}) = (\int x \in \{t-b..t-a\}. f (t-x) \partial \text{lborel})$
using *lborel-set-integral-Icc-affine-neg*[**where** $c=-1$] **by simp**

lemma *lborel-set-integral-Ici-affine-neg*:

fixes $f :: \text{real} \Rightarrow 'a :: \{\text{banach, second-countable-topology}\}$ **and** $b c :: \text{real}$
assumes $c < 0$
shows $(\int x \in \{..b\}. f x \partial \text{lborel}) = -c *_R (\int x \in \{(b-t)/c..t/c\}. f (t + c*x) \partial \text{lborel})$
unfolding *set-lebesgue-integral-def* **using** *assms*
apply (*rewrite indicator-Ici-affine-neg-inverse, simp*)
using *lborel-integral-real-affine*[**where** $c=c$] **by force**

corollary *lborel-set-integral-Ici-reverse*:

fixes $f :: \text{real} \Rightarrow 'a :: \{\text{banach, second-countable-topology}\}$ **and** $b :: \text{real}$
shows $(\int x \in \{..b\}. f x \partial \text{lborel}) = (\int x \in \{t-b..t\}. f (t-x) \partial \text{lborel})$
using *lborel-set-integral-Ici-affine-neg*[**where** $c=-1$] **by simp**

lemma *lborel-set-integral-Iic-affine-neg*:

fixes $f :: \text{real} \Rightarrow 'a :: \{\text{banach, second-countable-topology}\}$ **and** $a\ c :: \text{real}$
assumes $c < 0$
shows $(\int x \in \{a..\}. f\ x\ \partial \text{lborel}) = -c *_R (\int x \in \{..(a-t)/c\}. f\ (t + c*x)\ \partial \text{lborel})$
unfolding *set-lebesgue-integral-def* **using** *assms*
apply (*rewrite indicator-Iic-affine-neg-inverse, simp*)
using *lborel-integral-real-affine* [**where** $c=c$] **by force**

corollary *lborel-set-integral-Iic-reverse*:

fixes $f :: \text{real} \Rightarrow 'a :: \{\text{banach, second-countable-topology}\}$ **and** $a :: \text{real}$
shows $(\int x \in \{a..\}. f\ x\ \partial \text{lborel}) = (\int x \in \{..t-a\}. f\ (t-x)\ \partial \text{lborel})$
using *lborel-set-integral-Iic-affine-neg* [**where** $c=-1$] **by simp**

lemma *set-integrable-Ici-equiv-aux*:

fixes $f :: \text{real} \Rightarrow 'a :: \{\text{banach, second-countable-topology}\}$ **and** $a\ b :: \text{real}$
assumes $\bigwedge c\ d. \text{set-integrable lborel } \{c..d\} f\ a \leq b$
shows $\text{set-integrable lborel } \{a..\} f \iff \text{set-integrable lborel } \{b..\} f$

proof

assume $\text{set-integrable lborel } \{a..\} f$
thus $\text{set-integrable lborel } \{b..\} f$ **by** (*rule set-integrable-subset; simp add: assms*)

next

assume $\text{set-integrable lborel } \{b..\} f$
moreover have $\text{set-integrable lborel } \{a..b\} f$ **using** *assms* **by blast**
moreover have $\{a..\} = \{a..b\} \cup \{b..\}$ **using** *assms* **by auto**
ultimately show $\text{set-integrable lborel } \{a..\} f$ **using** *set-integrable-Un* **by force**

qed

corollary *set-integrable-Ici-equiv*:

fixes $f :: \text{real} \Rightarrow 'a :: \{\text{banach, second-countable-topology}\}$ **and** $a\ b :: \text{real}$
assumes $\bigwedge c\ d. \text{set-integrable lborel } \{c..d\} f$
shows $\text{set-integrable lborel } \{a..\} f \iff \text{set-integrable lborel } \{b..\} f$
using *set-integrable-Ici-equiv-aux* *assms* **by** (*smt (verit)*)

lemma *set-integrable-Iic-equiv*:

fixes $f :: \text{real} \Rightarrow \text{real}$ **and** $a\ b :: \text{real}$
assumes $\bigwedge c\ d. \text{set-integrable lborel } \{c..d\} f$
shows $\text{set-integrable lborel } \{..a\} f \iff \text{set-integrable lborel } \{..b\} f$ (**is** *?LHS* \iff *?RHS*)

proof –

have *?LHS* $\iff \text{set-integrable lborel } \{-a..\} (\lambda x. f\ (-x))$
using *set-integrable-Ici-reverse* [**where** $t=0$] **by force**
also have $\dots \iff \text{set-integrable lborel } \{-b..\} (\lambda x. f\ (-x))$

proof –

have $\bigwedge c\ d. \text{set-integrable lborel } \{c..d\} (\lambda x. f\ (-x))$
apply (*rewrite at* $\{\square..\}$ *minus-minus* [*THEN sym*])
apply (*rewrite at* $\{..\square\}$ *minus-minus* [*THEN sym*])
using *assms set-integrable-Icc-reverse* [**where** $t=0$] **by force**
thus *?thesis* **by** (*rule set-integrable-Ici-equiv*)

qed

also have ... \longleftrightarrow ?RHS using *set-integrable-Ici-reverse*[where $t=0$] by force
 finally show ?thesis .
 qed

2.3 Alternative Integral Test

lemma *nn-integral-suminf-Ico-real-nat*:

fixes $a::real$ and $f :: real \Rightarrow ennreal$

assumes $f \in borel\text{-}measurable\ lborel$

shows $(\int^{+x \in \{a..\}} f\ x\ \partial lborel) = (\sum k. \int^{+x \in \{a+k..<a+k+1\}} f\ x\ \partial lborel)$

apply (*rewrite Ico-real-nat-union*[THEN *sym*])

using *Ico-real-nat-disjoint assms* by (*intro nn-integral-disjoint-family; simp*)

lemma *set-integrable-iff-bounded*:

fixes $f :: 'a \Rightarrow 'b::\{banach, second-countable-topology\}$

assumes $A \in sets\ M$

shows *set-integrable* $M\ A\ f \longleftrightarrow$ *set-borel-measurable* $M\ A\ f \wedge (\int^{+x \in A. norm\ (f\ x)\ \partial M}) < \infty$

unfolding *set-integrable-def set-borel-measurable-def* using *integrable-iff-bounded*

by (*smt (verit, ccfv-threshold) indicator-mult-ennreal indicator-pos-le*

mult.commute nn-integral-cong norm-scaleR)

theorem *set-integrable-iff-summable*:

fixes $a::real$ and $f :: real \Rightarrow real$

assumes *antimono-on* $\{a..\} f \wedge x. a \leq x \implies f\ x \geq 0$ $f \in borel\text{-}measurable\ lborel$

shows *set-integrable* $lborel\ \{a..\} f \longleftrightarrow$ *summable* $(\lambda k. f\ (a+k))$

proof

assume *asm*: *set-integrable* $lborel\ \{a..\} f$

have [*measurable*]: $(\lambda x. ennreal\ (f\ x)) \in borel\text{-}measurable\ lborel$ using *assms* by *simp*

have $\forall k \geq 0. norm\ (f\ (a+(k+1)::nat)) \leq (\int x \in \{a+k..<a+k+1\}. f\ x\ \partial lborel)$

proof -

{ **fix** $k::nat$

have $norm\ (f\ (a+(k+1)::nat)) = f\ (a+k+1)$

using *assms* by (*smt (verit) of-nat-0-le-iff of-nat-1 of-nat-add real-norm-def*)

also have ... = $(\int x \in \{a+k..<a+k+1\}. f\ (a+k+1)\ \partial lborel)$

unfolding *set-lebesgue-integral-def* by *simp*

also have ... $\leq (\int x \in \{a+k..<a+k+1\}. f\ x\ \partial lborel)$

apply (*rule set-integral-mono, simp*)

apply (*rule set-integrable-restrict-space*[of $lborel\ \{a..\}$], *simp add: asm*)

apply (*rewrite sets-restrict-space, force*)

using *assms* **unfolding** *mono-on-def monotone-on-def* by *simp*

finally have $norm\ (f\ (a+(k+1)::nat)) \leq (\int x \in \{a+k..<a+k+1\}. f\ x\ \partial lborel)$

. }

thus ?thesis by *simp*

qed

moreover have *summable* $(\lambda k. \int x \in \{a+k..<a+k+1\}. f\ x\ \partial lborel)$

proof -

```

have (∫+ x∈{a..}. ennreal (f x) ∂lborel) ≠ ∞
  using asm unfolding set-integrable-def apply simp
  by (smt (verit) indicator-mult-ennreal infinity-ennreal-def mult.commute
      nn-integral-cong real-integrable-def)
thus ?thesis
  apply (rewrite in asm nn-integral-suminf-Ico-real-nat, simp)
  apply (rule summable-suminf-not-top)
  using assms apply (intro set-integral-nonneg, force)
  apply (rewrite set-nn-integral-eq-set-integral[THEN sym], simp add: assms)
  by (rule set-integrable-subset[of lborel {a..}], simp-all add: asm) force
qed
ultimately have summable (λk. f (a+(k+1::nat)))
  using summable-comparison-test by (smt (verit, del-Insts))
thus summable (λk. f (a+k)) using summable-iff-shift by blast
next
assume asm: summable (λk. f (a+k))
hence (∫+ x∈{a..}. ennreal |f x| ∂lborel) < ∞
proof -
  have (∫+ x∈{a..}. ennreal |f x| ∂lborel) = (∫+ x∈{a..}. ennreal (f x) ∂lborel)
  using assms by (metis abs-of-nonneg atLeast-iff indicator-simps(2) mult-eq-0-iff)
  also have ... = (∑ k. ∫+ x∈{a+k..using assms by (rewrite nn-integral-suminf-Ico-real-nat; simp)
  also have ... ≤ (∑ k. ∫+ x∈{a+k..proof -
  have ∧(k::nat) x. x∈{a+k..using assms unfolding monotone-on-def by auto
  thus ?thesis
  apply (intro suminf-le, simp-all)
  by (rule nn-integral-mono)
  (metis (no-types, opaque-lifting) atLeastLessThan-iff dual-order.refl en-
nreal-leI
indicator-simps(2) mult-eq-0-iff mult-mono zero-le)
qed
also have ... = (∑ k. ennreal (f (a+k)))
  apply (rule suminf-cong)
  by (rewrite nn-integral-cmult-indicator; simp)
also have ... < ∞
  unfolding infinity-ennreal-def apply (rewrite less-top[THEN sym])
  using asm assms by (smt (verit) of-nat-0-le-iff suminf-cong suminf-ennreal2
top-neq-ennreal)
  finally show ?thesis .
qed
moreover have set-borel-measurable lborel {a..} f
  using assms unfolding set-borel-measurable-def by simp
ultimately show set-integrable lborel {a..} f by (rewrite set-integrable-iff-bounded)
auto
qed

```

2.4 Interchange of Differentiation and Lebesgue Integration

definition *measurable-extension* :: 'a measure \Rightarrow 'b measure \Rightarrow ('a \Rightarrow 'b) \Rightarrow 'a \Rightarrow 'b **where**

measurable-extension $M\ N\ f =$
 (SOME $g. g \in M \rightarrow_M N \wedge (\exists S \in (\text{null-sets } M). \{x \in \text{space } M. f\ x \neq g\ x\} \subseteq S)$)

- The term *measurable-extension* is proposed by Reynald Affeldt.
- This function is used to make an almost-everywhere-defined function measurable.

lemma

fixes $f\ g$

assumes $g \in M \rightarrow_M N\ S \in \text{null-sets } M\ \{x \in \text{space } M. f\ x \neq g\ x\} \subseteq S$

shows *measurable-extensionI*: $AE\ x\ \text{in } M. f\ x = \text{measurable-extension } M\ N\ f\ x$

and

measurable-extensionI2: $AE\ x\ \text{in } M. g\ x = \text{measurable-extension } M\ N\ f\ x$ **and**

measurable-extension-measurable: $\text{measurable-extension } M\ N\ f \in \text{measurable } M\ N$

N

proof –

let $?G = \lambda g. g \in M \rightarrow_M N$ **and** $?S = \lambda g. \exists S \in \text{null-sets } M. \{x \in \text{space } M. f\ x \neq g\ x\} \subseteq S$

show $AE\ x\ \text{in } M. f\ x = \text{measurable-extension } M\ N\ f\ x$

unfolding *measurable-extension-def*

apply (rule *someI2*[of $\lambda g. ?G\ g \wedge ?S\ g\ g$])

using *assms* **apply** *blast*

using *AE-I'* **by** *auto*

moreover have $AE\ x\ \text{in } M. g\ x = f\ x$

using *assms* **by** (*smt* (*verit*, *best*) *AE-I'* *Collect-cong*)

ultimately show $AE\ x\ \text{in } M. g\ x = \text{measurable-extension } M\ N\ f\ x$ **by force**

show $\text{measurable-extension } M\ N\ f \in \text{measurable } M\ N$

unfolding *measurable-extension-def*

apply (rule *conjE*[of $?G\ g\ ?S\ g$])

using *assms* **apply** *auto*[1]

using *someI-ex*[of $\lambda g. ?G\ g \wedge ?S\ g$] **by** *auto*

qed

corollary *measurable-measurable-extension-AE*:

fixes f

assumes $f \in M \rightarrow_M N$

shows $AE\ x\ \text{in } M. f\ x = \text{measurable-extension } M\ N\ f\ x$

by (rule *measurable-extensionI*[**where** $g=f$ **and** $S=\{\}$]; *simp add*: *assms*)

definition *borel-measurable-extension* ::

'a measure \Rightarrow ('a \Rightarrow 'b::*topological-space*) \Rightarrow 'a \Rightarrow 'b **where**

borel-measurable-extension $M\ f = \text{measurable-extension } M\ \text{borel } f$

lemma

fixes $f\ g$

assumes $g \in \text{borel-measurable } M\ S \in \text{null-sets } M\ \{x \in \text{space } M. f\ x \neq g\ x\} \subseteq S$

shows *borel-measurable-extensionI*: $AE\ x\ in\ M. f\ x = borel-measurable-extension\ M\ f\ x$ **and**
borel-measurable-extensionI2: $AE\ x\ in\ M. g\ x = borel-measurable-extension\ M\ f\ x$ **and**
borel-measurable-extension-measurable: $borel-measurable-extension\ M\ f \in borel-measurable\ M$
unfolding *borel-measurable-extension-def* **using** *assms*
apply –
using *measurable-extensionI* **apply** *blast*
using *measurable-extensionI2* **apply** *blast*
using *measurable-extension-measurable* **by** *blast*

corollary *borel-measurable-measurable-extension-AE*:
fixes *f*
assumes $f \in borel-measurable\ M$
shows $AE\ x\ in\ M. f\ x = borel-measurable-extension\ M\ f\ x$
using *assms measurable-measurable-extension-AE* **unfolding** *borel-measurable-extension-def*
by *auto*

definition *set-borel-measurable-extension* ::
 $'a\ measure \Rightarrow 'a\ set \Rightarrow ('a \Rightarrow 'b::topological-space) \Rightarrow 'a \Rightarrow 'b$
where *set-borel-measurable-extension* $M\ A\ f = borel-measurable-extension\ (restrict-space\ M\ A)\ f$

lemma
fixes $f\ g :: 'a \Rightarrow 'b::real-normed-vector$ **and** A
assumes $A \in sets\ M\ set-borel-measurable\ M\ A\ g\ S \in null-sets\ M\ \{x \in A. f\ x \neq g\ x\} \subseteq S$
shows *set-borel-measurable-extensionI*:
 $AE\ x \in A\ in\ M. f\ x = set-borel-measurable-extension\ M\ A\ f\ x$ **and**
 $set-borel-measurable-extensionI2$:
 $AE\ x \in A\ in\ M. g\ x = set-borel-measurable-extension\ M\ A\ f\ x$ **and**
 $set-borel-measurable-extension-measurable$:
 $set-borel-measurable\ M\ A\ (set-borel-measurable-extension\ M\ A\ f)$
proof –
have $g \in borel-measurable\ (restrict-space\ M\ A)$
using *assms* **by** (*rewrite set-borel-measurable-restrict-space-iff; simp*)
moreover **have** $S \cap A \in null-sets\ (restrict-space\ M\ A)$
using *assms null-sets-restrict-space* **by** (*metis Int-lower2 null-set-Int2*)
moreover **have** $\{x \in space\ (restrict-space\ M\ A). f\ x \neq g\ x\} \subseteq S \cap A$
using *assms* **by** (*rewrite space-restrict-space2; simp*)
ultimately show $AE\ x \in A\ in\ M. f\ x = set-borel-measurable-extension\ M\ A\ f\ x$
and
 $AE\ x \in A\ in\ M. g\ x = set-borel-measurable-extension\ M\ A\ f\ x$ **and**
 $set-borel-measurable\ M\ A\ (set-borel-measurable-extension\ M\ A\ f)$
unfolding *set-borel-measurable-extension-def* **using** *assms*
apply –
apply (*rewrite AE-restrict-space-iff[THEN sym], simp*)
apply (*rule borel-measurable-extensionI[of g - S ∩ A]; simp*)

apply (*rewrite AE-restrict-space-iff*[*THEN sym*], *simp*)
apply (*rule borel-measurable-extensionI2*[*of g - S ∩ A*]; *simp*)
apply (*rewrite set-borel-measurable-restrict-space-iff*[*THEN sym*], *simp*)
by (*rule borel-measurable-extension-measurable*[*of g - S ∩ A*]; *simp*)
qed

corollary *set-borel-measurable-measurable-extension-AE*:

fixes $f :: 'a \Rightarrow 'b :: \text{real-normed-vector}$ **and** A
assumes *set-borel-measurable* $M A f A \in \text{sets } M$
shows $AE\ x \in A$ in M . $f\ x = \text{set-borel-measurable-extension } M A f\ x$
using *set-borel-measurable-restrict-space-iff*
borel-measurable-measurable-extension-AE AE-restrict-space-iff
unfolding *set-borel-measurable-extension-def*
by (*smt (verit) AE-cong sets.Int-space-eq2 assms*)

proposition *interchange-deriv-LINT-general*:

fixes $a\ b :: \text{real}$ **and** $f :: \text{real} \Rightarrow 'a \Rightarrow \text{real}$ **and** $g :: 'a \Rightarrow \text{real}$
assumes *f-integ*: $\bigwedge r. r \in \{a <..<b\} \implies \text{integrable } M (f\ r)$ **and**
f-diff: $AE\ x$ in M . $(\lambda r. f\ r\ x)$ *differentiable-on* $\{a <..<b\}$ **and**
Df-bound: $AE\ x$ in M . $\forall r \in \{a <..<b\}. |\text{deriv } (\lambda r. f\ r\ x)\ r| \leq g\ x$ *integrable* $M\ g$
shows $\bigwedge r. r \in \{a <..<b\} \implies ((\lambda r. \int x. f\ r\ x\ \partial M)$ *has-real-derivative*
 $\int x. \text{borel-measurable-extension } M (\lambda x. \text{deriv } (\lambda r. f\ r\ x)\ r)\ x\ \partial M)$ (at r)
proof –

Preparation

have *f-msr*: $\bigwedge r. r \in \{a <..<b\} \implies f\ r \in \text{borel-measurable } M$ **using** *f-integ* **by**
auto
from *f-diff* **obtain** $N1$ **where** *N1-null*: $N1 \in \text{null-sets } M$ **and**
 $\bigwedge x. x \in \text{space } M - N1 \implies (\lambda s. f\ s\ x)$ *differentiable-on* $\{a <..<b\}$
by (*smt (verit) AE-E3*)
hence *f-diffN1*: $\bigwedge x. x \in \text{space } M - N1 \implies (\lambda s. f\ s\ x)$ *differentiable-on* $\{a <..<b\}$
by (*meson Diff-iff sets.sets-into-space subset-eq*)
from *Df-bound* **obtain** $N2$ **where** *N2-null*: $N2 \in \text{null-sets } M$ **and**
 $\bigwedge x. x \in \text{space } M - N2 \implies \forall r \in \{a <..<b\}. |\text{deriv } (\lambda s. f\ s\ x)\ r| \leq g\ x$
by (*smt (verit) AE-E3*)
hence *Df-boundN2*: $\bigwedge x. x \in \text{space } M - N2 \implies \forall r \in \{a <..<b\}. |\text{deriv } (\lambda s. f\ s\ x)$
 $r| \leq g\ x$
by (*meson Diff-iff sets.sets-into-space subset-eq*)
define N **where** $N \equiv N1 \cup N2$
let $?CN = \text{space } M - N$
have *N-null*: $N \in \text{null-sets } M$ **and** *N-msr*: $N \in \text{sets } M$
unfolding *N-def* **using** *N1-null N2-null* **by** *auto*
have *f-diffCN*: $\bigwedge x. x \in ?CN \implies (\lambda s. f\ s\ x)$ *differentiable-on* $\{a <..<b\}$
unfolding *N-def* **using** *f-diffN1* **by** *simp*
define $Df :: \text{real} \Rightarrow 'a \Rightarrow \text{real}$ **where**
 $Df\ r\ x \equiv \text{indicator } (\{a <..<b\} \times ?CN) (r, x) * \text{deriv } (\lambda s. f\ s\ x)\ r$ **for** $r\ x$
have *Df-boundCN*: $\bigwedge x. x \in ?CN \implies \forall r \in \{a <..<b\}. |Df\ r\ x| \leq g\ x$
unfolding *Df-def N-def* **using** *Df-boundN2* **by** *simp*

Main Part of the Proof

fix r **assume** $r\text{-ab}$: $r \in \{a < \cdot < b\}$
then obtain e **where** $e\text{-pos}$: $e > 0$ **and** $ball\text{-ab}$: $ball\ r\ e \subseteq \{a < \cdot < b\}$
by (*meson openE open-greaterThanLessThan*)
have $\bigwedge d::nat \Rightarrow real$. $[\forall i. d\ i \in UNIV - \{0\}; d \longrightarrow 0] \Longrightarrow$
 $((\lambda h. ((\int x. f\ (r+h)\ x\ \partial M) - \int x. f\ r\ x\ \partial M) / h) \circ d) \longrightarrow$
 $\int x. borel\text{-measurable-extension}\ M\ (\lambda y. deriv\ (\lambda s. f\ s\ y)\ r)\ x\ \partial M$
proof –
fix $d::nat \Rightarrow real$ **assume** $d\text{-neq0}$: $\forall i. d\ i \in UNIV - \{0\}$ **and** $d\text{-to0}$: $d \longrightarrow 0$
then obtain m **where** $\forall i \geq m. |d\ i - 0| < e$ **using** *LIMSEQ-def e-pos*
dist-real-def **by** *metis*
hence $rd\text{-ab}$: $\bigwedge n. r + d\ (n+m) \in \{a < \cdot < b\}$ **using** *dist-real-def ball-ab* **by** (*simp*
add: subset-eq)
hence $fd\text{-msr}$: $\bigwedge n. (\lambda x. (f\ (r + d\ (n+m))\ x) - f\ r\ x) / d\ (n+m) \in borel\text{-measurable}$
 M
using $r\text{-ab}$ **by** (*measurable; (intro f-msr)?; simp*)
hence $limf\text{-msr}$: $(\lambda x. lim\ (\lambda n. (f\ (r + d\ (n+m))\ x) - f\ r\ x) / d\ (n+m)) \in$
borel-measurable M
by *measurable*
moreover have $limf\text{-Df}$: $\bigwedge x. x \in ?CN \Longrightarrow (\lambda n. (f\ (r + d\ (n+m))\ x) - f\ r\ x)$
 $/ d\ (n+m) \longrightarrow Df\ r\ x$
proof –
fix x **assume** $x\text{-CN}$: $x \in ?CN$
hence $(\lambda s. f\ s\ x)$ *field-differentiable (at r)*
using $f\text{-diffCN}\ r\text{-ab}$
by (*metis at-within-open differentiable-on-eq-field-differentiable-real*
open-greaterThanLessThan)
hence $((\lambda h. (f\ (r+h)\ x) - f\ r\ x) / h) \longrightarrow Df\ r\ x$ (*at 0*)
apply (*rewrite in asm DERIV-deriv-iff-field-differentiable[THEN sym]*)
unfolding $Df\text{-def}$ **using** $r\text{-ab}\ x\text{-CN}$ **by** (*simp add: DERIV-def*)
hence $(\lambda i. (f\ (r + d\ i)\ x) - f\ r\ x) / d\ i \longrightarrow Df\ r\ x$
apply (*rewrite in asm tendsto-at-iff-sequentially*)
apply (*rule allE'[where x=d], simp*)
unfolding $comp\text{-def}$ **using** $d\text{-neq0}\ d\text{-to0}$ **by** *simp*
thus $(\lambda n. (f\ (r + d\ (n+m))\ x) - f\ r\ x) / d\ (n+m) \longrightarrow Df\ r\ x$
by (*rule LIMSEQ-ignore-initial-segment[where k=m]*)
qed
ultimately have $Df\text{-eq}$:
 $\bigwedge x. Df\ r\ x = indicator\ ?CN\ x * lim\ (\lambda n. (f\ (r + d\ (n+m))\ x) - f\ r\ x) / d$
 $(n+m)$
proof –
fix x
show $Df\ r\ x = indicator\ ?CN\ x * lim\ (\lambda n. (f\ (r + d\ (n+m))\ x) - f\ r\ x) / d$
 $(n+m)$
proof (*cases <x ∈ ?CN>*)
case *True*
hence $lim\ (\lambda n. (f\ (r + d\ (n+m))\ x) - f\ r\ x) / d\ (n+m) = Df\ r\ x$
by (*intro limI, rule limf-Df*)
thus *?thesis* **using** *True* **by** *simp*
next

```

    case False
    thus ?thesis unfolding Df-def by simp
  qed
  qed
  hence Df-msr:  $Df\ r \in \text{borel-measurable } M$ 
  apply (rewrite in  $\lambda x. \sqsupset Df\text{-eq}$ )
  apply (measurable; (rule limf-msr)?)
  using N-null unfolding null-sets-def by force
  have (( $\lambda h. ((\int x. f\ (r+h)\ x\ \partial M) - \int x. f\ r\ x\ \partial M) / h$ )  $\circ d$ )  $\longrightarrow$ 
     $\int x. \lim (\lambda n. (f\ (r + d\ (n+m))\ x - f\ r\ x) / d\ (n+m))\ \partial M$ 
  proof -
    have ( $\lambda n. \int x. (f\ (r + d\ (n+m))\ x - f\ r\ x) / d\ (n+m)\ \partial M$ )  $\longrightarrow$ 
       $\int x. \lim (\lambda n. (f\ (r + d\ (n+m))\ x - f\ r\ x) / d\ (n+m))\ \partial M$ 
    proof - — by Lebesgue's Dominated Convergence Theorem
      have AE x in M. ( $\lambda n. (f\ (r + d\ (n+m))\ x - f\ r\ x) / d\ (n+m)$ )  $\longrightarrow$ 
         $\lim (\lambda n. (f\ (r + d\ (n+m))\ x - f\ r\ x) / d\ (n+m))$ 
      using limf-Df Df-eq N-null by (smt (verit) DiffI AE-I' limI mem-Collect-eq
subset-eq)
      moreover have  $\bigwedge n. AE\ x\ \text{in } M. \text{norm } ((f\ (r + d\ (n+m))\ x - f\ r\ x) / d\ (n+m)) \leq g\ x$ 
      proof -
        fix n
        { fix x assume x-CN:  $x \in ?CN$ 
          let ?I =  $\{r..(r + d\ (n+m))\} \cup \{(r + d\ (n+m))..r\}$ 
          have f-diffI: ( $\lambda s. f\ s\ x$ ) differentiable-on ?I
          apply (rule differentiable-on-subset[where  $t = \{a <.. < b\}$ ], rule f-diffCN,
rule x-CN)
          using r-ab rd-ab[of n] by (rewrite Un-subset-iff, auto)
          hence continuous-on ?I ( $\lambda s. f\ s\ x$ ) ( $\lambda s. f\ s\ x$ ) differentiable-on interior ?I
          apply -
          using differentiable-imp-continuous-on apply blast
          by (metis differentiable-on-subset interior-subset)
          then obtain t where t-01:  $t \in \{0 <.. < 1\}$  and
            f-MVT:  $f\ (r + d\ (n+m))\ x - f\ r\ x = d\ (n+m) * \text{deriv } (\lambda s. f\ s\ x)\ (r + t * (d\ (n+m)))$ 
            by (rule MVT-order-free)
            hence  $0 < t\ t < 1$  by simp-all
            hence rtd-ab:  $r + t * (d\ (n+m)) \in \{a <.. < b\}$ 
            using r-ab rd-ab[of n]
          by simp (smt (verit, ccfv-threshold) mult-less-cancel-left mult-less-cancel-right2)
          have  $d\ (n+m) * \text{deriv } (\lambda s. f\ s\ x)\ (r + t * (d\ (n+m))) =$ 
             $d\ (n+m) * Df\ (r + t * (d\ (n+m)))\ x$ 
          proof -
            have  $r + t * (d\ (n+m)) \in \{a <.. < b\}$ 
            using r-ab rd-ab[of n] t-01
            by (smt (verit) ball-eq-greaterThanLessThan dist-real-def
greaterThanLessThan-eq-iff greaterThanLessThan-eq-ball mem-ball
mult-le-cancel-right1 mult-minus-right mult-pos-neg)
            thus ?thesis unfolding Df-def using x-CN by simp
          }
        }
      }
    }
  }

```

qed
with *f-MVT* **have** $(f (r + d (n+m)) x - f r x) / d (n+m) = Df (r + t * (d (n+m))) x$
using *d-neq0* **by** *simp*
moreover **have** $|Df (r + t * (d (n+m))) x| \leq g x$ **using** *Df-boundCN*
x-CN rtd-ab **by** *simp*
ultimately **have** $|(f (r + d (n+m)) x - f r x) / d (n+m)| \leq g x$ **by**
simp }
thus *AE x in M. norm* $((f (r + d (n+m)) x - f r x) / d (n+m)) \leq g x$
unfolding *real-norm-def* **using** *AE-I' N-null*
by $(smt (verit, ccfv-threshold) Diff-iff mem-Collect-eq subsetI)$
qed
ultimately **show** $((\lambda n. \int x. (f (r + d (n+m)) x - f r x) / d (n+m)) \partial M)$
 \longrightarrow
 $\int x. \lim (\lambda n. (f (r + d (n+m)) x - f r x) / d (n+m)) \partial M$
using *limf-msr fd-msr Df-bound*
by $(intro \textit{integral-dominated-convergence}[\mathbf{where} \ w=g], \textit{simp-all})$
qed
moreover **have** $\bigwedge n. ((\int x. f (r + d (n+m)) x \partial M) - \int x. f r x \partial M) / d (n+m) =$
 $\int x. (f (r + d (n+m)) x - f r x) / d (n+m) \partial M$
using *d-neq0* **apply** *simp*
by $(rewrite \textit{Bochner-Integration.integral-diff};$
 $(rule \textit{f-integ} | \textit{simp}); (rule \textit{rd-ab} | rule \textit{r-ab}))$
ultimately **show** *?thesis*
unfolding *comp-def* **using** *d-neq0*
apply -
by $(rule \textit{LIMSEQ-offset}[\mathbf{where} \ k=m]) \textit{simp}$
qed
moreover **have** $(\int x. \lim (\lambda n. (f (r + d (n+m)) x - f r x) / d (n+m)) \partial M)$
 $=$
 $\int x. \textit{borel-measurable-extension} \ M (\lambda y. \textit{deriv} (\lambda s. f s y) r) x \partial M$
proof -
have $(\int x. \lim (\lambda n. (f (r + d (n+m)) x - f r x) / d (n+m)) \partial M) = \int x. Df$
 $r x \partial M$
proof -
have *AE x in M. lim* $(\lambda n. (f (r + d (n+m)) x - f r x) / d (n+m)) = Df$
 $r x$
proof -
{ **fix** *x* **assume** *x-CN: x ∈ ?CN*
hence *lim* $(\lambda n. (f (r + d (n+m)) x - f r x) / d (n+m)) = Df r x$ **by**
 $(\textit{simp add: Df-eq})$ }
thus *?thesis* **using** *AE-I' N-null* **by** $(smt (verit, del-insts) DiffI mem-Collect-eq$
 $subsetI)$
qed
thus *?thesis* **using** *limf-msr Df-msr* **by** $(intro \textit{integral-cong-AE}; \textit{simp})$
qed
also **have** $\dots = \int x. \textit{borel-measurable-extension} \ M (\lambda y. \textit{deriv} (\lambda s. f s y) r)$
 $x \partial M$

proof –
have $AE\ x\ in\ M.\ Df\ r\ x = borel-measurable-extension\ M\ (\lambda y.\ deriv\ (\lambda s.\ f\ s\ y)\ r)\ x$ **and**
 $borel-measurable-extension\ M\ (\lambda y.\ deriv\ (\lambda s.\ f\ s\ y)\ r) \in borel-measurable\ M$
proof –
have $\{x \in space\ M.\ deriv\ (\lambda s.\ f\ s\ x)\ r \neq Df\ r\ x\} \subseteq N$
proof –
{ fix x **assume** $x \in ?CN$
hence $deriv\ (\lambda s.\ f\ s\ x)\ r = Df\ r\ x$ **unfolding** $Df-def$ **using** $r-ab$ **by**
 $simp$ **}**
thus $?thesis$ **by** $blast$
qed
thus $AE\ x\ in\ M.\ Df\ r\ x = borel-measurable-extension\ M\ (\lambda y.\ deriv\ (\lambda s.\ f\ s\ y)\ r)\ x$ **and**
 $borel-measurable-extension\ M\ (\lambda y.\ deriv\ (\lambda s.\ f\ s\ y)\ r) \in borel-measurable\ M$
using $Df-msr\ N-null$
apply –
apply $(rule\ borel-measurable-extensionI2[where\ S=N];\ simp)$
by $(rule\ borel-measurable-extension-measurable[where\ g=Df\ r];\ simp)$
qed
thus $?thesis$ **using** $Df-msr$ **by** $(intro\ integral-cong-AE;\ simp)$
qed
finally show $?thesis$.
qed
ultimately show $((\lambda h.\ ((\int x.\ f\ (r+h)\ x\ \partial M) - \int x.\ f\ r\ x\ \partial M) / h) \circ d)$
 \longrightarrow
 $\int x.\ borel-measurable-extension\ M\ (\lambda y.\ deriv\ (\lambda s.\ f\ s\ y)\ r)\ x\ \partial M$
using $tendsto-cong-limit$ **by** $simp$
qed
thus $(\lambda s.\ \int x.\ f\ s\ x\ \partial M)$ $has-real-derivative$
 $\int x.\ borel-measurable-extension\ M\ (\lambda y.\ deriv\ (\lambda s.\ f\ s\ y)\ r)\ x\ \partial M)$ $(at\ r)$
by $(rewrite\ DERIV-def,\ rewrite\ tendsto-at-iff-sequentially)\ simp$
qed

proposition $interchange-deriv-LINT$:
fixes $a\ b :: real$ **and** $f :: real \Rightarrow 'a \Rightarrow real$ **and** $g :: 'a \Rightarrow real$
assumes $\bigwedge r.\ r \in \{a <..<b\} \implies integrable\ M\ (f\ r)$ **and**
 $AE\ x\ in\ M.\ (\lambda r.\ f\ r\ x)$ $differentiable-on\ \{a <..<b\}$ **and**
 $\bigwedge r.\ r \in \{a <..<b\} \implies (\lambda x.\ (deriv\ (\lambda r.\ f\ r\ x)\ r)) \in borel-measurable\ M$ **and**
 $AE\ x\ in\ M.\ \forall r \in \{a <..<b\}.\ |deriv\ (\lambda r.\ f\ r\ x)\ r| \leq g\ x$ $integrable\ M\ g$
shows $\bigwedge r.\ r \in \{a <..<b\} \implies ((\lambda r.\ \int x.\ f\ r\ x\ \partial M)$ $has-real-derivative$
 $\int x.\ deriv\ (\lambda r.\ f\ r\ x)\ r\ \partial M)$ $(at\ r)$

proof –
fix r **assume** $r-ab: r \in \{a <..<b\}$
hence $Df-msr: (\lambda x.\ deriv\ (\lambda s.\ f\ s\ x)\ r) \in borel-measurable\ M$ **using** $assms$ **by**
 $simp$
have $(\lambda s.\ \int x.\ f\ s\ x\ \partial M)$ $has-real-derivative$

$\int x. \text{borel-measurable-extension } M (\lambda y. \text{deriv } (\lambda s. f s y) r) x \partial M) (at r)$
using *assms r-ab* **by** (*intro interchange-deriv-LINT-general; simp*)
moreover have $(\int x. \text{borel-measurable-extension } M (\lambda y. \text{deriv } (\lambda s. f s y) r) x \partial M) =$
 $\int x. \text{deriv } (\lambda s. f s x) r \partial M$
apply (*rule integral-cong-AE*)
apply (*rule borel-measurable-extension-measurable*
[**where** $g = \lambda y. \text{deriv } (\lambda s. f s y) r$ **and** $S = \{\}$], *simp-all add: Df-msr*)
using *borel-measurable-measurable-extension-AE Df-msr* **by** (*smt (verit) AE-cong*)
ultimately show $((\lambda r. \int x. f r x \partial M) \text{ has-real-derivative } \int x. \text{deriv } (\lambda r. f r x) r \partial M) (at r)$
by *simp*
qed

proposition *interchange-deriv-LINT-set-general:*

fixes $a b :: \text{real}$ **and** $f :: \text{real} \Rightarrow 'a \Rightarrow \text{real}$ **and** $g :: 'a \Rightarrow \text{real}$ **and** $A :: 'a \text{ set}$
assumes *A-msr: A ∈ sets M* **and**
f-integ: $\bigwedge r. r \in \{a < .. < b\} \implies \text{set-integrable } M A (f r)$ **and**
f-diff: AE x ∈ A in M. $(\lambda r. f r x)$ differentiable-on $\{a < .. < b\}$ **and**
Df-bound: AE x ∈ A in M. $\forall r \in \{a < .. < b\}. |\text{deriv } (\lambda r. f r x) r| \leq g x$ *set-integrable*
M A g

shows $\bigwedge r. r \in \{a < .. < b\} \implies ((\lambda r. \int x \in A. f r x \partial M) \text{ has-real-derivative } (\int x \in A. \text{set-borel-measurable-extension } M A (\lambda x. \text{deriv } (\lambda r. f r x) r) x \partial M)) (at r)$

proof –

let $?M-A = \text{restrict-space } M A$
have $\bigwedge r. r \in \{a < .. < b\} \implies \text{integrable } ?M-A (f r)$
using *A-msr f-integ set-integrable-restrict-space-iff* **by** *auto*
moreover have *AE x in ?M-A. $(\lambda r. f r x)$ differentiable-on $\{a < .. < b\}$*
using *AE-restrict-space-iff A-msr f-diff* **by** (*metis sets.Int-space-eq2*)
moreover have *AE x in ?M-A. $\forall r \in \{a < .. < b\}. |\text{deriv } (\lambda r. f r x) r| \leq g x$* **and**
integrable ?M-A g
using *A-msr Df-bound set-integrable-restrict-space-iff*
apply –
by (*simp add: AE-restrict-space-iff, auto*)
ultimately have $\bigwedge r. r \in \{a < .. < b\} \implies ((\lambda r. \text{integral}^L ?M-A (f r)) \text{ has-real-derivative } \text{integral}^L ?M-A (\text{borel-measurable-extension } ?M-A (\lambda x. \text{deriv } (\lambda r. f r x) r))) (at r)$
by (*rule interchange-deriv-LINT-general*[**where** $M = \text{restrict-space } M A$]) *auto*
thus $\bigwedge r. r \in \{a < .. < b\} \implies ((\lambda r. \int x \in A. f r x \partial M) \text{ has-real-derivative } (\int x \in A. \text{set-borel-measurable-extension } M A (\lambda x. \text{deriv } (\lambda r. f r x) r) x \partial M)) (at r)$
unfolding *set-borel-measurable-extension-def* **using** *assms*
by (*rewrite set-lebesgue-integral-restrict-space, simp*)
qed

proposition *interchange-deriv-LINT-set:*

fixes $a b :: \text{real}$ **and** $f :: \text{real} \Rightarrow 'a \Rightarrow \text{real}$ **and** $g :: 'a \Rightarrow \text{real}$ **and** $A :: 'a \text{ set}$
assumes $A \in \text{sets } M$ **and**

$\bigwedge r. r \in \{a < .. < b\} \implies \text{set-integrable } M A (f r)$ **and**
 $AE x \in A \text{ in } M. (\lambda r. f r x) \text{ differentiable-on } \{a < .. < b\}$ **and**
 $\bigwedge r. r \in \{a < .. < b\} \implies \text{set-borel-measurable } M A (\lambda x. (\text{deriv } (\lambda r. f r x) r))$ **and**
 $AE x \in A \text{ in } M. \forall r \in \{a < .. < b\}. |\text{deriv } (\lambda r. f r x) r| \leq g x \text{ set-integrable } M A g$
shows $\bigwedge r. r \in \{a < .. < b\} \implies ((\lambda r. \int x \in A. f r x \partial M) \text{ has-real-derivative } (\int x \in A. \text{deriv } (\lambda r. f r x) r \partial M)) \text{ (at } r)$
proof –
fix r **assume** $r\text{-ab}: r \in \{a < .. < b\}$
hence $Df\text{-msr}: \text{set-borel-measurable } M A (\lambda x. \text{deriv } (\lambda s. f s x) r)$ **using** $assms$
by $simp$
have $((\lambda s. \int x \in A. f s x \partial M) \text{ has-real-derivative } (\int x \in A. \text{set-borel-measurable-extension } M A (\lambda y. \text{deriv } (\lambda s. f s y) r) x \partial M)) \text{ (at } r)$
using $assms\ r\text{-ab}$ **by** $(\text{intro interchange-deriv-LINT-set-general}; simp)$
moreover **have** $(\int x \in A. \text{set-borel-measurable-extension } M A (\lambda y. \text{deriv } (\lambda s. f s y) r) x \partial M) = (\int x \in A. \text{deriv } (\lambda s. f s x) r \partial M)$
apply $(\text{rule set-lebesgue-integral-cong-AE2}, simp\ add: assms)$
apply $(\text{rule set-borel-measurable-extension-measurable } [\text{where } g = \lambda y. \text{deriv } (\lambda s. f s y) r \text{ and } S = \{\}], simp\text{-all add: } Df\text{-msr } assms)$
using $\text{set-borel-measurable-measurable-extension-AE } Df\text{-msr } assms$ **by** $(\text{smt (verit) AE-cong})$
ultimately show
 $((\lambda r. \int x \in A. f r x \partial M) \text{ has-real-derivative } (\int x \in A. \text{deriv } (\lambda r. f r x) r \partial M)) \text{ (at } r)$
by $simp$
qed

3 Additional Lemmas for the *HOL-Probability* Library

lemma $(\text{in } \text{finite-borel-measure})$
fixes $F :: \text{real} \Rightarrow \text{real}$
assumes $\text{nondec } F : \bigwedge x y. x \leq y \implies F x \leq F y$ **and**
 $\text{right-cont-}F : \bigwedge a. \text{continuous (at-right } a) F$ **and**
 $\text{lim-}F\text{-at-bot} : (F \longrightarrow 0) \text{ at-bot}$ **and**
 $\text{lim-}F\text{-at-top} : (F \longrightarrow m) \text{ at-top}$ **and**
 $m : 0 \leq m$
shows $\text{emeasure-interval-measure-Ioi: emeasure (interval-measure } F) \{x < ..\} = m - F x$
and $\text{measure-interval-measure-Ioi: measure (interval-measure } F) \{x < ..\} = m - F x$
proof –
interpret $F\text{-FM}: \text{finite-measure interval-measure } F$
using $\text{finite-borel-measure.axioms(1) finite-borel-measure-interval-measure lim-}F\text{-at-bot lim-}F\text{-at-top } m \text{ nondec } F \text{ right-cont-}F$ **by** $blast$
have $UNIV = \{..x\} \cup \{x < ..\}$ **by** $auto$
moreover **have** $\{..x\} \cap \{x < ..\} = \{\}$ **by** $auto$

ultimately have $\text{emeasure (interval-measure } F) \text{ UNIV} =$
 $\text{emeasure (interval-measure } F) \{..x\} + \text{emeasure (interval-measure } F) \{x<..\}$
by (*simp add: plus-emeasure*)
moreover have $\text{emeasure (interval-measure } F) \text{ UNIV} = m$
using *assms interval-measure-UNIV by presburger*
ultimately show $\star: \text{emeasure (interval-measure } F) \{x<..\} = m - F x$
using *assms emeasure-interval-measure-Iic*
by (*metis ennreal-add-diff-cancel-left ennreal-minus measure-interval-measure-Iic*
measure-nonneg top-neq-ennreal)
hence $\text{ennreal (measure (interval-measure } F) \{x<..\}) = m - F x$
using *emeasure-eq-measure by (metis emeasure-eq-ennreal-measure top-neq-ennreal)*
moreover have $\bigwedge x. F x \leq m$
using *lim-F-at-top nondecF by (intro mono-at-top-le[where f=F]; simp add:*
mono-def)
ultimately show $\text{measure (interval-measure } F) \{x<..\} = m - F x$
using *ennreal-inj F-FM.emeasure-eq-measure by force*
qed

lemma (in prob-space) cond-prob-nonneg[*simp*]: $\text{cond-prob } M P Q \geq 0$
by (*auto simp: cond-prob-def*)

lemma (in prob-space) cond-prob-whole-1: $\text{cond-prob } M P P = 1$ **if** $\text{prob } \{\omega \in$
 $\text{space } M. P \omega\} \neq 0$
unfolding *cond-prob-def using that by simp*

lemma (in prob-space) cond-prob-0-null: $\text{cond-prob } M P Q = 0$ **if** $\text{prob } \{\omega \in \text{space}$
 $M. Q \omega\} = 0$
unfolding *cond-prob-def using that by simp*

lemma (in prob-space) cond-prob-AE-prob:
assumes $\{\omega \in \text{space } M. P \omega\} \in \text{events } \{\omega \in \text{space } M. Q \omega\} \in \text{events}$
and *AE ω in $M. Q \omega$*
shows $\text{cond-prob } M P Q = \text{prob } \{\omega \in \text{space } M. P \omega\}$
proof –
let $?setP = \{\omega \in \text{space } M. P \omega\}$
let $?setQ = \{\omega \in \text{space } M. Q \omega\}$
have [*simp*]: $\text{prob } ?setQ = 1$ **using** *assms prob-Collect-eq-1 by simp*
hence $\text{cond-prob } M P Q = \text{prob } (?setP \cap ?setQ)$
unfolding *cond-prob-def by (simp add: Collect-conj-eq2)*
also have $\dots = \text{prob } ?setP$
proof (*rule antisym*)
show $\text{prob } (?setP \cap ?setQ) \leq \text{prob } ?setP$
using *assms finite-measure-mono inf-sup-ord(1) by blast*
next
show $\text{prob } ?setP \leq \text{prob } (?setP \cap ?setQ)$
proof –
have $\text{prob } (?setP \cap ?setQ) = \text{prob } ?setP + \text{prob } ?setQ - \text{prob } (?setP \cup$
 $?setQ)$
using *assms by (smt (verit) finite-measure-Diff' finite-measure-Union'*

```

sup-commute)
  also have ... = prob ?setP + (1 - prob (?setP ∪ ?setQ)) by simp
  also have ... ≥ prob ?setP by simp
  finally show ?thesis .
qed
qed
finally show ?thesis .
qed

```

3.1 More Properties of *cdf*'s

```

context finite-borel-measure
begin

```

```

lemma cdf-diff-eq2:
  assumes  $x \leq y$ 
  shows  $\text{cdf } M \ y - \text{cdf } M \ x = \text{measure } M \ \{x < .. y\}$ 
proof (cases ⟨ $x = y$ ⟩)
  case True
  thus ?thesis by force
next
  case False
  hence  $x < y$  using assms by simp
  thus ?thesis by (rule cdf-diff-eq)
qed

```

```
end
```

```

context prob-space
begin

```

```

lemma cdf-distr-measurable [measurable]:
  assumes [measurable]: random-variable borel X
  shows  $\text{cdf } (\text{distr } M \ \text{borel } X) \in \text{borel-measurable borel}$ 
proof (rule borel-measurable-mono)
  show mono (cdf (distr M borel X))
  unfolding mono-def
  using finite-borel-measure.cdf-nondecreasing
  by (simp add: real-distribution.finite-borel-measure-M)
qed

```

```

lemma cdf-distr-P:
  assumes random-variable borel X
  shows  $\text{cdf } (\text{distr } M \ \text{borel } X) \ x = \mathcal{P}(\omega \text{ in } M. X \ \omega \leq x)$ 
  unfolding cdf-def apply (rewrite measure-distr; (simp add: assms)?)
  unfolding vimage-def by (rule arg-cong[where f=prob], force)

```

```

lemma cdf-continuous-distr-P-lt:
  assumes random-variable borel X isCont (cdf (distr M borel X)) x

```



```

shows cdf (distr M borel X) x =  $\mathcal{P}(\omega \text{ in } M. X \omega < x)$ 
proof -
have  $\mathcal{P}(\omega \text{ in } M. X \omega < x) = \text{measure } (distr M borel X) \{..<x\}$ 
  apply (rewrite measure-distr, simp-all add: assms)
  unfolding vimage-def by simp (smt (verit) Collect-cong Int-def mem-Collect-eq)
also have  $\dots = \text{measure } (distr M borel X) (\{..<x\} \cup \{x\})$ 
  apply (rewrite finite-measure.measure-zero-union, simp-all add: assms finite-measure-distr)
  using finite-borel-measure.isCont-cdf real-distribution.finite-borel-measure-M
  assms by blast
also have  $\dots = \text{measure } (distr M borel X) \{..x\}$  by (metis ivl-disj-un-singleton(2))
also have  $\dots = \text{cdf } (distr M borel X) x$  unfolding cdf-def by simp
finally show ?thesis by simp
qed

```

lemma cdf-distr-diff-P:

```

assumes  $x \leq y$ 
  and random-variable borel X
shows cdf (distr M borel X) y - cdf (distr M borel X) x =  $\mathcal{P}(\omega \text{ in } M. x < X \omega$ 
 $\wedge X \omega \leq y)$ 
proof -
interpret distrX-FBM: finite-borel-measure distr M borel X
  using real-distribution.finite-borel-measure-M real-distribution-distr assms by
  simp
have cdf (distr M borel X) y - cdf (distr M borel X) x = measure (distr M borel
X)  $\{x<..y\}$ 
  by (rewrite distrX-FBM.cdf-diff-eq2; simp add: assms)
also have  $\dots = \mathcal{P}(\omega \text{ in } M. x < X \omega \wedge X \omega \leq y)$ 
  apply (rewrite measure-distr; (simp add: assms) ?)
  unfolding vimage-def by (rule arg-cong[where f=prob], force)
finally show ?thesis .
qed

```

lemma cdf-distr-max:

```

fixes c::real
assumes [measurable]: random-variable borel X
shows cdf (distr M borel ( $\lambda x. \max (X x) c$ )) u = cdf (distr M borel X) u *
indicator  $\{c..\} u$ 
proof (cases  $\langle c \leq u \rangle$ )
case True
  thus ?thesis
    unfolding cdf-def
    apply (rewrite measure-distr; simp?)+
    by (smt (verit) Collect-cong atMost-iff vimage-def)
next
case False
  thus ?thesis
    unfolding cdf-def
    apply (rewrite measure-distr; simp?)+
    by (smt (verit, best) Int-emptyI atMost-iff measure-empty vimage-eq)

```

qed

lemma *cdf-distr-min:*

```
fixes c::real
assumes [measurable]: random-variable borel X
shows cdf (distr M borel ( $\lambda x. \min (X x) c$ )) u =
  1 - (1 - cdf (distr M borel X) u) * indicator {.. $c$ } u
proof (cases  $\langle c \leq u \rangle$ )
  case True
  thus ?thesis
    unfolding cdf-def
    using real-distribution.finite-borel-measure-M real-distribution-distr
    apply (rewrite measure-distr; simp?)
    by (smt (verit, del-Insts) Int-absorb1 atMost-iff prob-space subset-vimage-iff)
  next
  case False
  thus ?thesis
    unfolding cdf-def
    using real-distribution.finite-borel-measure-M real-distribution-distr
    apply (rewrite measure-distr; simp?)+
    using prob-space-axioms assms
    by (smt (verit) Collect-cong Int-def atMost-iff prob-space prob-space.cdf-distr-P
vimage-eq)
qed
```

lemma *cdf-distr-floor-P:*

```
fixes X :: 'a  $\Rightarrow$  real
assumes [measurable]: random-variable borel X
shows cdf (distr M borel ( $\lambda x. \lfloor X x \rfloor$ )) u =  $\mathcal{P}(x \text{ in } M. X x < \lfloor u \rfloor + 1)$ 
unfolding cdf-def
apply (rewrite measure-distr; simp?)
apply (rule arg-cong[where f=prob])
unfolding vimage-def using floor-le-iff le-floor-iff by blast
```

lemma *expectation-nonneg-tail:*

```
assumes [measurable]: random-variable borel X
and X-nonneg:  $\bigwedge x. x \in \text{space } M \implies X x \geq 0$ 
defines F u  $\equiv$  cdf (distr M borel X) u
shows ( $\int^+ x. \text{ennreal } (X x) \partial M$ ) = ( $\int^+ u \in \{0..\}. \text{ennreal } (1 - F u) \partial \text{lborel}$ )
proof -
let ?distrX = distr M borel X
have ( $\int^+ x. \text{ennreal } (X x) \partial M$ ) = ( $\int^+ u. \text{ennreal } u \partial ?\text{distrX}$ )
  by (rewrite nn-integral-distr; simp)
also have ... = ( $\int^+ u \in \{0..\}. \text{ennreal } u \partial ?\text{distrX}$ )
  by (rule nn-integral-distr-set; simp add: X-nonneg)
also have ... = ( $\int^+ u \in \{0..\}. (\int^+ v \in \{0..\}. \text{indicator } \{..<u\} v \partial \text{lborel}) \partial ?\text{distrX}$ )
proof -
have  $\bigwedge u::\text{real}. u \in \{0..\} \implies \text{ennreal } u = (\int^+ v \in \{0..\}. \text{indicator } \{..<u\} v \partial \text{lborel})$ 
  apply (rewrite indicator-inter-arith[THEN sym])
```

apply (*rewrite nn-integral-indicator, measurable, simp*)
by (*metis atLeastLessThan-def diff-zero emeasure-lborel-Ico inf commute*)
thus *?thesis* **by** (*metis (no-types, lifting) indicator-eq-0-iff mult-eq-0-iff*)
qed
also have $\dots = (\int^+ v \in \{0..\}. (\int^+ u \in \{0..\}. \text{indicator } \{.. < u\} v \partial^? \text{distr} X) \partial \text{lborel})$
proof –
have $(\int^+ u \in \{0..\}. (\int^+ v \in \{0..\}. \text{indicator } \{.. < u\} v \partial \text{lborel}) \partial^? \text{distr} X) =$
 $\int^+ u. (\int^+ v. \text{indicator } \{.. < u\} v * \text{indicator } \{0..\} v * \text{indicator } \{0..\} u \partial \text{lborel})$
 $\partial^? \text{distr} X$
by (*rewrite nn-integral-multc; simp*)
also have $\dots =$
 $\int^+ v. (\int^+ u. \text{indicator } \{.. < u\} v * \text{indicator } \{0..\} v * \text{indicator } \{0..\} u$
 $\partial^? \text{distr} X) \partial \text{lborel}$
apply (*rewrite pair-sigma-finite.Fubini'; simp?*)
using *pair-sigma-finite.intro assms*
prob-space-distr prob-space-imp-sigma-finite sigma-finite-lborel
apply *blast*
by *measurable auto*
also have $\dots = (\int^+ v \in \{0..\}. (\int^+ u \in \{0..\}. \text{indicator } \{.. < u\} v \partial^? \text{distr} X)$
 $\partial \text{lborel})$
apply (*rewrite nn-integral-multc[THEN sym]; measurable; simp?*)
apply (*rule nn-integral-cong*)
using *mult.assoc mult.commute* **by** *metis*
finally show *?thesis* **by** *simp*
qed
also have $\dots = (\int^+ v \in \{0..\}. (\int^+ u. \text{indicator } \{v < ..\} u \partial^? \text{distr} X) \partial \text{lborel})$
apply (*rule nn-integral-cong*)
apply (*rewrite nn-integral-multc[THEN sym], measurable; (simp del: nn-integral-indicator)?*)
apply (*rule nn-integral-cong*)
using *lessThan-iff greaterThan-iff atLeast-iff indicator-simps*
by (*smt (verit, del-insts) mult-1 mult-eq-0-iff*)
also have $\dots = (\int^+ v \in \{0..\}. \text{ennreal } (1 - F v) \partial \text{lborel})$
apply (*rule nn-integral-cong, simp*)
apply (*rewrite emeasure-distr; simp?*)
apply (*rewrite vimage-compl-atMost[THEN sym]*)
unfolding *F-def cdf-def*
apply (*rewrite measure-distr; simp?*)
apply (*rewrite prob-compl[THEN sym], simp*)
by (*metis (no-types, lifting) Diff-Compl Diff-Diff-Int Int-commute emeasure-eq-measure*)
finally show *?thesis* .
qed

lemma *expectation-nonpos-tail:*

assumes [*measurable*]: *random-variable borel X*
and *X-nonpos*: $\bigwedge x. x \in \text{space } M \implies X x \leq 0$
defines $F u \equiv \text{cdf } (\text{distr } M \text{ borel } X) u$
shows $(\int^+ x. \text{ennreal } (- X x) \partial M) = (\int^+ u \in \{..0\}. \text{ennreal } (F u) \partial \text{lborel})$
proof –
let $? \text{distr} X = \text{distr } M \text{ borel } X$

have $(\int^+ x. \text{ennreal } (- X x) \partial M) = (\int^+ u. \text{ennreal } (-u) \partial^? \text{distr} X)$
by *(rewrite nn-integral-distr; simp)*
also have $\dots = (\int^+ u \in \{..0\}. \text{ennreal } (-u) \partial^? \text{distr} X)$
proof –
have $[\text{simp}]: \{..0::\text{real}\} \cup \{0<..\} = \text{UNIV}$ **by** *force*
have $(\int^+ u. \text{ennreal } (-u) \partial^? \text{distr} X) =$
 $(\int^+ u \in \{..0\}. \text{ennreal } (-u) \partial^? \text{distr} X) + (\int^+ u \in \{0<..\}. \text{ennreal } (-u) \partial^? \text{distr} X)$
by *(rewrite nn-integral-disjoint-pair[THEN sym], simp-all, force)*
also have $\dots = (\int^+ u \in \{..0\}. \text{ennreal } (-u) \partial^? \text{distr} X)$
apply *(rewrite nn-integral-zero'[of $\lambda u. \text{ennreal } (-u) * \text{indicator } \{0<..\} u$]; simp?)*
unfolding *indicator-def* **using** *always-eventually ennreal-lt-0* **by** *force*
finally show *?thesis* .
qed
also have $\dots = (\int^+ u \in \{..0\}. (\int^+ v \in \{..0\}. \text{indicator } \{u..\} v \partial \text{lborel}) \partial^? \text{distr} X)$
proof –
have $\bigwedge u::\text{real}. u \in \{..0\} \implies \text{ennreal } (-u) = (\int^+ v \in \{..0\}. \text{indicator } \{u..\} v \partial \text{lborel})$
by *(rewrite indicator-inter-arith[THEN sym]) simp*
thus *?thesis* **by** *(metis (no-types, lifting) indicator-eq-0-iff mult-eq-0-iff)*
qed
also have $\dots = (\int^+ v \in \{..0\}. (\int^+ u \in \{..0\}. \text{indicator } \{u..\} v \partial^? \text{distr} X) \partial \text{lborel})$
proof –
have $(\int^+ u \in \{..0\}. (\int^+ v \in \{..0\}. \text{indicator } \{u..\} v \partial \text{lborel}) \partial^? \text{distr} X) =$
 $\int^+ u. (\int^+ v. \text{indicator } \{u..\} v * \text{indicator } \{..0\} v * \text{indicator } \{..0\} u \partial \text{lborel})$
 $\partial^? \text{distr} X$
by *(rewrite nn-integral-multc; simp)*
also have $\dots =$
 $\int^+ v. (\int^+ u. \text{indicator } \{u..\} v * \text{indicator } \{..0\} v * \text{indicator } \{..0\} u \partial^? \text{distr} X)$
 ∂lborel
apply *(rewrite pair-sigma-finite.Fubini'; simp?)*
using *pair-sigma-finite.intro assms*
prob-space-distr prob-space-imp-sigma-finite sigma-finite-lborel
apply *blast*
by *measurable auto*
also have $\dots = (\int^+ v \in \{..0\}. (\int^+ u \in \{..0\}. \text{indicator } \{u..\} v \partial^? \text{distr} X) \partial \text{lborel})$
apply *(rewrite nn-integral-multc[THEN sym]; measurable; simp?)*
apply *(rule nn-integral-cong)+*
using *mult.assoc mult.commute* **by** *metis*
finally show *?thesis* **by** *simp*
qed
also have $\dots = (\int^+ v \in \{..0\}. (\int^+ u. \text{indicator } \{..v\} u \partial^? \text{distr} X) \partial \text{lborel})$
apply *(rule nn-integral-cong)*
apply *(rewrite nn-integral-multc[THEN sym], measurable; (simp del: nn-integral-indicator)?)+*
apply *(rule nn-integral-cong)*
using *atMost-iff atLeast-iff indicator-simps* **by** *(smt (verit, del-insts) mult-1 mult-eq-0-iff)*
also have $\dots = (\int^+ v \in \{..0\}. \text{ennreal } (F v) \partial \text{lborel})$
apply *(rule nn-integral-cong, simp)*

```

apply (rewrite emeasure-distr; simp?)
unfolding F-def cdf-def
by (rewrite measure-distr; simp add: emeasure-eq-measure)
finally show ?thesis .
qed

proposition expectation-tail:
assumes [measurable]: integrable M X
defines F u  $\equiv$  cdf (distr M borel X) u
shows expectation X = (LBINT u:{0..}. 1 - F u) - (LBINT u:{..0}. F u)
proof -
have expectation X = expectation ( $\lambda x$ . max (X x) 0) + expectation ( $\lambda x$ . min (X
x) 0)
using integrable-max integrable-min
apply (rewrite Bochner-Integration.integral-add[THEN sym], measurable)
by (rule Bochner-Integration.integral-cong; simp)
also have ... = expectation ( $\lambda x$ . max (X x) 0) - expectation ( $\lambda x$ . - min (X x)
0) by force
also have ... = (LBINT u:{0..}. 1 - F u) - (LBINT u:{..0}. F u)
proof -
have expectation ( $\lambda x$ . max (X x) 0) = (LBINT u:{0..}. 1 - F u)
proof -
have expectation ( $\lambda x$ . max (X x) 0) = enn2real ( $\int^+ x$ . ennreal (max (X x)
0)  $\partial M$ )
by (rule integral-eq-nn-integral; simp)
also have ... = enn2real ( $\int^+ u \in \{0..\}$ . ennreal (1 - F u)  $\partial$ borel)
apply (rewrite expectation-nonneg-tail; simp?)
apply (rewrite cdf-distr-max, simp)
unfolding F-def
by (metis (opaque-lifting) indicator-simps mult.commute mult-1 mult-eq-0-iff)
also have ... = enn2real ( $\int^+ u$ . ennreal ((1 - F u) * indicator {0..} u)
 $\partial$ borel)
by (simp add: indicator-mult-ennreal mult.commute)
also have ... = (LBINT u:{0..}. 1 - F u)
apply (rewrite integral-eq-nn-integral[THEN sym], simp add: F-def)
unfolding F-def using real-distribution.cdf-bounded-prob apply force
unfolding set-lebesgue-integral-def by (rule Bochner-Integration.integral-cong;
simp)
finally show ?thesis .
qed
moreover have expectation ( $\lambda x$ . - min (X x) 0) = (LBINT u:{..0}. F u)
proof -
have expectation ( $\lambda x$ . - min (X x) 0) = enn2real ( $\int^+ x$ . ennreal (- min (X
x) 0)  $\partial M$ )
by (rule integral-eq-nn-integral; simp)
also have ... = enn2real ( $\int^+ u \in \{..0\}$ . ennreal (F u)  $\partial$ borel)
proof -
let ?distrminX = distr M borel ( $\lambda x$ . min (X x) 0)
have [simp]: sym-diff {..0} {..<0} = {0::real} by force

```

```

    have enn2real (∫+x. ennreal (- min (X x) 0) ∂M) =
enn2real (∫+u∈{..0}. ennreal (cdf ?distrminX u) ∂lborel)
    by (rewrite expectation-nonpos-tail; simp)
  also have ... = enn2real (∫+u∈{..<0}. ennreal (cdf ?distrminX u) ∂lborel)
    by (rewrite nn-integral-null-delta, auto)
  also have ... = enn2real (∫+u∈{..<0}. ennreal (F u) ∂lborel)
    apply (rewrite cdf-distr-min, simp)
    apply (rule arg-cong[where f=enn2real], rule nn-integral-cong)
  unfolding F-def by (smt (verit) indicator-simps mult-cancel-left1 mult-eq-0-iff)
  also have ... = enn2real (∫+u∈{..0}. ennreal (F u) ∂lborel)
    by (rewrite nn-integral-null-delta, auto simp add: sup-commute)
  finally show ?thesis .
qed
also have ... = enn2real (∫+u. ennreal (F u * indicator {..0} u) ∂lborel)
  using mult.commute indicator-mult-ennreal by metis
also have ... = (LBINT u:{..0}. F u)
  apply (rewrite integral-eq-nn-integral[THEN sym], simp add: F-def)
  unfolding F-def
  using finite-borel-measure.cdf-nonneg real-distribution.finite-borel-measure-M
apply simp
  unfolding set-lebesgue-integral-def by (rule Bochner-Integration.integral-cong;
simp)
  finally show ?thesis .
qed
  ultimately show ?thesis by simp
qed
  finally show ?thesis .
qed

```

proposition *distributed-deriv-cdf*:

```

  assumes [measurable]: random-variable borel X
  defines F u ≡ cdf (distr M borel X) u
  assumes finite S ∧ x. x ∉ S ⇒ (F has-real-derivative f x) (at x)
    and continuous-on UNIV F f ∈ borel-measurable lborel
  shows distributed M lborel X f

```

proof –

```

  have FBM: finite-borel-measure (distr M borel X)
    using real-distribution.finite-borel-measure-M real-distribution-distr assms by
simp
  then interpret distrX-FBM: finite-borel-measure distr M borel X .
  have FBML: finite-borel-measure (distr M lborel X) using FBM distr-borel-lborel
by smt
  then interpret distrLX-FBM: finite-borel-measure distr M lborel X .
  have [simp]: (λx. ennreal (f x)) ∈ borel-measurable borel using assms by simp
  moreover have distr M lborel X = density lborel f

```

proof –

```

  have ∧ a b. a ≤ b ⇒ emeasure (distr M lborel X) {a<..b} < ⊤
    using distrLX-FBM.emeasure-real less-top-ennreal by blast
  moreover have ∧ a b. a ≤ b ⇒

```

```

    emeasure (distr M lborel X) {a<..b} = emeasure (density lborel f) {a<..b}
  proof -
    fix a b :: real assume a ≤ b
    hence [simp]: sym-diff {a<..b} {a..b} = {a} by force
    have emeasure (density lborel f) {a<..b} = (∫+x∈{a<..b}. ennreal (f x)
    ∂lborel)
      by (rewrite emeasure-density; simp)
    also have ... = (∫+x∈{a..b}. ennreal (f x) ∂lborel) by (rewrite nn-integral-null-delta,
    auto)
    also have ... = ∫+x. ennreal (indicat-real {a..b} x * f x) ∂lborel
      by (metis indicator-mult-ennreal mult commute)
    also have ... = ennreal (F b - F a)
  proof -
    define g where g x = (if x ∈ S then 0 else f x) for x :: real
    have [simp]: ∧x. g x ≥ 0
      unfolding g-def
      apply (split if-split, auto)
      apply (rule mono-on-imp-deriv-nonneg[of UNIV F], auto)
    unfolding F-def mono-on-def using distrX-FBM.cdf-nondecreasing apply
    blast
      using assms unfolding F-def by force
    have (∫+x. ennreal (indicat-real {a..b} x * f x) ∂lborel)
      = ∫+x. ennreal (indicat-real {a..b} x * g x) ∂lborel
      apply (rule nn-integral-cong-AE)
      apply (rule AE-mp[where P= λx. x ∉ S])
      using assms finite-imp-null-set-lborel AE-not-in apply blast
      unfolding g-def by simp
    also have ... = ennreal (F b - F a)
      apply (rewrite nn-integral-has-integral-lebesgue, simp)
      apply (rule fundamental-theorem-of-calculus-strong[of S], auto simp: ⟨a
    ≤ b⟩ g-def assms)
      using has-real-derivative-iff-has-vector-derivative assms apply presburger
      using assms continuous-on-subset subsetI by fastforce
      finally show ?thesis .
  qed
  also have ... = emeasure (distr M lborel X) {a <.. b}
    apply (rewrite distrX-FBM.emeasure-Ioc, simp add: ⟨a ≤ b⟩)
    unfolding F-def cdf-def
    apply (rewrite ennreal-minus[THEN sym], simp)+
    by (metis distr-borel-lborel)
  finally show emeasure (distr M lborel X) {a<..b} = emeasure (density lborel
  f) {a<..b}
    by simp
  qed
  ultimately show ?thesis by (intro measure-eqI-Ioc; simp)
  qed
  ultimately show ?thesis unfolding distributed-def by auto
  qed

```

end

Define the conditional probability space. This is obtained by rescaling the restricted space of a probability space.

3.2 Conditional Probability Space

```
lemma restrict-prob-space:  
  assumes measure-space  $\Omega$   $A$   $\mu$   $a \in A$   
    and  $0 < \mu a$   $\mu a < \infty$   
  shows prob-space (scale-measure ( $1 / \mu a$ ) (restrict-space (measure-of  $\Omega$   $A$   $\mu$ )  
   $a$ ))  
proof  
  let  $?M = \text{measure-of } \Omega \ A \ \mu$   
  let  $?Ma = \text{restrict-space } ?M \ a$   
  let  $?rMa = \text{scale-measure } (1 / \mu a) \ ?Ma$   
  have emeasure  $?rMa$  (space  $?rMa$ ) =  $1 / \mu a * \text{emeasure } ?Ma$  (space  $?rMa$ ) by  
  simp  
  also have  $\dots = 1 / \mu a * \text{emeasure } ?M$  (space  $?rMa$ )  
    using assms  
    apply (rewrite emeasure-restrict-space)  
    apply (simp add: measure-space-def sigma-algebra.sets-measure-of-eq)  
    by (simp add: space-restrict-space space-scale-measure) simp  
  also have  $\dots = 1 / \mu a * \text{emeasure } ?M$  (space  $?Ma$ ) by (rewrite space-scale-measure)  
  simp  
  also have  $\dots = 1 / \mu a * \text{emeasure } ?M \ a$   
    using assms  
    apply (rewrite space-restrict-space2)  
    by (simp add: measure-space-closed)  
  also have  $\dots = 1$   
    using assms measure-space-def  
    apply (rewrite emeasure-measure-of-sigma, blast+)  
    by (simp add: ennreal-divide-times)  
  finally show emeasure  $?rMa$  (space  $?rMa$ ) =  $1$  .  
qed
```

```
definition cond-prob-space :: 'a measure  $\Rightarrow$  'a set  $\Rightarrow$  'a measure (infix  $\lfloor$  200)  
  where  $M \lfloor A \equiv \text{scale-measure } (1 / \text{emeasure } M \ A) \ (\text{restrict-space } M \ A)$ 
```

```
context prob-space  
begin
```

```
lemma cond-prob-space-whole[simp]:  $M \lfloor \text{space } M = M$   
  unfolding cond-prob-space-def using emeasure-space-1 by simp
```

```
lemma cond-prob-space-correct:  
  assumes  $A \in \text{events } \text{prob } A > 0$   
  shows prob-space ( $M \lfloor A$ )  
  unfolding cond-prob-space-def
```


apply (*subst*(2) *measure-of-of-measure*[of M , *THEN sym*])
using *assms*
by (*intro restrict-prob-space*; (*simp add: measure-space*)?; *simp-all add: emeasure-eq-measure*)

lemma *space-cond-prob-space*:
assumes $A \in \text{events}$
shows $\text{space } (M \downarrow A) = A$
unfolding *cond-prob-space-def* **using** *assms* **by** (*simp add: space-scale-measure*)

lemma *sets-cond-prob-space*: $\text{sets } (M \downarrow A) = (\cap) A$ ‘ *events*
unfolding *cond-prob-space-def* **by** (*metis sets-restrict-space sets-scale-measure*)

lemma *measure-cond-prob-space*:
assumes $A \in \text{events}$ $B \in \text{events}$
and $\text{prob } A > 0$
shows $\text{measure } (M \downarrow A) (B \cap A) = \text{prob } (B \cap A) / \text{prob } A$
proof –
have $1 / \text{emeasure } M A = \text{ennreal } (1 / \text{prob } A)$
using *assms* **by** (*smt (verit) divide-ennreal emeasure-eq-measure ennreal-1*)
hence $\text{measure } (M \downarrow A) (B \cap A) = (1 / \text{prob } A) * \text{measure } (\text{restrict-space } M A) (B \cap A)$
unfolding *cond-prob-space-def* **using** *measure-scale-measure* **by force**
also have $\dots = (1 / \text{prob } A) * \text{prob } (B \cap A)$
using *measure-restrict-space assms* **by** (*metis inf.cobounded2 sets.Int-space-eq2*)
also have $\dots = \text{prob } (B \cap A) / \text{prob } A$ **by** *simp*
finally show *?thesis* .

qed

corollary *measure-cond-prob-space-subset*:
assumes $A \in \text{events}$ $B \in \text{events}$ $B \subseteq A$
and $\text{prob } A > 0$
shows $\text{measure } (M \downarrow A) B = \text{prob } B / \text{prob } A$
proof –
have $B = B \cap A$ **using** *assms* **by auto**
moreover have $\text{measure } (M \downarrow A) (B \cap A) = \text{prob } (B \cap A) / \text{prob } A$
using *assms measure-cond-prob-space* **by** *simp*
ultimately show *?thesis* **by auto**

qed

lemma *cond-cond-prob-space*:
assumes $A \in \text{events}$ $B \in \text{events}$ $B \subseteq A$ $\text{prob } B > 0$
shows $(M \downarrow A) \downarrow B = M \downarrow B$
proof (*rule measure-eqI*)
have $pA\text{-pos}[simp]$: $\text{prob } A > 0$ **using** *assms* **by** (*smt (verit, ccfv-SIG) finite-measure-mono*)
interpret *MA-PS*: *prob-space* $M \downarrow A$ **using** *cond-prob-space-correct assms* **by** *simp*
interpret *MB-PS*: *prob-space* $M \downarrow B$ **using** *cond-prob-space-correct assms* **by** *simp*
have $1 / \text{emeasure } M A = \text{ennreal } (1 / \text{prob } A)$

using pA -pos **by** (smt (verit, ccfv-SIG) divide-ennreal emeasure-eq-measure ennreal-1)
hence [simp]: $0 < MA$ -PS.prob B
using *assms* pA -pos
by (metis divide-eq-0-iff measure-cond-prob-space-subset zero-less-measure-iff)
have [simp]: $B \in MA$ -PS.events
using *assms* **by** (rewrite sets-cond-prob-space, unfold image-def) blast
have [simp]: finite-measure $((M \setminus A) \setminus B)$
by (rule prob-space.finite-measure, rule prob-space.cond-prob-space-correct, simp-all add: MA -PS.prob-space-axioms)
show sets-MAB: sets $((M \setminus A) \setminus B) = sets (M \setminus B)$
apply (rewrite prob-space.sets-cond-prob-space)
using MA -PS.prob-space-axioms **apply** presburger
apply (rewrite sets-cond-prob-space, unfold image-def)+
using *assms* **by** blast
show $\bigwedge C. C \in sets ((M \setminus A) \setminus B) \implies emeasure ((M \setminus A) \setminus B) C = emeasure (M \setminus B) C$
proof –
fix C **assume** $C \in sets ((M \setminus A) \setminus B)$
hence $C \in sets (M \setminus B)$ **using** sets-MAB **by** simp
from this **obtain** D **where** $D \in events$ $C = B \cap D$
by (rewrite in asm sets-cond-prob-space, auto)
hence [simp]: $C \in events$ **and** [simp]: $C \subseteq B$ **and** [simp]: $C \subseteq A$ **using** *assms*
by auto
hence [simp]: $C \in MA$ -PS.events
using *assms* **by** (rewrite sets-cond-prob-space, unfold image-def) blast
show $emeasure ((M \setminus A) \setminus B) C = emeasure (M \setminus B) C$
apply (rewrite finite-measure.emeasure-eq-measure, simp)+
apply (rewrite ennreal-inj, simp-all)
apply (rewrite prob-space.measure-cond-prob-space-subset, simp-all add: *assms* prob-space-axioms MA -PS.prob-space-axioms)+
using pA -pos **by** fastforce
qed
qed

lemma cond-prob-space-prob:

assumes[measurable]: Measurable.pred M P Measurable.pred M Q
and $\mathcal{P}(x \text{ in } M. Q x) > 0$
shows $measure (M \setminus \{x \in space M. Q x\}) \{x \in space M. P x \wedge Q x\} = \mathcal{P}(x \text{ in } M. P x \mid Q x)$
proof –
let ?SetP = $\{x \in space M. P x\}$
let ?SetQ = $\{x \in space M. Q x\}$
have $measure (M \setminus ?SetQ) \{x \in space M. P x \wedge Q x\} = measure (M \setminus ?SetQ) (?SetP \cap ?SetQ)$
by (smt (verit, ccfv-SIG) Collect-cong Int-def mem-Collect-eq)
also **have** $\dots = prob (?SetP \cap ?SetQ) / prob ?SetQ$
using *assms* **by** (rewrite measure-cond-prob-space; simp?)
also **have** $\dots = \mathcal{P}(x \text{ in } M. P x \mid Q x)$

unfolding *cond-prob-def assms* **by** (*smt (verit) Collect-cong Int-def mem-Collect-eq*)
finally show *?thesis* .
qed

lemma *cond-prob-space-cond-prob*:

assumes [*measurable*]: *Measurable.pred M P Measurable.pred M Q*
and $\mathcal{P}(x \text{ in } M. Q x) > 0$
shows $\mathcal{P}(x \text{ in } M. P x \mid Q x) = \mathcal{P}(x \text{ in } (M \downarrow \{x \in \text{space } M. Q x\}). P x)$

proof –

let *?setQ* = $\{x \in \text{space } M. Q x\}$
have $\mathcal{P}(x \text{ in } M. P x \mid Q x) = \text{measure } (M \downarrow ?setQ) \{x \in \text{space } M. P x \wedge Q x\}$
using *cond-prob-space-prob assms* **by** *simp*
also have $\dots = \mathcal{P}(x \text{ in } (M \downarrow ?setQ). P x)$

proof –

have $\{x \in \text{space } M. P x \wedge Q x\} = \{x \in \text{space } (M \downarrow ?setQ). P x\}$
using *space-cond-prob-space assms* **by** *force*
thus *?thesis* **by** *simp*

qed

finally show *?thesis* .

qed

lemma *cond-prob-neg*:

assumes[*measurable*]: *Measurable.pred M P Measurable.pred M Q*
and $\mathcal{P}(x \text{ in } M. Q x) > 0$
shows $\mathcal{P}(x \text{ in } M. \neg P x \mid Q x) = 1 - \mathcal{P}(x \text{ in } M. P x \mid Q x)$

proof –

let *?setP* = $\{x \in \text{space } M. P x\}$
let *?setQ* = $\{x \in \text{space } M. Q x\}$

interpret *setQ-PS*: *prob-space M \downarrow ?setQ* **using** *cond-prob-space-correct assms* **by**
simp

have [*simp*]: $\{x \in \text{space } (M \downarrow ?setQ). P x\} \in \text{setQ-PS.events}$

proof –

have $\{x \in \text{space } (M \downarrow ?setQ). P x\} = ?setQ \cap ?setP$ **using** *space-cond-prob-space*
by *force*

thus *?thesis* **using** *sets-cond-prob-space* **by** *simp*

qed

have $\mathcal{P}(x \text{ in } M. \neg P x \mid Q x) = \mathcal{P}(x \text{ in } M \downarrow ?setQ. \neg P x)$

by (*rewrite cond-prob-space-cond-prob; simp add: assms*)

also have $\dots = 1 - \mathcal{P}(x \text{ in } M \downarrow ?setQ. P x)$ **by** (*rewrite setQ-PS.prob-neg,*
simp-all add: assms)

also have $\dots = 1 - \mathcal{P}(x \text{ in } M. P x \mid Q x)$

by (*rewrite cond-prob-space-cond-prob; simp add: assms*)

finally show *?thesis* .

qed

lemma *random-variable-cond-prob-space*:

assumes $A \in \text{events}$ *prob A > 0*

and [*measurable*]: *random-variable borel X*

shows $X \in \text{borel-measurable } (M \downarrow A)$

```

proof (rule borel-measurableI)
  fix S :: 'b set
  assume [measurable]: open S
  show X - ' S ∩ space (M ↓ A) ∈ sets (M ↓ A)
    apply (rewrite space-cond-prob-space, simp add: assms)
    apply (rewrite sets-cond-prob-space, simp add: image-def)
    apply (rule bexI[of - X - ' S ∩ space M]; measurable)
    using sets.Int-space-eq2 Int-commute assms by auto
qed

lemma AE-cond-prob-space-iff:
  assumes A ∈ events prob A > 0
  shows (AE x in M ↓ A. P x) ↔ (AE x in M. x ∈ A → P x)
proof -
  have [simp]: 1 / emeasure M A > 0
    using assms divide-ennreal emeasure-eq-measure ennreal-1
    by (smt (verit) divide-pos-pos ennreal-eq-0-iff not-gr-zero)
  show ?thesis
    unfolding cond-prob-space-def
    apply (rewrite AE-scale-measure-iff, simp)
    by (rewrite AE-restrict-space-iff; simp add: assms)
qed

lemma integral-cond-prob-space-nn:
  assumes A ∈ events prob A > 0
    and [measurable]: random-variable borel X
    and nonneg: AE x in M. x ∈ A → 0 ≤ X x
  shows integralL (M ↓ A) X = expectation (λx. indicator A x * X x) / prob A
proof -
  have [simp]: X ∈ borel-measurable (M ↓ A)
    by (rule random-variable-cond-prob-space; (simp add: assms))
  have [simp]: AE x in (M ↓ A). 0 ≤ X x
    by (rewrite AE-cond-prob-space-iff; simp add: assms)
  have [simp]: random-variable borel (λx. indicat-real A x * X x)
    using borel-measurable-indicator assms by force
  have [simp]: AE x in M. 0 ≤ indicat-real A x * X x using nonneg by fastforce
  have integralL (M ↓ A) X = enn2real (∫+ x. ennreal (X x) ∂(M ↓ A))
    by (rewrite integral-eq-nn-integral; simp)
  also have ... = enn2real (1 / prob A * ∫+ x. ennreal (X x) ∂(restrict-space M
A))
    unfolding cond-prob-space-def
    apply (rewrite nn-integral-scale-measure, simp add: measurable-restrict-space1)
    using divide-ennreal emeasure-eq-measure ennreal-1 assms by smt
  also have ... = enn2real (1 / prob A * (∫+ x. ennreal (indicator A x * X x)
∂M))
    apply (rewrite nn-integral-restrict-space, simp add: assms)
    by (metis indicator-mult-ennreal mult.commute)
  also have ... = integralL M (λx. indicator A x * X x) / prob A
    apply (rewrite integral-eq-nn-integral; simp?)

```

by (*simp add: divide-nonneg-pos enn2real-mult*)
 finally show *?thesis* by *simp*
 qed

end

Define the complementary cumulative distribution function, also known as tail distribution. The theory presented below is a slight modification of the subsection "Properties of cdf's" in the theory *Distribution-Functions*.

3.3 Complementary Cumulative Distribution Function

definition *ccdf* :: *real measure* \Rightarrow *real* \Rightarrow *real*
 where *ccdf* *M* $\equiv \lambda x. \text{measure } M \{x <..\}$
 — complementary cumulative distribution function (tail distribution)

lemma *ccdf-def2*: *ccdf* *M* *x* = *measure* *M* $\{x <..\}$
 by (*simp add: ccdf-def*)

context *finite-borel-measure*
begin

lemma *add-cdf-ccdf*: *cdf* *M* *x* + *ccdf* *M* *x* = *measure* *M* (*space* *M*)

proof —

have $\{..x\} \cup \{x <..\} = \text{UNIV}$ by *auto*
 moreover have $\{..x\} \cap \{x <..\} = \{\}$ by *auto*
 ultimately have *emeasure* *M* $\{..x\} + \text{emeasure } M \{x <..\} = \text{emeasure } M \text{UNIV}$
 using *plus-emeasure M-is-borel atMost-borel greaterThan-borel* by *metis*
 hence *measure* *M* $\{..x\} + \text{measure } M \{x <..\} = \text{measure } M \text{UNIV}$
 using *finite-emeasure-space emeasure-eq-measure ennreal-inj*
 by (*smt (verit, ccfv-SIG) ennreal-plus measure-nonneg*)
 thus *?thesis* **unfolding** *cdf-def ccdf-def* **using** *borel-UNIV* by *simp*

qed

lemma *ccdf-cdf*: *ccdf* *M* = $(\lambda x. \text{measure } M (\text{space } M) - \text{cdf } M x)$
 by (*rule ext*) (*smt add-cdf-ccdf*)

lemma *cdf-ccdf*: *cdf* *M* = $(\lambda x. \text{measure } M (\text{space } M) - \text{ccdf } M x)$
 by (*rule ext*) (*smt add-cdf-ccdf*)

lemma *isCont-cdf-ccdf*: *isCont* (*cdf* *M*) *x* \longleftrightarrow *isCont* (*ccdf* *M*) *x*

proof

show *isCont* (*cdf* *M*) *x* \implies *isCont* (*ccdf* *M*) *x* by (*rewrite cdf-cdf*) *auto*

next

show *isCont* (*ccdf* *M*) *x* \implies *isCont* (*cdf* *M*) *x* by (*rewrite cdf-ccdf*) *auto*

qed

lemma *isCont-ccdf*: *isCont* (*ccdf* *M*) *x* \longleftrightarrow *measure* *M* $\{x\} = 0$
 using *isCont-cdf-ccdf isCont-cdf* by *simp*

lemma *continuous-cdf-ccdf*:
 shows $\text{continuous } F \text{ (cdf } M) \longleftrightarrow \text{continuous } F \text{ (ccdf } M)$
 (is ?LHS \longleftrightarrow ?RHS)
proof
 assume ?LHS
 thus ?RHS using *continuous-diff continuous-const* by (rewrite *ccdf-cdf*) blast
next
 assume ?RHS
 thus ?LHS using *continuous-diff continuous-const* by (rewrite *cdf-ccdf*) blast
qed

lemma *has-real-derivative-cdf-ccdf*:
 $(\text{cdf } M \text{ has-real-derivative } D) F \longleftrightarrow (\text{ccdf } M \text{ has-real-derivative } -D) F$
proof
 assume $(\text{cdf } M \text{ has-real-derivative } D) F$
 thus $(\text{ccdf } M \text{ has-real-derivative } -D) F$
 using *ccdf-cdf DERIV-const Deriv.field-differentiable-diff* by fastforce
next
 assume $(\text{ccdf } M \text{ has-real-derivative } -D) F$
 thus $(\text{cdf } M \text{ has-real-derivative } D) F$
 using *cdf-ccdf DERIV-const Deriv.field-differentiable-diff* by fastforce
qed

lemma *differentiable-cdf-ccdf*: $(\text{cdf } M \text{ differentiable } F) \longleftrightarrow (\text{ccdf } M \text{ differentiable } F)$
unfolding *differentiable-def*
apply (rewrite *has-real-derivative-iff[THEN sym]*)
apply (rewrite *has-real-derivative-cdf-ccdf*)
by (*metis verit-minus-simplify(4)*)

lemma *deriv-cdf-ccdf*:
 assumes *cdf M differentiable at x*
 shows $\text{deriv } (\text{cdf } M) x = - \text{deriv } (\text{ccdf } M) x$
 using *has-real-derivative-cdf-ccdf differentiable-cdf-ccdf assms*
 by (*simp add: DERIV-deriv-iff-real-differentiable DERIV-imp-deriv*)

lemma *ccdf-diff-eq2*:
 assumes $x \leq y$
 shows $\text{ccdf } M x - \text{ccdf } M y = \text{measure } M \{x <..y\}$
proof –
 have $\text{ccdf } M x - \text{ccdf } M y = \text{cdf } M y - \text{cdf } M x$ using *add-cdf-ccdf* by (*smt (verit)*)
 also have $\dots = \text{measure } M \{x <..y\}$ using *cdf-diff-eq2 assms* by *simp*
 finally show ?thesis .
qed

lemma *ccdf-nonincreasing*: $x \leq y \implies \text{ccdf } M x \geq \text{ccdf } M y$
 using *add-cdf-ccdf cdf-nondecreasing* by *smt*

```

lemma ccdf-nonneg:  $ccdf\ M\ x \geq 0$ 
  using add-cdf-ccdf cdf-bounded by smt

lemma ccdf-bounded:  $ccdf\ M\ x \leq measure\ M\ (space\ M)$ 
  using add-cdf-ccdf cdf-nonneg by smt

lemma ccdf-lim-at-top:  $(ccdf\ M \longrightarrow 0)$  at-top
proof –
  have  $((\lambda x. measure\ M\ (space\ M) - cdf\ M\ x) \longrightarrow measure\ M\ (space\ M) -$ 
measure\ M\ (space\ M)) at-top
    apply (intro tendsto-intros)
    by (rule cdf-lim-at-top)
  thus ?thesis
    by (rewrite cdf-cdf) simp
qed

lemma ccdf-lim-at-bot:  $(ccdf\ M \longrightarrow measure\ M\ (space\ M))$  at-bot
proof –
  have  $((\lambda x. measure\ M\ (space\ M) - cdf\ M\ x) \longrightarrow measure\ M\ (space\ M) - 0)$ 
at-bot
    apply (intro tendsto-intros)
    by (rule cdf-lim-at-bot)
  thus ?thesis
    by (rewrite cdf-cdf) simp
qed

lemma ccdf-is-right-cont: continuous (at-right a) (ccdf M)
proof –
  have continuous (at-right a) ( $\lambda x. measure\ M\ (space\ M) - cdf\ M\ x$ )
    apply (intro continuous-intros)
    by (rule cdf-is-right-cont)
  thus ?thesis by (rewrite cdf-cdf) simp
qed

end

context prob-space
begin

lemma ccdf-distr-measurable [measurable]:
  assumes [measurable]: random-variable borel X
  shows  $ccdf\ (distr\ M\ borel\ X) \in borel\text{-measurable}\ borel$ 
  using real-distribution.finite-borel-measure-M by (rewrite finite-borel-measure.cdf-cdf;
simp)

lemma ccdf-distr-P:
  assumes random-variable borel X
  shows  $ccdf\ (distr\ M\ borel\ X)\ x = \mathcal{P}(\omega\ in\ M. X\ \omega > x)$ 

```

unfolding *ccdf-def* **apply** (*rewrite measure-distr*; (*simp add: assms*)?)
unfolding *vimage-def* **by** (*rule arg-cong[where f=prob]*) *force*

lemma *ccdf-continuous-distr-P-ge*:

assumes *random-variable borel X isCont (ccdf (distr M borel X)) x*
shows *ccdf (distr M borel X) x = $\mathcal{P}(\omega \text{ in } M. X \omega \geq x)$*

proof –

have *ccdf (distr M borel X) x = measure (distr M borel X) {x<..}* **unfolding**
ccdf-def **by** *simp*

also have $\dots = \text{measure (distr M borel X) } (\{x<..\} \cup \{x\})$

apply (*rewrite finite-measure.measure-zero-union, simp-all add: assms finite-measure-distr*)

using *finite-borel-measure.isCont-ccdf real-distribution.finite-borel-measure-M*

assms **by** *blast*

also have $\dots = \text{measure (distr M borel X) } \{x..\}$ **by** (*metis Un-commute ivl-disj-un-singleton(1)*)

also have $\dots = \mathcal{P}(\omega \text{ in } M. X \omega \geq x)$

apply (*rewrite measure-distr, simp-all add: assms*)

unfolding *vimage-def* **by** *simp (smt (verit) Collect-cong Int-def mem-Collect-eq)*

finally show *?thesis* .

qed

lemma *ccdf-distr-diff-P*:

assumes $x \leq y$

and *random-variable borel X*

shows *ccdf (distr M borel X) x - ccdf (distr M borel X) y = $\mathcal{P}(\omega \text{ in } M. x < X \omega \wedge X \omega \leq y)$*

proof –

interpret *distrX-FBM: finite-borel-measure distr M borel X*

using *real-distribution.finite-borel-measure-M real-distribution-distr assms* **by**
simp

have *ccdf (distr M borel X) x - ccdf (distr M borel X) y = measure (distr M borel X) {x<..y}*

by (*rewrite distrX-FBM.ccdf-diff-eq2; simp add: assms*)

also have $\dots = \mathcal{P}(\omega \text{ in } M. x < X \omega \wedge X \omega \leq y)$

apply (*rewrite measure-distr; (simp add: assms)*?)

unfolding *vimage-def* **by** (*rule arg-cong[where f=prob], force*)

finally show *?thesis* .

qed

end

context *real-distribution*

begin

lemma *ccdf-bounded-prob: $\bigwedge x. \text{ccdf } M \ x \leq 1$*

by (*subst prob-space[THEN sym], rule ccdf-bounded*)

lemma *ccdf-lim-at-bot-prob: (ccdf M \longrightarrow 1) at-bot*

by (*subst prob-space[THEN sym], rule ccdf-lim-at-bot*)

end

Introduce the hazard rate. This notion will be used to define the force of mortality.

3.4 Hazard Rate

context *prob-space*

begin

definition *hazard-rate* :: (*'a* \Rightarrow *real*) \Rightarrow *real* \Rightarrow *real*

where *hazard-rate* *X* *t* \equiv

$\text{Lim } (\text{at-right } 0) (\lambda dt. \mathcal{P}(x \text{ in } M. t < X x \wedge X x \leq t + dt \mid X x > t) / dt)$

— Here, *X* is supposed to be a random variable.

lemma *hazard-rate-0-ccdf-0*:

assumes [*measurable*]: *random-variable borel X*

and *ccdf (distr M borel X) t = 0*

shows *hazard-rate X t = 0*

— Note that division by 0 is calculated as 0 in Isabelle/HOL.

proof —

have $\bigwedge dt. \mathcal{P}(x \text{ in } M. t < X x \wedge X x \leq t + dt \mid X x > t) = 0$

unfolding *cond-prob-def* **using** *ccdf-distr-P* **assms** **by** *simp*

hence *hazard-rate X t = Lim (at-right 0) ($\lambda dt::real. 0$)*

unfolding *hazard-rate-def* **by** (*rewrite Lim-cong; simp*)

also have $\dots = 0$ **by** (*rule tendsto-Lim; simp*)

finally show *?thesis* .

qed

lemma *hazard-rate-deriv-cdf*:

assumes [*measurable*]: *random-variable borel X*

and (*cdf (distr M borel X)*) *differentiable at t*

shows *hazard-rate X t = deriv (cdf (distr M borel X)) t / cdf (distr M borel X) t*

proof (*cases* $\langle \text{cdf } (\text{distr } M \text{ borel } X) \text{ } t = 0 \rangle$)

case *True*

with *hazard-rate-0-ccdf-0* **show** *?thesis* **by** *simp*

next

case *False*

let *?F = cdf (distr M borel X)*

have $\forall_F dt \text{ in } \text{at-right } 0. \mathcal{P}(x \text{ in } M. t < X x \wedge X x \leq t + dt \mid X x > t) / dt =$
 $(?F (t + dt) - ?F t) / dt / \text{cdf } (\text{distr } M \text{ borel } X) \text{ } t$

apply (*rule eventually-at-rightI*[**where** *b=1*]; *simp*)

unfolding *cond-prob-def*

apply (*rewrite cdf-distr-diff-P; simp*)

apply (*rewrite cdf-distr-P*[*THEN sym*], *simp*)

by (*smt (verit) Collect-cong mult.commute*)

moreover have $(\lambda dt. (?F (t + dt) - ?F t) / dt / \text{cdf } (\text{distr } M \text{ borel } X) \text{ } t)$

\longrightarrow

$\text{deriv } ?F t / \text{ccdf } (\text{distr } M \text{ borel } X) t) (\text{at-right } 0)$
apply (rule *tendsto-intros*, *simp-all add: False*)
apply (rule *Lim-at-imp-Lim-at-within*)
using *DERIV-deriv-iff-real-differentiable* *assms* *DERIV-def* **by** *blast*
ultimately show *?thesis*
unfolding *hazard-rate-def* **using** *tendsto-cong* **by** (*intro tendsto-Lim; force*)
qed

lemma *deriv-cdf-hazard-rate*:
assumes [*measurable*]: *random-variable borel X*
and (*cdf (distr M borel X)*) *differentiable at t*
shows $\text{deriv } (\text{cdf } (\text{distr } M \text{ borel } X)) t = \text{ccdf } (\text{distr } M \text{ borel } X) t * \text{hazard-rate } X t$
proof –
interpret *distrX-FBM: finite-borel-measure distr M borel X*
using *real-distribution.finite-borel-measure-M* *real-distribution-distr* **assms** **by** *simp*
show *?thesis*
proof (*cases* $\langle \text{ccdf } (\text{distr } M \text{ borel } X) t = 0 \rangle$)
case *True*
hence $\text{cdf } (\text{distr } M \text{ borel } X) t = 1$
using *distrX-FBM.cdf-ccdf*
by *simp (metis assms(1) distrX-FBM.borel-UNIV prob-space.prob-space prob-space-distr)*
moreover obtain *D* **where** (*cdf (distr M borel X) has-real-derivative D*) (*at t*)
using *assms real-differentiable-def* **by** *atomize-elim blast*
ultimately have (*cdf (distr M borel X) has-real-derivative 0*) (*at t*)
using *assms*
by (*smt (verit) DERIV-local-max real-distribution.cdf-bounded-prob real-distribution-distr*)
thus *?thesis* **using** *True* **by** (*simp add: DERIV-imp-deriv*)
next
case *False*
thus *?thesis* **using** *hazard-rate-deriv-cdf* **assms** **by** *simp*
qed
qed

lemma *hazard-rate-deriv-ccdf*:
assumes [*measurable*]: *random-variable borel X*
and (*ccdf (distr M borel X)*) *differentiable at t*
shows $\text{hazard-rate } X t = - \text{deriv } (\text{ccdf } (\text{distr } M \text{ borel } X)) t / \text{ccdf } (\text{distr } M \text{ borel } X) t$
proof –
interpret *distrX-FBM: finite-borel-measure distr M borel X*
using *real-distribution.finite-borel-measure-M* *real-distribution-distr* **assms** **by** *simp*
show *?thesis*
using *hazard-rate-deriv-cdf distrX-FBM.deriv-cdf-ccdf* *assms* *distrX-FBM.differentiable-cdf-ccdf*
by *presburger*

qed

lemma *deriv-ccdf-hazard-rate*:

assumes [*measurable*]: *random-variable borel X*

and (*ccdf (distr M borel X)*) *differentiable at t*

shows *deriv (ccdf (distr M borel X)) t = - ccdf (distr M borel X) t * hazard-rate X t*

proof –

interpret *distrX-FBM: finite-borel-measure distr M borel X*

using *real-distribution.finite-borel-measure-M real-distribution-distr assms* **by**

simp

show *?thesis*

using *deriv-ccdf-hazard-rate distrX-FBM.deriv-ccdf-ccdf assms distrX-FBM.differentiable-ccdf-ccdf*

by *simp*

qed

lemma *hazard-rate-deriv-ln-ccdf*:

assumes [*measurable*]: *random-variable borel X*

and (*ccdf (distr M borel X)*) *differentiable at t*

and *ccdf (distr M borel X) t ≠ 0*

shows *hazard-rate X t = - deriv (λt. ln (ccdf (distr M borel X) t)) t*

proof –

interpret *distrX-FBM: finite-borel-measure distr M borel X*

using *real-distribution.finite-borel-measure-M real-distribution-distr assms* **by**

simp

let *?srul = ccdf (distr M borel X)*

have *?srul t > 0* **using** *distrX-FBM.ccdf-nonneg assms* **by** (*smt (verit)*)

moreover **have** (*?srul has-real-derivative (deriv ?srul t)*) (*at t*)

using *DERIV-deriv-iff-real-differentiable assms* **by** *blast*

ultimately **have** (*(λt. ln (?srul t)) has-real-derivative 1 / ?srul t * deriv ?srul t*) (*at t*)

by (*rule derivative-intros*)

hence *deriv (λt. ln (?srul t)) t = deriv ?srul t / ?srul t* **by** (*simp add: DERIV-imp-deriv*)

also **have** *... = - hazard-rate X t* **using** *hazard-rate-deriv-ccdf assms* **by** *simp*

finally **show** *?thesis* **by** *simp*

qed

lemma *hazard-rate-has-integral*:

assumes [*measurable*]: *random-variable borel X*

and *t ≤ u*

and (*ccdf (distr M borel X)*) *piecewise-differentiable-on {t<..}*

and *continuous-on {t..u}* (*ccdf (distr M borel X)*)

and $\bigwedge s. s \in \{t..u\} \implies \text{ccdf (distr M borel X) } s \neq 0$

shows

(hazard-rate X has-integral ln (ccdf (distr M borel X) t / ccdf (distr M borel X) u)) {t..u}

proof –

interpret *distrX-FBM: finite-borel-measure distr M borel X*

```

using real-distribution.finite-borel-measure-M real-distribution-distr assms by
simp
let ?srval = cdf (distr M borel X)
have [simp]:  $\bigwedge s. t \leq s \wedge s \leq u \implies ?srval\ s > 0$ 
using distrX-FBM.cdf-nonneg assms by (smt (verit) atLeastAtMost-iff)
have (deriv ( $\lambda s. - \ln (?srval\ s)$ ) has-integral  $-\ln (?srval\ u) - - \ln (?srval\ t)$ )
{t..u}
proof -
have continuous-on {t..u} ( $\lambda s. - \ln (?srval\ s)$ )
by (rule continuous-intros, rule continuous-on-ln, auto simp add: assms)
moreover hence ( $\lambda s. - \ln (?srval\ s)$ ) piecewise-differentiable-on {t<..}
proof -
have ?srval ' {t<..}  $\subseteq$  {0<..}
proof -
{ fix s assume s  $\in$  {t<..}
hence ?srval s  $\neq$  0 using assms by simp
moreover have ?srval s  $\geq$  0 using distrX-FBM.cdf-nonneg by simp
ultimately have ?srval s  $>$  0 by simp }
thus ?thesis by auto
qed
hence ( $\lambda r. - \ln r$ )  $\circ$  ?srval piecewise-differentiable-on {t<..}
apply (intro differentiable-on-piecewise-compose, simp add: assms)
apply (rule derivative-intros)
apply (rule differentiable-on-subset[of  $\ln$  {0<..}], simp-all)
apply (rewrite differentiable-on-eq-field-differentiable-real, auto)
unfolding field-differentiable-def using DERIV-ln by (metis has-field-derivative-at-within)
thus ?thesis unfolding comp-def by simp
qed
ultimately show ?thesis by (intro FTC-real-deriv-has-integral; simp add:
assms)
qed
hence ln: (deriv ( $\lambda s. - \ln (?srval\ s)$ ) has-integral  $\ln (?srval\ t / ?srval\ u)$ ) {t..u}
by simp (rewrite ln-div; force simp: assms)
thus ((hazard-rate X) has-integral  $\ln (?srval\ t / ?srval\ u)$ ) {t..u}
proof -
from assms obtain S0 where finS0: finite S0 and
diffS0:  $\bigwedge s. s \in \{t<..\} - S0 \implies ?srval$  differentiable at s within {t<..}
unfolding piecewise-differentiable-on-def by blast
from this obtain S where finite S and  $\bigwedge s. s \in \{t..u\} - S \implies ?srval$  differ-
entiable at s
proof (atomize-elim)
let ?S = S0  $\cup$  {t, u}
have finite ?S using finS0 by simp
moreover have  $\forall s. s \in \{t..u\} - ?S \implies$  cdf (distr M borel X) differentiable
at s
proof -
{ fix s assume s-in: s  $\in$  {t..u} - ?S
hence ?srval differentiable at s within {t<..} using diffS0 by simp
hence ?srval differentiable at s

```

```

    using s-in by (rewrite at-within-open[THEN sym], simp-all) }
  thus ?thesis by blast
qed
ultimately show
   $\exists S. \text{finite } S \wedge (\forall s. s \in \{t..u\} - S \longrightarrow \text{ccdf } (\text{distr } M \text{ borel } X) \text{ differentiable at } s)$ 
  by blast
qed
thus ?thesis
  apply (rewrite has-integral-spike-finite-eq [of S - deriv ( $\lambda s. - \ln$  (?srvl s))], simp)
  apply (rewrite hazard-rate-deriv-ln-ccdf, simp-all add: assms)
  apply (rewrite deriv-minus, simp-all)
  apply (rewrite in asm differentiable-eq-field-differentiable-real)
  apply (rewrite comp-def[THEN sym], rule field-differentiable-compose[of ?srvl], simp-all)
  unfolding field-differentiable-def apply (rule exI, rule DERIV-ln, simp)
  using ln by simp
qed
qed

```

corollary *hazard-rate-integrable*:

```

  assumes [measurable]: random-variable borel X
    and  $t \leq u$ 
    and (ccdf (distr M borel X)) piecewise-differentiable-on  $\{t < .. < u\}$ 
    and continuous-on  $\{t..u\}$  (ccdf (distr M borel X))
    and  $\bigwedge s. s \in \{t..u\} \implies \text{ccdf } (\text{distr } M \text{ borel } X) s \neq 0$ 
  shows hazard-rate X integrable-on  $\{t..u\}$ 
  using has-integral-integrable hazard-rate-has-integral assms by blast

```

lemma *ccdf-exp-cumulative-hazard*:

```

  assumes [measurable]: random-variable borel X
    and  $t \leq u$ 
    and (ccdf (distr M borel X)) piecewise-differentiable-on  $\{t < .. < u\}$ 
    and continuous-on  $\{t..u\}$  (ccdf (distr M borel X))
    and  $\bigwedge s. s \in \{t..u\} \implies \text{ccdf } (\text{distr } M \text{ borel } X) s \neq 0$ 
  shows ccdf (distr M borel X) u / ccdf (distr M borel X) t =
    exp (- integral {t..u} (hazard-rate X))
  proof -
    interpret distrX-FBM: finite-borel-measure distr M borel X
      using real-distribution.finite-borel-measure-M real-distribution-distr assms by simp
    let ?srvl = ccdf (distr M borel X)
    have [simp]:  $\bigwedge s. t \leq s \wedge s \leq u \implies ?srvl s > 0$ 
      using distrX-FBM.ccdf-nonneg assms by (smt (verit) atLeastAtMost-iff)
    have integral {t..u} (hazard-rate X) =  $\ln$  (?srvl t / ?srvl u)
      using hazard-rate-has-integral has-integral-integrable-integral assms by auto
    also have  $\dots = - \ln$  (?srvl u / ?srvl t) using ln-div assms by simp
    finally have - integral {t..u} (hazard-rate X) =  $\ln$  (?srvl u / ?srvl t) by simp
  qed

```

```

thus ?thesis using assms by simp
qed

lemma hazard-rate-density-ccdf:
  assumes distributed M lborel X f
  and  $\bigwedge s. f s \geq 0 \ t < u$  continuous-on {t..u} f
  shows hazard-rate X t = f t / cdf (distr M borel X) t
proof (cases  $\langle \text{ccdf (distr M borel X) t} = 0 \rangle$ )
  case True
  thus ?thesis
  apply (rewrite hazard-rate-0-ccdf-0, simp-all)
  using assms(1) distributed-measurable by force
next
  case False
  have [simp]:  $t \leq u$  using assms by simp
  have [measurable]: random-variable borel X
  using assms distributed-measurable measurable-lborel1 by blast
  have [measurable]:  $f \in \text{borel-measurable lborel}$ 
  using assms distributed-real-measurable by metis
  have [simp]: integrable lborel f
  proof –
  have  $\text{prob } (X - \text{'UNIV } \cap \text{space } M) = \text{LINT } x | \text{lborel. indicat-real UNIV } x * f x$ 
  by (rule distributed-measure; simp add: assms)
  thus ?thesis
  using prob-space not-integrable-integral-eq by fastforce
qed
have  $((\lambda dt. (\text{LBINT } s:\{t..t+dt\}. f s) / dt) \longrightarrow f t)$  (at-right 0)
proof –
  have  $\bigwedge dt. (\int^+ x. \text{ennreal (indicat-real } \{t..t+dt\} x * f x) \partial \text{lborel}) < \infty$ 
  proof –
  fix dt :: real
  have  $(\int^+ x. \text{ennreal (indicat-real } \{t..t+dt\} x * f x) \partial \text{lborel}) =$ 
     $\text{set-nn-integral lborel } \{t..t+dt\} f$ 
  by (metis indicator-mult-ennreal mult commute)
  moreover have  $\text{emeasure } M (X - \text{' } \{t..t+dt\} \cap \text{space } M) = \text{set-nn-integral}$ 
     $\text{lborel } \{t..t+dt\} f$ 
  by (rule distributed-emeasure; simp add: assms)
  moreover have  $\text{emeasure } M (X - \text{' } \{t..t+dt\} \cap \text{space } M) < \infty$ 
  using emeasure-eq-measure ennreal-less-top infinity-ennreal-def by presburger
  ultimately show  $(\int^+ x. \text{ennreal (indicat-real } \{t..t+dt\} x * f x) \partial \text{lborel}) <$ 
 $\infty$  by simp
qed
hence  $\bigwedge dt. (\text{LBINT } s:\{t..t+dt\}. f s) = \text{integral } \{t..t+dt\} f$ 
  apply (intro set-borel-integral-eq-integral)
  unfolding set-integrable-def
  apply (rule integrableI-nonneg; simp)
  by (rule AE-I2, simp add: assms)
  moreover have  $((\lambda dt. (\text{integral } \{t..t+dt\} f) / dt) \longrightarrow f t)$  (at-right 0)
  proof –

```

```

have (( $\lambda x$ .  $\text{integral } \{t..x\} f$ )  $\text{has-real-derivative } f t$ ) (at  $t$  within  $\{t..u\}$ )
  by (rule  $\text{integral-has-real-derivative}$ ; simp add:  $\text{assms}$ )
moreover have (at  $t$  within  $\{t..u\}$ ) = (at  $(t+0)$  within  $(+ )t \text{ ' } \{0..u-t\}$ ) by
simp
  ultimately have
    (( $\lambda x$ .  $\text{integral } \{t..x\} f$ )  $\circ (+ )t$   $\text{has-real-derivative } f t$ ) (at  $0$  within  $\{0..u-t\}$ )
    by (metis  $\text{DERIV-at-within-shift-lemma}$ )
  hence (( $\lambda dt$ . ( $\text{integral } \{t..t+dt\} f$ ) /  $dt$ )  $\longrightarrow f t$ ) (at  $0$  within  $\{0..u-t\}$ )
    using  $\text{has-field-derivative-iff}$  by force
  thus ?thesis using  $\text{at-within-Icc-at-right assms}$  by smt
qed
ultimately show ?thesis by simp
qed
moreover have  $\bigwedge dt$ .  $dt > 0 \implies \mathcal{P}(x \text{ in } M. X x \in \{t <.. t+dt\}) = (\text{LBINT } s:\{t..t+dt\}. f s)$ 
proof -
  fix  $dt :: \text{real}$  assume  $dt > 0$ 
  hence [simp]:  $\text{sym-diff } \{t <.. t + dt\} \{t..t + dt\} = \{t\}$  by force
  have  $\text{prob } (X \text{ - ' } \{t <.. t+dt\} \cap \text{space } M) = \int s. \text{indicator } \{t <.. t+dt\} s * f s$ 
 $\partial \text{lborel}$ 
  by (rule  $\text{distributed-measure}$ ; simp add:  $\text{assms}$ )
  hence  $\mathcal{P}(x \text{ in } M. X x \in \{t <.. t+dt\}) = (\text{LBINT } s:\{t <.. t+dt\}. f s)$ 
  unfolding  $\text{set-lebesgue-integral-def vimage-def Int-def}$  by simp (smt (verit)
Collect-cong)
  moreover have  $(\text{LBINT } s:\{t <.. t+dt\}. f s) = (\text{LBINT } s:\{t..t+dt\}. f s)$ 
  by (rule  $\text{set-integral-null-delta}$ ; force)
  ultimately show  $\mathcal{P}(x \text{ in } M. X x \in \{t <.. t+dt\}) = (\text{LBINT } s:\{t..t+dt\}. f s)$ 
by simp
qed
ultimately have (( $\lambda dt$ .  $\mathcal{P}(x \text{ in } M. t < X x \wedge X x \leq t + dt)$  /  $dt$ )  $\longrightarrow f t$ )
(at-right  $0$ )
  by simp (smt (verit)  $\text{Lim-cong-within greaterThan-iff}$ )
  hence (( $\lambda dt$ .  $\mathcal{P}(x \text{ in } M. t < X x \wedge X x \leq t + dt \mid X x > t)$  /  $dt$ )  $\longrightarrow$ 
 $f t$  /  $\text{cdf } (\text{distr } M \text{ borel } X) t$ ) (at-right  $0$ )
  unfolding  $\text{cond-prob-def}$ 
  apply (rewrite  $\text{cdf-distr-P[THEN sym]}$ ; simp)
  apply (rewrite  $\text{mult.commute}$ , rewrite  $\text{divide-divide-eq-left[THEN sym]}$ )
  by (rule  $\text{tendsto-intros}$ ; (simp add:  $\text{False}$ )?) (smt (verit)  $\text{Collect-cong Lim-cong-within}$ )
thus ?thesis unfolding  $\text{hazard-rate-def}$  by (intro  $\text{tendsto-Lim}$ ; simp)
qed

end

end
theory Interest
  imports Preliminaries
begin

```

4 Theory of Interest

locale *interest* =

fixes $i :: \text{real}$ — i stands for an interest rate.

assumes *v-futr-pos*: $1 + i > 0$ — Assume that the future value is positive.

begin

definition *i-nom* :: $\text{nat} \Rightarrow \text{real}$ ($\$i\{-\}$ [0] 200)

where $\$i\{m\} \equiv m * ((1+i).\wedge(1/m) - 1)$ — nominal interest rate

definition *i-force* :: real ($\$\delta$ 200)

where $\$\delta \equiv \ln(1+i)$ — force of interest

definition *d-nom* :: $\text{nat} \Rightarrow \text{real}$ ($\$d\{-\}$ [0] 200)

where $\$d\{m\} \equiv \$i\{m\} / (1 + \$i\{m\}/m)$ — discount rate

abbreviation *d-nom-yr* :: real ($\$d$ 200)

where $\$d \equiv \$d\{1\}$ — Post-fix *yr* stands for "year".

definition *v-pres* :: real ($\$v$ 200)

where $\$v \equiv 1 / (1+i)$ — present value factor

definition *ann* :: $\text{nat} \Rightarrow \text{nat} \Rightarrow \text{real}$ ($\$a\{-\}'\text{-}$ [0,101] 200)

where $\$a\{m\}\text{-}n \equiv \sum_{k < n*m} \$v.\wedge((k+1)::\text{nat})/m) / m$
— present value of an immediate annuity

abbreviation *ann-yr* :: $\text{nat} \Rightarrow \text{real}$ ($\$a'\text{-}$ [101] 200)

where $\$a\text{-}n \equiv \$a\{1\}\text{-}n$

definition *acc* :: $\text{nat} \Rightarrow \text{nat} \Rightarrow \text{real}$ ($\$\$s\{-\}'\text{-}$ [0,101] 200)

where $\$\$s\{m\}\text{-}n \equiv \sum_{k < n*m} (1+i).\wedge((k::\text{nat})/m) / m$
— future value of an immediate annuity
— The name *acc* stands for "accumulation".

abbreviation *acc-yr* :: $\text{nat} \Rightarrow \text{real}$ ($\$\$s'\text{-}$ 200)

where $\$\$s\text{-}n \equiv \$\$s\{1\}\text{-}n$

definition *ann-due* :: $\text{nat} \Rightarrow \text{nat} \Rightarrow \text{real}$ ($\$a'''\{-\}'\text{-}$ [0,101] 200)

where $\$a'''\{m\}\text{-}n \equiv \sum_{k < n*m} \$v.\wedge((k::\text{nat})/m) / m$
— present value of an annuity-due

abbreviation *ann-due-yr* :: $\text{nat} \Rightarrow \text{real}$ ($\$a'''\text{-}$ [101] 200)

where $\$a'''\text{-}n \equiv \$a'''\{1\}\text{-}n$

definition *acc-due* :: $\text{nat} \Rightarrow \text{nat} \Rightarrow \text{real}$ ($\$\$s'''\{-\}'\text{-}$ [0,101] 200)

where $\$\$s'''\{m\}\text{-}n \equiv \sum_{k < n*m} (1+i).\wedge((k+1)::\text{nat})/m) / m$
— future value of an annuity-due

abbreviation *acc-due-yr* :: $\text{nat} \Rightarrow \text{real}$ ($\$\$s'''\text{-}$ [101] 200)

where $\$s''\text{-}n \equiv \$s''\{1\}\text{-}n$

definition $\text{ann-cont} :: \text{real} \Rightarrow \text{real}$ ($\$a'''\text{-}$ [101] 200)
where $\$a'\text{-}n \equiv \text{integral } \{0..n\} (\lambda t::\text{real}. \$v.\hat{t})$
— present value of a continuous annuity

definition $\text{acc-cont} :: \text{real} \Rightarrow \text{real}$ ($\$s'''\text{-}$ [101] 200)
where $\$s'\text{-}n \equiv \text{integral } \{0..n\} (\lambda t::\text{real}. (1+i).\hat{t})$
— future value of a continuous annuity

definition $\text{perp} :: \text{nat} \Rightarrow \text{real}$ ($\$a\hat{\{-}\}\text{'-}\infty$ [0] 200)
where $\$a\hat{\{m\}}\text{'-}\infty \equiv 1 / \$i\hat{\{m\}}$
— present value of a perpetual annuity

abbreviation $\text{perp-yr} :: \text{real}$ ($\$a'\text{-}\infty$ 200)
where $\$a\text{-}\infty \equiv \$a\hat{\{1\}}\text{'-}\infty$

definition $\text{perp-due} :: \text{nat} \Rightarrow \text{real}$ ($\$a'''\hat{\{-}\}\text{'-}\infty$ [0] 200)
where $\$a''\hat{\{m\}}\text{'-}\infty \equiv 1 / \$d\hat{\{m\}}$
— present value of a perpetual annuity-due

abbreviation $\text{perp-due-yr} :: \text{real}$ ($\$a'''\text{-}\infty$ 200)
where $\$a''\text{-}\infty \equiv \$a''\hat{\{1\}}\text{'-}\infty$

definition $\text{ann-incr} :: \text{nat} \Rightarrow \text{nat} \Rightarrow \text{nat} \Rightarrow \text{real}$ ($\$(I\hat{\{-}\}a)\hat{\{-}\}\text{'-}$ [0,0,101] 200)
where $\$(I\hat{\{l\}}a)\hat{\{m\}}\text{'-}n \equiv \sum k < n * m. \$v.\hat{(k+1::nat)/m)} * \lceil l * (k+1::nat) / m \rceil / (l * m)$
— present value of an increasing annuity
— This is my original definition.
— Here, l represents the number of increments per unit time.

abbreviation $\text{ann-incr-lvl} :: \text{nat} \Rightarrow \text{nat} \Rightarrow \text{real}$ ($\$(Ia)\hat{\{-}\}\text{'-}$ [0,101] 200)
where $\$(Ia)\hat{\{m\}}\text{'-}n \equiv \$(I\hat{\{1\}}a)\hat{\{m\}}\text{'-}n$
— The post-fix lvl stands for "level".

abbreviation $\text{ann-incr-yr} :: \text{nat} \Rightarrow \text{real}$ ($\$(Ia)\text{'-}$ [101] 200)
where $\$(Ia)\text{'-}n \equiv \$(Ia)\hat{\{1\}}\text{'-}n$

definition $\text{acc-incr} :: \text{nat} \Rightarrow \text{nat} \Rightarrow \text{nat} \Rightarrow \text{real}$ ($\$(I\hat{\{-}\}s)\hat{\{-}\}\text{'-}$ [0,0,101] 200)
where $\$(I\hat{\{l\}}s)\hat{\{m\}}\text{'-}n \equiv \sum k < n * m. (1+i).\hat{(n-(k+1::nat)/m)} * \lceil l * (k+1::nat) / m \rceil / (l * m)$
— future value of an increasing annuity

abbreviation $\text{acc-incr-lvl} :: \text{nat} \Rightarrow \text{nat} \Rightarrow \text{real}$ ($\$(Is)\hat{\{-}\}\text{'-}$ [0,101] 200)
where $\$(Is)\hat{\{m\}}\text{'-}n \equiv \$(I\hat{\{1\}}s)\hat{\{m\}}\text{'-}n$

abbreviation $\text{acc-incr-yr} :: \text{nat} \Rightarrow \text{real}$ ($\$(Is)\text{'-}$ [101] 200)
where $\$(Is)\text{'-}n \equiv \$(Is)\hat{\{1\}}\text{'-}n$

definition $ann\text{-}due\text{-}incr :: nat \Rightarrow nat \Rightarrow nat \Rightarrow real (\$(I\{-}a''''')\{-}\}'-- [0,0,101] 200)$

where $\$(I\{l\}a'')\{-}m\}\text{-}n \equiv \sum k < n * m. \$v. \wedge((k::nat)/m) * [l*(k+1::nat)/m] / (l*m)$

abbreviation $ann\text{-}due\text{-}incr\text{-}lvl :: nat \Rightarrow nat \Rightarrow real (\$(Ia''''')\{-}\}'-- [0,101] 200)$

where $\$(Ia'')\{-}m\}\text{-}n \equiv \$(I\{1\}a'')\{-}m\}\text{-}n$

abbreviation $ann\text{-}due\text{-}incr\text{-}yr :: nat \Rightarrow real (\$(Ia''''')'\text{-} [101] 200)$

where $\$(Ia'')\text{-}n \equiv \$(Ia'')\{-}1\}\text{-}n$

definition $acc\text{-}due\text{-}incr :: nat \Rightarrow nat \Rightarrow nat \Rightarrow real (\$(I\{-}s''''')\{-}\}'-- [0,0,101] 200)$

where $\$(I\{l\}s'')\{-}m\}\text{-}n \equiv \sum k < n * m. (1+i). \wedge(n-(k::nat)/m) * [l*(k+1::nat)/m] / (l*m)$

abbreviation $acc\text{-}due\text{-}incr\text{-}lvl :: nat \Rightarrow nat \Rightarrow real (\$(Is''''')\{-}\}'-- [0,101] 200)$

where $\$(Is'')\{-}m\}\text{-}n \equiv \$(I\{1\}s'')\{-}m\}\text{-}n$

abbreviation $acc\text{-}due\text{-}incr\text{-}yr :: nat \Rightarrow real (\$(Is''''')'\text{-} [101] 200)$

where $\$(Is'')\text{-}n \equiv \$(Is'')\{-}1\}\text{-}n$

definition $perp\text{-}incr :: nat \Rightarrow nat \Rightarrow real (\$(I\{-}a'')\{-}\}'\text{-}\infty [0,0] 200)$

where $\$(I\{l\}a'')\{-}m\}\text{-}\infty \equiv \lim (\lambda n. \$(I\{l\}a'')\{-}m\}\text{-}n)$

abbreviation $perp\text{-}incr\text{-}lvl :: nat \Rightarrow real (\$(Ia'')\{-}\}'\text{-}\infty [0] 200)$

where $\$(Ia'')\{-}m\}\text{-}\infty \equiv \$(I\{1\}a'')\{-}m\}\text{-}\infty$

abbreviation $perp\text{-}incr\text{-}yr :: real (\$(Ia'')'\text{-}\infty 200)$

where $\$(Ia'')\text{-}\infty \equiv \$(Ia'')\{-}1\}\text{-}\infty$

definition $perp\text{-}due\text{-}incr :: nat \Rightarrow nat \Rightarrow real (\$(I\{-}a''''')\{-}\}'\text{-}\infty [0,0] 200)$

where $\$(I\{l\}a'')\{-}m\}\text{-}\infty \equiv \lim (\lambda n. \$(I\{l\}a'')\{-}m\}\text{-}n)$

abbreviation $perp\text{-}due\text{-}incr\text{-}lvl :: nat \Rightarrow real (\$(Ia''''')\{-}\}'\text{-}\infty [0] 200)$

where $\$(Ia'')\{-}m\}\text{-}\infty \equiv \$(I\{1\}a'')\{-}m\}\text{-}\infty$

abbreviation $perp\text{-}due\text{-}incr\text{-}yr :: real (\$(Ia''''')'\text{-}\infty 200)$

where $\$(Ia'')\text{-}\infty \equiv \$(Ia'')\{-}1\}\text{-}\infty$

lemma $v\text{-}futr\text{-}m\text{-}pos: 1 + \$i\{-}m\}/m > 0$ **if** $m \neq 0$ **for** $m::nat$

using $v\text{-}futr\text{-}pos$ $i\text{-}nom\text{-}def$ **by force**

lemma $i\text{-}nom\text{-}1[simp]: \$i\{-}1\} = i$

using $v\text{-}futr\text{-}pos$ $i\text{-}nom\text{-}def$ **by force**

lemma $i\text{-}nom\text{-}eff: (1 + \$i\{-}m\}/m)\wedge m = 1 + i$ **if** $m \neq 0$ **for** $m::nat$

unfolding $i\text{-}nom\text{-}def$ **using** $less\text{-}imp\text{-}neq$ $v\text{-}futr\text{-}pos$ **that**

apply ($simp$, $subst$ $powr\text{-}realpow[THEN sym]$, $simp$)

by (subst powr-powr, simp)

lemma *i-nom-i*: $1 + i^{\wedge}m/m = (1+i)^{\wedge}(1/m)$ if $m \neq 0$ for $m::nat$
unfolding *i-nom-def* by (simp add: that)

lemma *i-nom-0-iff-i-0*: $i^{\wedge}m = 0 \iff i = 0$ if $m \neq 0$ for $m::nat$
proof

assume $i^{\wedge}m = 0$

hence $\star: (1+i)^{\wedge}(1/m) = (1+i)^{\wedge}0$

unfolding *i-nom-def* using *v-futr-pos* that by simp

show $i = 0$

proof (rule ccontr)

assume $i \neq 0$

hence $1/m = 0$ using *powr-inj* \star *v-futr-pos* by smt

thus *False* using that by simp

qed

next

assume $i = 0$

thus $i^{\wedge}m = 0$

unfolding *i-nom-def* by simp

qed

lemma *i-nom-pos-iff-i-pos*: $i^{\wedge}m > 0 \iff i > 0$ if $m \neq 0$ for $m::nat$
proof

assume $i^{\wedge}m > 0$

hence $\star: (1+i)^{\wedge}(1/m) > 1^{\wedge}(1/m)$

unfolding *i-nom-def* using *v-futr-pos* that by (simp add: zero-less-mult-iff)

thus $i > 0$

using *powr-less-cancel2*[of $1/m$ 1 $1+i$] *v-futr-pos* that by simp

next

assume $i > 0$

hence $(1+i)^{\wedge}(1/m) > 1^{\wedge}(1/m)$

using *powr-less-mono2* *v-futr-pos* that by simp

thus $i^{\wedge}m > 0$

unfolding *i-nom-def* using that by (simp add: zero-less-mult-iff)

qed

lemma *e-delta*: $\exp(\delta) = 1 + i$

unfolding *i-force-def* by (simp add: *v-futr-pos*)

lemma *delta-0-iff-i-0*: $\delta = 0 \iff i = 0$

proof

assume $\delta = 0$

thus $i = 0$

using *e-delta* by auto

next

assume $i = 0$

thus $\delta = 0$

unfolding *i-force-def* by simp

qed

lemma *lim-i-nom*: $(\lambda m. \$i\{m\}) \longrightarrow \δ

proof –

let $?f = \lambda h. ((1+i).\hat{\sim}h - 1) / h$

have *D1ipwr*: *DERIV* $(\lambda h. (1+i).\hat{\sim}h) 0 := \δ

unfolding *i-force-def*

using *has-real-derivative-powr2*[*OF v-futr-pos*, **where** $x=0$] *v-futr-pos* **by** *simp*

hence *limf*: $(?f \longrightarrow \$\delta)$ (*at 0*)

unfolding *DERIV-def* using *v-futr-pos* **by** *auto*

hence $(\lambda m. \$i\{Suc\ m\}) \longrightarrow \δ

unfolding *i-nom-def* using *tendsto-at-iff-sequentially*[*of ?f \\$\delta 0 R*, *THEN iffD1*]

apply *simp*

apply (*drule-tac* $x=\lambda m. 1 / Suc\ m$ **in** *spec*, *simp*, *drule mp*)

subgoal using *lim-1-over-n LIMSEQ-Suc* **by** *force*

by (*simp add: o-def mult.commute*)

thus *?thesis*

by (*simp add: LIMSEQ-imp-Suc*)

qed

lemma *d-nom-0-iff-i-0*: $\$d\{m\} = 0 \longleftrightarrow i = 0$ **if** $m \neq 0$ **for** $m::nat$

proof –

have $\$d\{m\} = 0 \longleftrightarrow \$i\{m\} = 0$

unfolding *d-nom-def* using *v-futr-m-pos* **by** (*smt (verit) divide-eq-0-iff of-nat-0*)

thus *?thesis*

using *i-nom-0-iff-i-0 that* **by** *auto*

qed

lemma *d-nom-pos-iff-i-pos*: $\$d\{m\} > 0 \longleftrightarrow i > 0$ **if** $m \neq 0$ **for** $m::nat$

proof –

have $\$d\{m\} > 0 \longleftrightarrow \$i\{m\} > 0$

unfolding *d-nom-def* using *zero-less-divide-iff i-nom-pos-iff-i-pos v-futr-m-pos that* **by** *smt*

thus *?thesis*

using *i-nom-pos-iff-i-pos that* **by** *auto*

qed

lemma *d-nom-i-nom*: $1 - \$d\{m\}/m = 1 / (1 + \$i\{m\}/m)$ **if** $m \neq 0$ **for** $m::nat$

proof –

have $1 - \$d\{m\}/m = 1 - (\$i\{m\}/m) / (1 + \$i\{m\}/m)$

by (*simp add: d-nom-def*)

also have $\dots = 1 / (1 + \$i\{m\}/m)$

using *v-futr-m-pos*

by (*smt (verit, ccfv-SIG) add-divide-distrib that div-self*)

finally show *?thesis* .

qed

lemma *lim-d-nom*: $(\lambda m. \$d\{m\}) \longrightarrow \δ

proof –
have $(\lambda m. \$i^{\wedge}\{m\}/m) \longrightarrow 0$
using *lim-i-nom tendsto-divide-0 tendsto-of-nat* **by** *blast*
hence $(\lambda m. 1 + \$i^{\wedge}\{m\}/m) \longrightarrow 1$
by (*metis add.right-neutral tendsto-add-const-iff*)
thus *?thesis*
unfolding *d-nom-def* **using** *lim-i-nom tendsto-divide div-by-1* **by** *fastforce*
qed

lemma *v-pos*: $\$v > 0$
unfolding *v-pres-def* **using** *v-futr-pos* **by** *auto*

lemma *v-1-iff-i-0*: $\$v = 1 \longleftrightarrow i = 0$

proof
assume $\$v = 1$
thus $i = 0$
unfolding *v-pres-def* **by** *simp*
next
assume $i = 0$
thus $\$v = 1$
unfolding *v-pres-def* **by** *simp*
qed

lemma *v-lt-1-iff-i-pos*: $\$v < 1 \longleftrightarrow i > 0$

proof
assume $\$v < 1$
thus $i > 0$
unfolding *v-pres-def* **by** (*simp add: v-futr-pos*)
next
assume $i > 0$
thus $\$v < 1$
unfolding *v-pres-def* **by** (*simp add: v-futr-pos*)
qed

lemma *v-i-nom*: $\$v = (1 + \$i^{\wedge}\{m\}/m)^{\wedge{-}m}$ **if** $m \neq 0$ **for** $m::nat$

proof –
have $\$v = (1 + i)^{\wedge{-}1}$
unfolding *v-pres-def* **using** *v-futr-pos powr-real-def that* **by** (*simp add: powr-neg-one*)
also have $\dots = ((1 + \$i^{\wedge}\{m\}/m)^{\wedge m})^{\wedge{-}1}$
using *i-nom-eff that* **by** *presburger*
also have $\dots = (1 + \$i^{\wedge}\{m\}/m)^{\wedge{-}m}$
using *powr-powr powr-realpow[THEN sym] v-futr-m-pos that* **by** *simp*
finally show *?thesis* .
qed

lemma *i-v*: $1 + i = \$v^{\wedge{-}1}$

unfolding *v-pres-def powr-real-def* **using** *v-futr-pos powr-neg-one* **by** *simp*

lemma *i-v-powr*: $(1 + i)^{\wedge a} = \$v^{\wedge{-}a}$ **for** $a::real$

by (subst i-v, subst powr-powr, simp)

lemma v-delta: $\ln (\$v) = - \δ

unfolding i-force-def v-pres-def **using** v-futr-pos **by** (simp add: ln-div)

lemma is-derive-vpow: $DERIV (\lambda t. \$v. \hat{\ }t) t :> - \$\delta * \$v. \hat{\ }t$

using v-delta has-real-derivative-powr2 v-pos **by** (metis mult.commute)

lemma d-nom-v: $\$d \hat{\ }m = m * (1 - \$v. \hat{\ } (1/m))$ **if** $m \neq 0$ **for** $m :: nat$

proof –

have $\$d \hat{\ }m = m * (1 - 1 / (1 + \$i \hat{\ }m / m))$

using d-nom-i-nom[THEN sym] **that** **by** force

also have $\dots = m * (1 - 1 / (1 + i. \hat{\ } (1/m)))$

using i-nom-i **that** powr-minus-divide **by** simp

also have $\dots = m * (1 - \$v. \hat{\ } (1/m))$

using v-pres-def v-futr-pos powr-divide **by** simp

finally show ?thesis .

qed

lemma d-nom-i-nom-v: $\$d \hat{\ }m = \$i \hat{\ }m * \$v. \hat{\ } (1/m)$ **if** $m \neq 0$ **for** $m :: nat$

unfolding d-nom-def v-pres-def **using** i-nom-i powr-divide v-futr-pos **that** **by** auto

lemma a-calc: $\$a \hat{\ }m - n = (1 - \$v \hat{\ }n) / \$i \hat{\ }m$ **if** $m \neq 0$ $i \neq 0$ **for** $n m :: nat$

proof –

have $\wedge l :: nat. l/m = (1/m) * l$ **by** simp

hence $\star: \wedge l :: nat. \$v. \hat{\ } (l/m) = (\$v. \hat{\ } (1/m)) \hat{\ } l$

using powr-powr powr-realpow v-pos **by** (metis powr-gt-zero)

hence $\$a \hat{\ }m - n = (\sum k < n * m. (\$v. \hat{\ } (1/m)) \hat{\ } (k+1 :: nat)) / m$

unfolding ann-def **by** presburger

also have $\dots = \$v. \hat{\ } (1/m) * (\sum k < n * m. (\$v. \hat{\ } (1/m)) \hat{\ } k) / m$

by (simp, subst sum-divide-distrib[THEN sym], subst sum-distrib-left[THEN sym], simp)

also have $\dots = \$v. \hat{\ } (1/m) * (((\$v. \hat{\ } (1/m)) \hat{\ } (n * m) - 1) / (\$v. \hat{\ } (1/m) - 1)) / m$

apply (subst geometric-sum[of $\$v. \hat{\ } (1/m)$ $n * m$]; simp?)

using powr-zero-eq-one[of $\$v$] v-pos v-1-iff-i-0 powr-inj **that**

by (smt (verit, del-insts) divide-eq-0-iff of-nat-eq-0-iff)

also have $\dots = ((\$v. \hat{\ } (1/m)) \hat{\ } (n * m) - 1) / (m * (\$v. \hat{\ } (1/m) - 1) / \$v. \hat{\ } (1/m))$

by (simp add: field-simps)

also have $\dots = (\$v \hat{\ }n - 1) / (m * (1 - 1 / \$v. \hat{\ } (1/m)))$

apply (subst \star [of $n * m :: nat$, THEN sym], simp only: of-nat-simps)

apply (subst nonzero-mult-div-cancel-right[where 'a=real, of m n], simp add: that)

apply (subst powr-realpow[OF v-pos])

apply (subst times-divide-eq-right[of - - $\$v. \hat{\ } (1/m)$, THEN sym])

using v-pos **by** (subst diff-divide-distrib[of - - $\$v. \hat{\ } (1/m)$], simp)

also have $\dots = (1 - \$v \hat{\ }n) / (m * (1 / \$v. \hat{\ } (1/m) - 1))$

using minus-divide-divide **by** (smt mult-minus-right)

also have $\dots = (1 - v^n) / i^m$
 unfolding *i-nom-def v-pres-def* using *v-futr-pos powr-divide* by *auto*
 finally show *?thesis* .
 qed

lemma *a-calc-i-0*: $\$a^m{-}n = n$ if $m \neq 0$ $i = 0$ for $n m :: nat$
 unfolding *ann-def v-pres-def* using *that* by *simp*

lemma *s-calc-i-0*: $\$s^m{-}n = n$ if $m \neq 0$ $i = 0$ for $n m :: nat$
 unfolding *acc-def* using *that* by *simp*

lemma *s-a*: $\$s^m{-}n = (1+i)^n * \$a^m{-}n$ if $m \neq 0$ for $n m :: nat$
 proof -

have $(1+i)^n * \$a^m{-}n = (\sum k < n*m. (1+i)^n * (v.^((k+1::nat)/m) / m))$
 unfolding *ann-def* using *sum-distrib-left* by *blast*
 also have $\dots = (\sum k < n*m. (1+i).^((n*m - Suc k)/m) / m)$
 proof -
 have $\bigwedge k :: nat. k < n*m \implies (1+i)^n * (v.^((k+1::nat)/m) / m) = (1+i).^((n*m - Suc k)/m) / m$
 - *unfolding v-pres-def*
 apply (*subst powr-realpow[THEN sym], simp add: v-futr-pos*)
 apply (*subst inverse-powr, simp add: v-futr-pos*)
 apply (*subst times-divide-eq-right, subst powr-add[THEN sym], simp add: that*)
 by (*subst of-nat-diff, simp add: Suc-le-eq, simp add: diff-divide-distrib that*)
 thus *?thesis* by (*meson lessThan-iff sum.cong*)
 qed
 also have $\dots = (\sum k < n*m. (1+i).^((k)/m) / m)$
 apply (*subst atLeast0LessThan[THEN sym]*) +
 by (*subst sum.atLeastLessThan-rev[THEN sym, of - n*m 0, simplified add-0-right], simp*)
 also have $\dots = \$s^m{-}n$
 unfolding *acc-def* by *simp*
 finally show *?thesis ..*
 qed

lemma *s-calc*: $\$s^m{-}n = ((1+i)^n - 1) / i^m$ if $m \neq 0$ $i \neq 0$ for $n m :: nat$
 using *that v-futr-pos*
 apply (*subst s-a, simp, subst a-calc; simp?*)
 apply (*rule disjI2*)
 apply (*subst right-diff-distrib, simp*)
 apply (*rule left-right-inverse-power*)
 unfolding *v-pres-def* by *auto*

lemma *a''-a*: $\$a''^m{-}n = (1+i).^((1/m) * \$a^m{-}n)$ if $m \neq 0$ for $m :: nat$
 unfolding *ann-def ann-due-def*
 apply (*subst sum-distrib-left, subst times-divide-eq-right, simp*)
 by (*subst i-v, subst powr-powr, subst powr-add[THEN sym], simp, subst add-divide-distrib,*

simp)

lemma $a-a''$: $\$a\{m\}-n = \$v.\wedge(1/m) * \$a''\{m\}-n$ **if** $m \neq 0$ **for** $m::nat$
unfolding *ann-def ann-due-def*
apply (*subst sum-distrib-left, subst times-divide-eq-right, simp*)
by (*subst pour-add[THEN sym], subst add-divide-distrib, simp*)

lemma $a''-calc-i-0$: $\$a''\{m\}-n = n$ **if** $m \neq 0$ $i = 0$ **for** $n m :: nat$
unfolding *ann-due-def v-pres-def* **using** *that* **by** *simp*

lemma $s''-calc-i-0$: $\$s''\{m\}-n = n$ **if** $m \neq 0$ $i = 0$ **for** $n m :: nat$
unfolding *acc-due-def* **using** *that* **by** *simp*

lemma $a''-calc$: $\$a''\{m\}-n = (1 - \$v\hat{n}) / \$d\{m\}$ **if** $m \neq 0$ $i \neq 0$ **for** $n m :: nat$

proof –

have $\$a''\{m\}-n = (1+i).\wedge(1/m) * ((1 - \$v\hat{n}) / \$i\{m\})$
using $a''-a$ *a-calc times-divide-eq-right* **that** **by** *simp*
also have $\dots = (1 - \$v\hat{n}) / (\$v.\wedge(1/m) * \$i\{m\})$
by (*subst i-v, subst pour-pour, simp, subst pour-minus-divide, simp*)
also have $\dots = (1 - \$v\hat{n}) / \$d\{m\}$
using *d-nom-i-nom-v* **that** **by** *simp*
finally show *?thesis* .

qed

lemma $s''-s$: $\$s''\{m\}-n = (1+i).\wedge(1/m) * \$s\{m\}-n$ **if** $m \neq 0$ **for** $m::nat$
unfolding *acc-def acc-due-def*
apply (*subst sum-distrib-left, subst times-divide-eq-right*)
by (*subst pour-add[THEN sym], simp, subst add-divide-distrib, simp*)

lemma $s-s''$: $\$s\{m\}-n = \$v.\wedge(1/m) * \$s''\{m\}-n$ **if** $m \neq 0$ **for** $m::nat$
unfolding *acc-def acc-due-def v-pres-def* **using** *v-futr-pos*
apply (*subst sum-distrib-left, subst times-divide-eq-right, simp*)
by (*subst inverse-pour, simp, subst pour-add[THEN sym], subst add-divide-distrib, simp*)

lemma $s''-calc$: $\$s''\{m\}-n = ((1+i)\hat{n} - 1) / \$d\{m\}$ **if** $m \neq 0$ $i \neq 0$ **for** $n m :: nat$

proof –

have $\$s''\{m\}-n = (1+i).\wedge(1/m) * ((1+i)\hat{n} - 1) / \$i\{m\}$
using $s''-s$ *s-calc times-divide-eq-right* **that** **by** *simp*
also have $\dots = ((1+i)\hat{n} - 1) / (\$v.\wedge(1/m) * \$i\{m\})$
by (*subst i-v, subst pour-pour, simp, subst pour-minus-divide, simp*)
also have $\dots = ((1+i)\hat{n} - 1) / \$d\{m\}$
using *d-nom-i-nom-v* **that** **by** *simp*
finally show *?thesis* .

qed

lemma $s''-a''$: $\$s''\{m\}-n = (1+i)\hat{n} * \$a''\{m\}-n$ **if** $m \neq 0$ **for** $m::nat$


```

using that s''-s a''-a s-a by simp

lemma a'-calc: $a'-n = (1 - $v. ^n) / $δ if i ≠ 0 n ≥ 0 for n::real
  unfolding ann-cont-def
  apply (rule integral-unique)
  using has-integral-powr2-from-0[OF v-pos - that(2)] v-delta v-1-iff-i-0 that
  by (smt minus-divide-divide)

lemma a'-calc-i-0: $a'-n = n if i = 0 n ≥ 0 for n::real
  unfolding ann-cont-def
  apply (subst iffD2[OF v-1-iff-i-0], simp add: that)
  by (simp add: integral-cong that)

lemma s'-calc: $s'-n = ((1+i). ^n - 1) / $δ if i ≠ 0 n ≥ 0 for n::real
  unfolding acc-cont-def
  apply (rule integral-unique)
  using has-integral-powr2-from-0[OF v-futr-pos - that(2)] i-force-def that
  by simp

lemma s'-calc-i-0: $s'-n = n if i = 0 n ≥ 0 for n::real
  unfolding acc-cont-def
  apply (subst ⟨i = 0⟩, simp)
  by (simp add: integral-cong that)

lemma s'-a': $s'-n = (1+i). ^n * $a'-n if n ≥ 0 for n::real
proof -
  have (1+i). ^n * $a'-n = integral {0..n} (λt. (1+i). ^{n-t})
    unfolding ann-cont-def
    using integrable-on-powr2-from-0-general[of $v n] v-pos v-futr-pos that
    apply (subst integral-mult, simp)
    apply (rule integral-cong)
    unfolding v-pres-def using inverse-powr powr-add[THEN sym] by smt
  also have ... = $s'-n
    unfolding acc-cont-def using v-futr-pos that
    apply (subst has-integral-interval-reverse[of 0 n, simplified, THEN integral-unique];
  simp?)
    by (rule continuous-on-powr; auto)
  finally show ?thesis ..
qed

lemma lim-m-a: (λm. $a^{m}-n) ⟶ $a'-n for n::nat
proof (rule LIMSEQ-imp-Suc)
  show (λm. $a^{Suc m}-n) ⟶ $a'-n
  proof (cases i = 0)
    case True
    show ?thesis
    using a-calc-i-0 a'-calc-i-0 True by simp
  next
  case False

```

```

show ?thesis
  using False v-pos delta-0-iff-i-0
  apply (subst a-calc; simp?)
  apply (subst a'-calc; simp?)
  apply (subst powr-realpow, simp)
  apply (rule tendsto-divide; simp?)
  by (rule LIMSEQ-Suc[OF lim-i-nom])
qed
qed

lemma lim-m-a'': ( $\lambda m. \$a''^{\sim\{m\}}-n$ )  $\longrightarrow$   $\$a'-n$  for  $n::nat$ 
proof (rule LIMSEQ-imp-Suc)
  show ( $\lambda m. \$a''^{\sim\{Suc\ m\}}-n$ )  $\longrightarrow$   $\$a'-n$ 
  proof (cases  $i = 0$ )
    case True
      show ?thesis
        using a''-calc-i-0 a'-calc-i-0 True by simp
    next
      case False
        show ?thesis
          using False v-pos delta-0-iff-i-0
          apply (subst a''-calc; simp?)
          apply (subst a'-calc; simp?)
          apply (subst powr-realpow, simp)
          apply (rule tendsto-divide; simp?)
          by (rule LIMSEQ-Suc[OF lim-d-nom])
  qed
qed

lemma lim-m-s: ( $\lambda m. \$s^{\sim\{m\}}-n$ )  $\longrightarrow$   $\$s'-n$  for  $n::nat$ 
proof (rule LIMSEQ-imp-Suc)
  show ( $\lambda m. \$s^{\sim\{Suc\ m\}}-n$ )  $\longrightarrow$   $\$s'-n$ 
  proof (cases  $i = 0$ )
    case True
      show ?thesis
        using s-calc-i-0 s'-calc-i-0 True by simp
    next
      case False
        show ?thesis
          using False v-futr-pos delta-0-iff-i-0
          apply (subst s-calc; simp?)
          apply (subst s'-calc; simp?)
          apply (subst powr-realpow, simp)
          apply (rule tendsto-divide; simp?)
          by (rule LIMSEQ-Suc[OF lim-i-nom])
  qed
qed

lemma lim-m-s'': ( $\lambda m. \$s''^{\sim\{m\}}-n$ )  $\longrightarrow$   $\$s'-n$  for  $n::nat$ 

```

```

proof (rule LIMSEQ-imp-Suc)
  show  $(\lambda m. \$s'' \wedge \{Suc\ m\} \cdot n) \longrightarrow \$s' \cdot n$ 
  proof (cases  $i = 0$ )
    case True
      show ?thesis
      using s''-calc-i-0 s'-calc-i-0 True by simp
    next
      case False
        show ?thesis
        using False v-futr-pos delta-0-iff-i-0
        apply (subst s''-calc; simp?)
        apply (subst s'-calc; simp?)
        apply (subst powr-realpow, simp)
        apply (rule tendsto-divide; simp?)
        by (rule LIMSEQ-Suc[OF lim-d-nom])
  qed
qed

```

```

lemma lim-n-a:  $(\lambda n. \$a \wedge \{m\} \cdot n) \longrightarrow \$a \wedge \{m\} \cdot \infty$  if  $m \neq 0$   $i > 0$  for  $m :: nat$ 
proof –
  have  $\$i \wedge \{m\} \neq 0$  using i-nom-pos-iff-i-pos that by smt
  moreover have  $(\lambda n. \$v \wedge n) \longrightarrow 0$ 
    using LIMSEQ-realpow-zero[of $v] v-pos v-lt-1-iff-i-pos that by simp
  ultimately show ?thesis
    using that apply (subst a-calc; simp?)
    unfolding perp-def apply (rule tendsto-divide; simp?)
    using tendsto-diff[where a=1 and b=0] by auto
qed

```

```

lemma lim-n-a'':  $(\lambda n. \$a'' \wedge \{m\} \cdot n) \longrightarrow \$a'' \wedge \{m\} \cdot \infty$  if  $m \neq 0$   $i > 0$  for  $m :: nat$ 
proof –
  have  $\$d \wedge \{m\} \neq 0$  using d-nom-pos-iff-i-pos that by smt
  moreover have  $(\lambda n. \$v \wedge n) \longrightarrow 0$ 
    using LIMSEQ-realpow-zero[of $v] v-pos v-lt-1-iff-i-pos that by simp
  ultimately show ?thesis
    using that apply (subst a''-calc; simp?)
    unfolding perp-due-def apply (rule tendsto-divide; simp?)
    using tendsto-diff[where a=1 and b=0] by auto
qed

```

```

lemma lsm-lam:  $(I \wedge \{l\} s) \wedge \{m\} \cdot n = (1+i) \wedge n * (I \wedge \{l\} a) \wedge \{m\} \cdot n$ 
if  $l \neq 0$   $m \neq 0$  for  $l\ n\ m :: nat$ 
unfolding acc-incr-def ann-incr-def v-pres-def using v-futr-pos powr-realpow
apply (subst inverse-powr, simp)
apply (subst sum-distrib-left)
by (subst minus-real-def, subst powr-add, subst times-divide-eq-right, subst mult.assoc,
simp)

```

```

lemma lam-calc:  $(Ia) \wedge \{m\} \cdot n = (\sum j < n. (j+1)/m * (\sum k = j*m .. < (j+1)*m. \$v \wedge ((k+1)/m)))$ 

```

if $m \neq 0$ **for** $n\ m :: \text{nat}$
proof –
let $?I = \{..<n\}$
let $?A = \lambda j. \{j*m..<(j+1)*m\}$
let $?g = \lambda k. \$v. \wedge((k+1)::\text{nat})/m * \lceil(k+1)::\text{nat}\rceil / m$
have $\$(Ia) \wedge\{m\}-n = (\sum j<n. \sum k=j*m..<(j+1)*m. \$v. \wedge((k+1)/m) * \lceil(k+1)/m\rceil / m)$
unfolding *ann-incr-def* **using** *seq-part-multiple* **that**
apply (*simp only: mult-1*)
by (*subst sum.UNION-disjoint*[of $?I\ ?A\ ?g$, *THEN sym*]; *simp*)
also have $\dots = (\sum j<n. (j+1)/m * (\sum k=j*m..<(j+1)*m. \$v. \wedge((k+1)/m)))$
proof –
{ **fix** $j\ k$
assume $j*m \leq k \wedge k < (j+1)*m$
hence $j*m < k+1 \wedge k+1 \leq (j+1)*m$ **by** *force*
hence $j < (k+1)/m \wedge (k+1)/m \leq j+1$
using *pos-less-divide-eq pos-divide-le-eq of-nat-less-iff of-nat-le-iff* **that**
by (*smt (verit) of-nat-le-0-iff of-nat-mult*)
hence $\lceil(k+1)/m\rceil = j+1$
by (*simp add: ceiling-unique*) **}**
hence $\bigwedge j\ k. j*m \leq k \wedge k < (j+1)*m \implies \lceil(k+1)/m\rceil = j+1$
by (*metis (no-types) of-nat-1 of-nat-add*)
with *v-pos* **show** *?thesis*
apply –
apply (*rule sum.cong, simp*)
apply (*subst sum-distrib-left, rule sum.cong; simp*)
by (*smt (verit, ccfv-SIG) of-int-1 of-int-diff of-int-of-nat-eq*)
qed
finally show *?thesis* .
qed

lemma *Ism-calc*: $\$(Is) \wedge\{m\}-n = (\sum j<n. (j+1)/m * (\sum k=j*m..<(j+1)*m. (1+i). \wedge(n-(k+1)/m)))$
if $m \neq 0$ **for** $n\ m :: \text{nat}$
using *v-pos* **that**
apply (*subst Iism-Ilam; simp*)
apply (*subst Iam-calc[simplified]; simp?*)
apply ((*subst sum-distrib-left, rule sum.cong; simp*))+
unfolding *v-pres-def* **using** *v-futr-pos*
apply (*subst inverse-powr; simp*)
apply (*subst powr-realpow[THEN sym], simp*)
by (*subst powr-add[THEN sym]; simp*)

lemma *Imam-calc-aux*: $\$(I\{m\}a) \wedge\{m\}-n = (\sum k<n*m. \$v. \wedge((k+1)/m) * (k+1) / m^2)$
if $m \neq 0$ **for** $m::\text{nat}$
unfolding *ann-incr-def power-def*
apply (*rule sum.cong, simp*)
apply (*subst of-nat-mult*)
using *v-pos* **that**

apply (*subst nonzero-mult-div-cancel-left*, *simp*)
by (*subst ceiling-of-nat*; *simp*)

lemma *Imam-calc*:

$\$(I^{\wedge}\{m\}a)^{\wedge}\{m\}-n = (\$v.^{\wedge}(1/m) * (1 - (n*m+1)*\$v.^{\wedge}n + n*m*\$v.^{\wedge}(n+1/m))) / (m*(1-\$v.^{\wedge}(1/m)))^{\wedge}2$

if $i \neq 0$ $m \neq 0$ **for** n $m :: \text{nat}$

proof –

have \star : $\$v.^{\wedge}(1/m) > 0$ **using** *v-pos* **by** *force*

hence $\$(I^{\wedge}\{m\}a)^{\wedge}\{m\}-n = (\sum k < n*m. (k+1)*(\$v.^{\wedge}(1/m))^{\wedge}(k+1)) / m^{\wedge}2$

using *that*

apply (*subst Imam-calc-aux*, *simp*)

apply (*subst sum-divide-distrib*[*THEN sym*], *simp*)

apply (*rule sum.cong*; *simp*)

using *powr-realpow*[*THEN sym*] *powr-powr* **by** (*simp add: add-divide-distrib powr-add*)

also have $\dots = \$v.^{\wedge}(1/m) * (\sum k < n*m. (k+1)*(\$v.^{\wedge}(1/m))^{\wedge}k) / m^{\wedge}2$

by (*subst sum-distrib-left*, *simp add: that*, *rule sum.cong*; *simp*)

also have $\dots = \$v.^{\wedge}(1/m) *$

$((1 - (n*m+1)*(\$v.^{\wedge}(1/m))^{\wedge}(n*m) + n*m*(\$v.^{\wedge}(1/m))^{\wedge}(n*m+1)) / (1 - \$v.^{\wedge}(1/m))^{\wedge}2) / m^{\wedge}2$

using *v-pos v-1-iff-i-0* **that** **by** (*subst geometric-increasing-sum*; *simp?*)

also have $\dots = (\$v.^{\wedge}(1/m) * (1 - (n*m+1)*\$v.^{\wedge}n + n*m*\$v.^{\wedge}(n+1/m))) / (m*(1-\$v.^{\wedge}(1/m)))^{\wedge}2$

using \star

apply (*subst powr-realpow*[*of* $\$v.^{\wedge}(1/m)$, *THEN sym*], *simp*)**+**

apply (*subst powr-powr*)**+**

apply (*subst times-divide-eq-right*[*THEN sym*], *subst divide-divide-eq-left*)

apply (*subst power-mult-distrib*)

using *powr-eq-one-iff-gen v-pos v-1-iff-i-0* **apply** (*simp add: field-simps*)

by (*subst powr-realpow*, *simp*)**+**, *simp*)

finally show *?thesis* .

qed

lemma *Imam-calc-i-0*: $\$(I^{\wedge}\{m\}a)^{\wedge}\{m\}-n = (n*m+1)*n / (2*m)$ **if** $i = 0$ $m \neq 0$ **for** n $m :: \text{nat}$

proof –

have $\$(I^{\wedge}\{m\}a)^{\wedge}\{m\}-n = (\sum k < n*m. \$v.^{\wedge}((k+1)/m) * (k+1)) / m^{\wedge}2$

by (*subst Imam-calc-aux*, *simp-all add: that*)

also have $\dots = (\sum k < n*m. k+1) / m^{\wedge}2$

apply (*subst v-1-iff-i-0*[*THEN iffD2*], *simp-all add: that*)

by (*subst sum-divide-distrib*[*THEN sym*], *simp*)

also have $\dots = (n*m*(n*m+1) \text{ div } 2) / m^{\wedge}2$

apply (*subst Suc-eq-plus1*[*THEN sym*], *subst sum-bounds-lt-plus1*[*of id*, *simplified*])

by (*subst Sum-Icc-nat*, *simp*)

also have $\dots = (n*m+1)*n / (2*m)$

apply (*subst real-of-nat-div*, *simp*)

using *that* **by** (*subst power2-eq-square*, *simp add: field-simps*)

finally show ?thesis .
qed

lemma *Imsm-calc*:

$\$(I\{m\}s)\{m\}-n = ((1+i).\wedge(n+1/m) - (n*m+1)*(1+i).\wedge(1/m) + n*m) / (m*((1+i).\wedge(1/m)-1))^2$
if $i \neq 0$ $m \neq 0$ for n $m :: nat$

proof -

have $\$(I\{m\}a)\{m\}-n = (\$v\wedge n * ((1+i).\wedge(n+1/m) - (n*m+1)*(1+i).\wedge(1/m) + n*m)) / (m*((1+i).\wedge(1/m)-1))^2$

proof -

have $\$(I\{m\}a)\{m\}-n = (\$v.\wedge(1/m) * (1 - (n*m+1)*\$v\wedge n + n*m*\$v.\wedge(n+1/m))) / (m*(1-\$v.\wedge(1/m)))^2$
using that by (subst *Imam-calc*; simp)

also have ... = $(1 - (n*m+1)*\$v\wedge n + n*m*\$v.\wedge(n+1/m)) / (\$v.\wedge(1/m)*(m*(\$v.\wedge(-1/m)-1))^2)$

apply (subgoal-tac $\$v.\wedge(-1/m) = 1 / \$v.\wedge(1/m)$, erule ssubst)

apply ((subst power2-eq-square)+, simp add: field-simps that)

by (simp add: powr-minus-divide)

also have ... =

$(\$v.\wedge(n+1/m) * (\$v.\wedge(-n-1/m) - (n*m+1)*\$v.\wedge(-1/m) + n*m)) / (\$v.\wedge(1/m)*(m*(\$v.\wedge(-1/m)-1))^2)$

apply (subgoal-tac $\$v.\wedge(-n-1/m) = 1 / \$v.\wedge(n+1/m)$ $\$v.\wedge(-1/m) = \$v\wedge n / \$v.\wedge(n+1/m)$)

apply ((erule ssubst)+, simp-all add: field-simps)

using v-pos

apply (simp add: powr-diff[THEN sym] powr-realpow[THEN sym])

by (smt powr-minus-divide)

also have ... =

$(\$v\wedge n * (\$v.\wedge(-n-1/m) - (n*m+1)*\$v.\wedge(-1/m) + n*m)) / ((m*(\$v.\wedge(-1/m)-1))^2)$

apply (subst powr-add[of - n 1/m])

using v-pos powr-realpow by simp

also have ... =

$(\$v\wedge n * ((1+i).\wedge(n+1/m) - (n*m+1)*(1+i).\wedge(1/m) + n*m)) / ((m*((1+i).\wedge(1/m)-1))^2)$

apply (subgoal-tac $-n-1/m = -(n+1/m) - 1/m = -(1/m)$, (erule ssubst)+)

apply (subst i-v-powr[THEN sym])+

by simp-all

finally show ?thesis .

qed

thus ?thesis

apply -

using that v-futr-pos

apply (subst *Imsm-Ilam*, simp)

apply (erule ssubst, simp)

apply (rule disjI2)

by (subst power-mult-distrib[THEN sym], simp add: v-pres-def)

qed

lemma *Imsm-calc-i-0*: $\$(I\{m\}s)\{m\}-n = (n*m+1)*n / (2*m)$ if $i = 0$ $m \neq 0$
for n $m :: nat$

using *that*
apply (*subst Ilsm-Ilam, simp*)
by (*subst Imam-calc-i-0; simp*)

lemma *Ila''m-Ilam*: $\$(I^{\wedge}\{l\}a'')^{\wedge}\{m\}-n = (1+i).\wedge(1/m) * \$(I^{\wedge}\{l\}a)^{\wedge}\{m\}-n$
if $l \neq 0$ **for** $m \neq 0$ **for** $l\ m\ n :: \text{nat}$
unfolding *ann-incr-def ann-due-incr-def* **using** *that*
apply (*subst i-v, subst powr-powr, simp*)
apply (*subst sum-distrib-left*)
apply (*rule sum.cong; simp*)
apply (*rule disjI2*)
by (*smt (verit) add-divide-distrib powr-add*)

lemma *Ia''m-calc*: $\$(Ia'')^{\wedge}\{m\}-n = (\sum j < n. (j+1)/m * (\sum k=j*m..<(j+1)*m. \$v.\wedge(k/m)))$
if $m \neq 0$ **for** $n\ m :: \text{nat}$
using *that*
apply (*subst Ila''m-Ilam; simp del: One-nat-def*)
apply (*subst Iam-calc; simp*)
apply (*subst sum-distrib-left*)
apply (*rule sum.cong; simp*)
apply (*subst sum-distrib-left*)
apply (*rule sum.cong; simp*)
apply (*subst i-v-powr*)
using *powr-add*[of $\$v$, *THEN sym*] **by** (*simp add: field-simps*)

lemma *Ima''m-calc-aux*: $\$(I^{\wedge}\{m\}a'')^{\wedge}\{m\}-n = (\sum k < n*m. \$v.\wedge(k/m) * (k+1) / m^{\wedge}2)$
if $m \neq 0$ **for** $m :: \text{nat}$
using *that*
apply (*subst Ila''m-Ilam, simp*)
apply (*subst Imam-calc-aux, simp*)
apply (*subst sum-distrib-left*)
apply (*rule sum.cong; simp*)
using *powr-add*[of $\$v$, *THEN sym*] *i-v-powr* **by** (*simp add: field-simps*)

lemma *Ima''m-calc*: $\$(I^{\wedge}\{m\}a'')^{\wedge}\{m\}-n = (1 - (n*m+1)*\$v^{\wedge}n + n*m*\$v.\wedge(n+1/m)) / (m*(1-\$v.\wedge(1/m)))^{\wedge}2$
if $i \neq 0$ $m \neq 0$ **for** $n\ m :: \text{nat}$
using *that v-pos*
apply (*subst Ila''m-Ilam, simp*)
apply (*subst Imam-calc; simp*)
apply (*rule disjI2*)
by (*subst i-v, subst powr-powr, subst powr-add*[*THEN sym*], *simp*)

lemma *Ils''m-Ilsm*: $\$(I^{\wedge}\{l\}s'')^{\wedge}\{m\}-n = (1+i).\wedge(1/m) * \$(I^{\wedge}\{l\}s)^{\wedge}\{m\}-n$
if $l \neq 0$ $m \neq 0$ **for** $l\ m\ n :: \text{nat}$
unfolding *acc-incr-def acc-due-incr-def* **using** *that*
apply (*subst sum-distrib-left*)

apply (rule *sum.cong*; *simp*)
apply (rule *disjI2*)
by (*subst powr-add*[*THEN sym*], *subst add-divide-distrib*, *simp*)

lemma *Ims''m-calc*:

$\$(I^{\wedge\{m\}}s')^{\wedge\{m\}}-n =$
 $(1+i).^{\wedge(1/m)} * ((1+i).^{\wedge(n+1/m)} - (n*m+1)*(1+i).^{\wedge(1/m)} + n*m) /$
 $(m*((1+i).^{\wedge(1/m)}-1))^{\wedge 2}$
if $i \neq 0$ $m \neq 0$ **for** n $m :: \text{nat}$
using that by (*simp add: Ils''m-Ilsm Imsm-calc*)

lemma *lim-Imam*: $(\lambda n. \$(I^{\wedge\{m\}}a)^{\wedge\{m\}}-n) \longrightarrow 1 / (\$i^{\wedge\{m\}}*\$d^{\wedge\{m\}})$ **if** $m \neq 0$ $i > 0$ **for** $m :: \text{nat}$

proof –

have $(\lambda n. \$(I^{\wedge\{m\}}a)^{\wedge\{m\}}-n) =$
 $(\lambda n. \$v.^{\wedge(1/m)} * (1 - (n*m+1)*\$v^{\wedge n} + n*m*\$v.^{\wedge(n+1/m)}) / (m*(1-\$v.^{\wedge(1/m)}))^{\wedge 2}$
using that by (*subst Imam-calc*; *simp*)
moreover have $(\lambda n. \$v.^{\wedge(1/m)} * (1 - (n*m+1)*\$v^{\wedge n} + n*m*\$v.^{\wedge(n+1/m)})$
 $/ (m*(1-\$v.^{\wedge(1/m)}))^{\wedge 2}$
 $\longrightarrow 1 / (\$i^{\wedge\{m\}}*\$d^{\wedge\{m\}})$

proof –

have $\star: |\$v| < 1$
using *v-lt-1-iff-i-pos v-pos* **that by force**
hence $(\lambda n. (n*m+1)*\$v^{\wedge n}) \longrightarrow 0$
apply (*subst tendsto-cong*[*of -* ($\lambda n. n*m*\$v^{\wedge n} + \$v^{\wedge n}$)])
apply (rule *always-eventually*, rule *allI*)
apply (*simp add: distrib-right*)
apply (*subgoal-tac* $0 = 0 + 0$, *erule ssubst*, *intro tendsto-intros*; *simp*)
apply (*subst mult.commute*, *subst mult.assoc*)
apply (*subgoal-tac* $0 = \text{real } m * 0$, *erule ssubst*, *intro tendsto-intros*; *simp?*)
by (rule *powser-times-n-limit-0*; *simp*)
moreover have $(\lambda n. n*m*\$v.^{\wedge(n+1/m)}) \longrightarrow 0$
apply (*subst tendsto-cong*[*of -* ($\lambda n. (m*\$v.^{\wedge(1/m)})*(n*\$v^{\wedge n})$)])
apply (rule *always-eventually*, rule *allI*)
apply (*subst powr-add*, *subst powr-realpow*; *simp add: v-pos*)
apply (*subgoal-tac* $0 = m*\$v.^{\wedge(1/m)} * 0$, *erule ssubst*, *intro tendsto-intros*;
simp?)
by (rule *powser-times-n-limit-0*, *simp add: \star*)
ultimately have $(\lambda n. \$v.^{\wedge(1/m)} * (1 - (n*m+1)*\$v^{\wedge n} + n*m*\$v.^{\wedge(n+1/m)})$
 $/ (m*(1-\$v.^{\wedge(1/m)}))^{\wedge 2}$
 $\longrightarrow \$v.^{\wedge(1/m)} * (1 - 0 + 0) / (m*(1-\$v.^{\wedge(1/m)}))^{\wedge 2}$
using *v-lt-1-iff-i-pos v-pos* **that by** (*intro tendsto-intros*; *simp*)
thus *?thesis*
unfolding *i-nom-def* **using** *v-pos* **that**
apply (*subst i-v-powr*, *subst powr-minus-divide*, *subst d-nom-v*; *simp*)
by (*subst(asm)*(2) *power2-eq-square*, *simp add: field-simps*)
qed
ultimately show *?thesis* **by** *simp*
qed

lemma *perp-incr-calc*: $\$(I^{\wedge}\{m\}a)^{\wedge}\{m\}-\infty = 1 / (\$i^{\wedge}\{m\}*\$d^{\wedge}\{m\})$ **if** $m \neq 0$ $i > 0$ **for** $m::nat$

unfolding *perp-incr-def* **by** (rule *limI*, rule *lim-Imam*; simp add: *that*)

lemma *lim-Ima''m*: $(\lambda n. \$(I^{\wedge}\{m\}a'')^{\wedge}\{m\}-n) \longrightarrow 1 / (\$d^{\wedge}\{m\})^{\wedge}2$ **if** $m \neq 0$ $i > 0$ **for** $m::nat$

unfolding *perp-due-incr-def* **using** *that*

apply (subst *Ila''m-Ilam*, simp, subst *mult.commute*, subst *i-v-powr*, subst *powr-minus-divide*)

apply (subgoal-tac $1/(\$d^{\wedge}\{m\})^{\wedge}2 = (1/(\$i^{\wedge}\{m\}*\$d^{\wedge}\{m\}))*(1/\$v.^{\wedge}(1/m))$, erule *ssubst*)

apply (intro *tendsto-intros*, simp add: *lim-Imam*)

by (subst *power2-eq-square*, subst(1) *d-nom-i-nom-v*; simp add: *field-simps that*)

lemma *perp-due-incr-calc*: $\$(I^{\wedge}\{m\}a'')^{\wedge}\{m\}-\infty = 1 / (\$d^{\wedge}\{m\})^{\wedge}2$ **if** $m \neq 0$ $i > 0$ **for** $m::nat$

unfolding *perp-due-incr-def* **by** (rule *limI*, rule *lim-Ima''m*; simp add: *that*)

end

end

theory *Survival-Model*

imports *HOL-Library.Rewrite* *HOL-Library.Extended-Nonnegative-Real* *HOL-Library.Extended-Real* *HOL-Probability.Probability Preliminaries*

begin

5 Survival Model

The survival model is built on the probability space \mathfrak{M} . Additionally, the random variable $X : space \mathfrak{M} \rightarrow \mathbb{R}$ is introduced, which represents the age at death.

locale *prob-space-actuary* = *MM-PS*: *prob-space* \mathfrak{M} **for** \mathfrak{M}

— Since the letter M may be used as a commutation function, adopt the letter \mathfrak{M} instead as a notation for the measure space.

locale *survival-model* = *prob-space-actuary* +

fixes $X :: 'a \Rightarrow real$

assumes $X-RV[simp]$: *MM-PS.random-variable* (borel :: real measure) X

and $X-pos-AE[simp]$: *AE* ξ in \mathfrak{M} . $X \xi > 0$

begin

5.1 General Theory of Survival Model

interpretation *distrX-RD*: *real-distribution* *distr* \mathfrak{M} borel X

using *MM-PS.real-distribution-distr* **by** *simp*

lemma $X-le-event[simp]$: $\{\xi \in space \mathfrak{M}. X \xi \leq x\} \in MM-PS.events$

by *measurable simp*

lemma *X-gt-event[simp]*: $\{\xi \in \text{space } \mathfrak{M}. X \xi > x\} \in \text{MM-PS.events}$
by *measurable simp*

lemma *X-compl-le-gt*: $\text{space } \mathfrak{M} - \{\xi \in \text{space } \mathfrak{M}. X \xi \leq x\} = \{\xi \in \text{space } \mathfrak{M}. X \xi > x\}$ **for** $x::\text{real}$

proof –

have $\text{space } \mathfrak{M} - \{\xi \in \text{space } \mathfrak{M}. X \xi \leq x\} = \text{space } \mathfrak{M} - (X - \{..x\})$ **by** *blast*
also have $\dots = (X - \{x<..\}) \cap \text{space } \mathfrak{M}$ **using** *vimage-compl-atMost* **by** *fastforce*

also have $\dots = \{\xi \in \text{space } \mathfrak{M}. X \xi > x\}$ **by** *blast*

finally show *?thesis* .

qed

lemma *X-compl-gt-le*: $\text{space } \mathfrak{M} - \{\xi \in \text{space } \mathfrak{M}. X \xi > x\} = \{\xi \in \text{space } \mathfrak{M}. X \xi \leq x\}$ **for** $x::\text{real}$

using *X-compl-le-gt* **by** *blast*

5.1.1 Introduction of Survival Function for X

Note that *ccdf* (*distr* \mathfrak{M} *borel* X) is the survival (distributive) function for X .

lemma *ccdfX-0-1*: *ccdf* (*distr* \mathfrak{M} *borel* X) $0 = 1$

apply (*rewrite* *MM-PS.ccdf-distr-P*, *simp*)

using *X-pos-AE* *MM-PS.prob-space*

using *MM-PS.prob-Collect-eq-1* *X-gt-event* **by** *presburger*

lemma *ccdfX-unborn-1*: *ccdf* (*distr* \mathfrak{M} *borel* X) $x = 1$ **if** $x \leq 0$

proof (*rule antisym*)

show *ccdf* (*distr* \mathfrak{M} *borel* X) $x \leq 1$ **using** *MM-PS.ccdf-distr-P* **by** *simp*

show *ccdf* (*distr* \mathfrak{M} *borel* X) $x \geq 1$

proof –

have *ccdf* (*distr* \mathfrak{M} *borel* X) $x \geq$ *ccdf* (*distr* \mathfrak{M} *borel* X) 0

using *finite-borel-measure.ccdf-nonincreasing* *distrX-RD.finite-borel-measure-M* **that** **by** *simp*

also have *ccdf* (*distr* \mathfrak{M} *borel* X) $0 = 1$ **using** *ccdfX-0-1* **that** **by** *simp*

finally show *?thesis* .

qed

qed

definition *death-pt* :: *ereal* (ψ)

where $\psi \equiv \text{Inf } (\text{ereal } \{x \in \mathbb{R}. \text{ccdf } (\text{distr } \mathfrak{M} \text{ borel } X) x = 0\})$

— This is my original notation, which is used to develop life insurance mathematics rigorously.

lemma *psi-nonneg*: $\psi \geq 0$

unfolding *death-pt-def*

proof (*rule Inf-greatest*)

fix $x'::\text{ereal}$

assume $x' \in \text{ereal } \{x \in \mathbb{R}. \text{cdf} (\text{distr } \mathfrak{M} \text{ borel } X) x = 0\}$
then obtain $x::\text{real}$ **where** $x' = \text{ereal } x$ **and** $\text{cdf} (\text{distr } \mathfrak{M} \text{ borel } X) x = 0$ **by**
blast
hence $\text{cdf} (\text{distr } \mathfrak{M} \text{ borel } X) 0 > \text{cdf} (\text{distr } \mathfrak{M} \text{ borel } X) x$ **using** cdfX-0-1
X-pos-AE **by** *simp*
hence $x \geq 0$
using *mono-invE finite-borel-measure.cdf-nonincreasing distrX-RD.finite-borel-measure-M*
X-RV
by (*smt(verit)*)
thus $x' \geq 0$ **using** $\langle x' = \text{ereal } x \rangle$ **by** *simp*
qed

lemma *cdfX-beyond-0*: $\text{cdf} (\text{distr } \mathfrak{M} \text{ borel } X) x = 0$ **if** $x > \psi$ **for** $x::\text{real}$
proof –
have $\text{ereal } \{y \in \mathbb{R}. \text{cdf} (\text{distr } \mathfrak{M} \text{ borel } X) y = 0\} \neq \{\}$ **using** *death-pt-def* **that**
by *force*
hence $\exists y' \in (\text{ereal } \{y \in \mathbb{R}. \text{cdf} (\text{distr } \mathfrak{M} \text{ borel } X) y = 0\}). y' < \text{ereal } x$
using *that unfolding death-pt-def* **by** (*rule cInf-lessD*)
then obtain y'
where $y' \in (\text{ereal } \{y \in \mathbb{R}. \text{cdf} (\text{distr } \mathfrak{M} \text{ borel } X) y = 0\})$ **and** $y' < \text{ereal } x$
by *blast*
then obtain $y::\text{real}$
where $y' = \text{ereal } y$ **and** $\text{cdf} (\text{distr } \mathfrak{M} \text{ borel } X) y = 0$ **and** $\text{ereal } y < \text{ereal } x$
by *blast*
hence $\text{cdf} (\text{distr } \mathfrak{M} \text{ borel } X) y = 0$ **and** $y < x$ **by** *simp-all*
hence $\text{cdf} (\text{distr } \mathfrak{M} \text{ borel } X) x \leq 0$
using *finite-borel-measure.cdf-nonincreasing distrX-RD.finite-borel-measure-M*
X-RV
by (*metis order-less-le*)
thus *?thesis* **using** *finite-borel-measure.cdf-nonneg distrX-RD.finite-borel-measure-M*
X-RV **by** *smt*
qed

lemma *cdfX-psi-0*: $\text{cdf} (\text{distr } \mathfrak{M} \text{ borel } X) (\text{real-of-ereal } \psi) = 0$ **if** $\psi < \infty$
proof –
have $|\psi| \neq \infty$ **using** *that psi-nonneg* **by** *simp*
then obtain $u::\text{real}$ **where** $\psi = \text{ereal } u$ **using** *ereal-real'* **by** *blast*
hence $\text{real-of-ereal } \psi = u$ **by** *simp*
moreover have $\text{cdf} (\text{distr } \mathfrak{M} \text{ borel } X) u = 0$
proof –
have $\bigwedge x::\text{real}. x \neq u \implies x \in \{u < ..\} \implies \text{cdf} (\text{distr } \mathfrak{M} \text{ borel } X) x = 0$
by (*rule cdfX-beyond-0, simp add: \langle \psi = \text{ereal } u \rangle*)
hence $(\text{cdf} (\text{distr } \mathfrak{M} \text{ borel } X) \longrightarrow 0)$ (*at-right u*)
apply –
by (*rule iffD2[OF Lim-cong-within[where ?g=(\lambda x.0)]], simp-all+*)
moreover have $(\text{cdf} (\text{distr } \mathfrak{M} \text{ borel } X) \longrightarrow \text{cdf} (\text{distr } \mathfrak{M} \text{ borel } X) u)$
(*at-right u*)
using *finite-borel-measure.cdf-is-right-cont distrX-RD.finite-borel-measure-M*
continuous-within X-RV **by** *blast*

ultimately show *?thesis* using *tendsto-unique trivial-limit-at-right-real* by
blast
 qed
 ultimately show *?thesis* by *simp*
 qed

lemma *ccdfX-0-equiv*: $ccdf (distr \mathfrak{M} \text{ borel } X) x = 0 \longleftrightarrow x \geq \psi$ for $x::real$
proof

assume $ccdf (distr \mathfrak{M} \text{ borel } X) x = 0$
 thus $ereal x \geq \psi$ **unfolding** *death-pt-def* by (*simp add: INF-lower*)
next
 assume $\psi \leq ereal x$
 hence $\psi = ereal x \vee \psi < ereal x$ **unfolding** *less-eq-ereal-def* by *auto*
 thus $ccdf (distr \mathfrak{M} \text{ borel } X) x = 0$
proof
 assume $\star: \psi = ereal x$
 hence $\psi < \infty$ by *simp*
 moreover have *real-of-ereal* $\psi = x$ using \star by *simp*
 ultimately show $ccdf (distr \mathfrak{M} \text{ borel } X) x = 0$ using *ccdfX-psi-0* by *simp*
next
 assume $\psi < ereal x$
 thus $ccdf (distr \mathfrak{M} \text{ borel } X) x = 0$ by (*rule ccdfX-beyond-0*)
 qed
 qed

lemma *psi-pos[simp]*: $\psi > 0$

proof (*rule not-le-imp-less, rule notI*)

show $\psi \leq (0::ereal) \implies False$

proof –

assume $\psi \leq (0::ereal)$

hence $ccdf (distr \mathfrak{M} \text{ borel } X) 0 = 0$ using *ccdfX-0-equiv* by (*simp add: zero-ereal-def*)

moreover have $ccdf (distr \mathfrak{M} \text{ borel } X) 0 = 1$ using *ccdfX-0-1* by *simp*

ultimately show *False* by *simp*

qed

qed

corollary *psi-pos'[simp]*: $\psi > ereal 0$

using *psi-pos zero-ereal-def* by *presburger*

5.1.2 Introduction of Future-Lifetime Random Variable $T(x)$

definition *alive* :: $real \Rightarrow 'a \text{ set}$

where $alive x \equiv \{\xi \in space \mathfrak{M}. X \xi > x\}$

lemma *alive-event[simp]*: $alive x \in MM\text{-}PS.events$ for $x::real$

unfolding *alive-def* by *simp*

lemma *X-alivex-measurable[measurable, simp]*: $X \in \text{borel-measurable } (\mathfrak{M} \downarrow alive$

x) for $x::real$
unfolding *cond-prob-space-def* by (*measurable; simp add: measurable-restrict-space1*)

definition *futr-life* :: $real \Rightarrow ('a \Rightarrow real) (T)$
where $T x \equiv (\lambda \xi. X \xi - x)$
— Note that $T(x) : space \mathfrak{M} \rightarrow \mathbf{R}$ represents the time until death of a person aged x .

lemma *T0-eq-X[simp]*: $T 0 = X$
unfolding *futr-life-def* by *simp*

lemma *Tx-measurable[measurable, simp]*: $T x \in borel\text{-}measurable \mathfrak{M}$ for $x::real$
unfolding *futr-life-def* by (*simp add: borel-measurable-diff*)

lemma *Tx-alivex-measurable[measurable, simp]*: $T x \in borel\text{-}measurable (\mathfrak{M} \downarrow alive\ x)$ for $x::real$
unfolding *futr-life-def* by (*simp add: borel-measurable-diff*)

lemma *alive-T*: $alive\ x = \{\xi \in space\ \mathfrak{M}. T\ x\ \xi > 0\}$ for $x::real$
unfolding *alive-def futr-life-def* by *force*

lemma *alivex-Tx-pos[simp]*: $0 < T\ x\ \xi$ if $\xi \in space\ (\mathfrak{M} \downarrow alive\ x)$ for $x::real$
using *MM-PS.space-cond-prob-space alive-T* that by *simp*

lemma *PT0-eq-PX-lborel*: $\mathcal{P}(\xi\ in\ \mathfrak{M}. T\ 0\ \xi \in A \mid T\ 0\ \xi > 0) = \mathcal{P}(\xi\ in\ \mathfrak{M}. X\ \xi \in A)$
if $A \in sets\ lborel$ for $A :: real\ set$
apply (*rewrite MM-PS.cond-prob-AE-prob, simp-all*)
using that *X-RV measurable-lborel1 predE pred-sets2* by *blast*

5.1.3 Actuarial Notations on the Survival Model

definition *survive* :: $real \Rightarrow real \Rightarrow real (\$p'\{-\&\} [0,0] 200)$
where $\$p'\{t\&x\} \equiv cdf\ (distr\ (\mathfrak{M} \downarrow alive\ x)\ borel\ (T\ x))\ t$
— the probability that a person aged x will survive for t years
— Note that the function $\$p'\{\cdot\&x\}$ is the survival function on $(\mathfrak{M} \downarrow alive\ x)$ for the random variable $T(x)$.
— The parameter t is usually nonnegative, but theoretically it can be negative.

abbreviation *survive-1* :: $real \Rightarrow real (\$p'\{- [101] 200)$
where $\$p'\{-x\} \equiv \$p'\{1\&x\}$

definition *die* :: $real \Rightarrow real \Rightarrow real (\$q'\{-\&\} [0,0] 200)$
where $\$q'\{t\&x\} \equiv cdf\ (distr\ (\mathfrak{M} \downarrow alive\ x)\ borel\ (T\ x))\ t$
— the probability that a person aged x will die within t years
— Note that the function $\$q'\{\cdot\&x\}$ is the cumulative distributive function on $(\mathfrak{M} \downarrow alive\ x)$ for the random variable $T(x)$.
— The parameter t is usually nonnegative, but theoretically it can be negative.

abbreviation $die-1 :: real \Rightarrow real$ ($\$q'- [101] 200$)
where $\$q-x \equiv \$q-\{1\&x\}$

definition $die-defer :: real \Rightarrow real \Rightarrow real \Rightarrow real$ ($\$q'-\{-|\&- \} [0,0,0] 200$)
where $\$q-\{f|t\&x\} = |\$q-\{f+t\&x\} - \$q-\{f\&x\}|$
— the probability that a person aged x will die within t years, deferred f years
— The parameters f and t are usually nonnegative, but theoretically they can be negative.

abbreviation $die-defer-1 :: real \Rightarrow real \Rightarrow real$ ($\$q'-\{-|\&- \} [0,0] 200$)
where $\$q-\{f|\&x\} \equiv \$q-\{f|1\&x\}$

definition $life-expect :: real \Rightarrow real$ ($\$e'\circ'- [101] 200$)
where $\$e'\circ-x \equiv integral^L (\mathfrak{M} \downarrow alive\ x) (T\ x)$
— complete life expectation
— Note that $\$e'\circ-x$ is calculated as 0 when $nn-integral (\mathfrak{M} \downarrow alive\ x) (T\ x) = \infty$.

definition $temp-life-expect :: real \Rightarrow real \Rightarrow real$ ($\$e'\circ'-\{-:- \} [0,0] 200$)
where $\$e'\circ-\{x:n\} \equiv integral^L (\mathfrak{M} \downarrow alive\ x) (\lambda\xi. min (T\ x\ \xi)\ n)$
— temporary complete life expectation

definition $curt-life-expect :: real \Rightarrow real$ ($\$e'- [101] 200$)
where $\$e-x \equiv integral^L (\mathfrak{M} \downarrow alive\ x) (\lambda\xi. \lfloor T\ x\ \xi \rfloor)$
— curtate life expectation
— Note that $\$e-x$ is calculated as 0 when $nn-integral (\mathfrak{M} \downarrow alive\ x) \lfloor T\ x \rfloor = \infty$.

definition $temp-curt-life-expect :: real \Rightarrow real \Rightarrow real$ ($\$e'-\{-:- \} [0,0] 200$)
where $\$e-\{x:n\} \equiv integral^L (\mathfrak{M} \downarrow alive\ x) (\lambda\xi. \lfloor min (T\ x\ \xi)\ n \rfloor)$
— temporary curtate life expectation
— In the definition n can be a real number, but in practice n is usually a natural number.

5.1.4 Properties of Survival Function for $T(x)$

context

fixes $x::real$
assumes $x-lt-psi[simp]: x < \psi$

begin

lemma $PXx-pos[simp]: \mathcal{P}(\xi\ in\ \mathfrak{M}. X\ \xi > x) > 0$

proof —

have $\mathcal{P}(\xi\ in\ \mathfrak{M}. X\ \xi > x) = cdf\ (distr\ \mathfrak{M}\ borel\ X)\ x$
unfolding $alive-def$ **using** $MM-PS.cdf-distr-P$ **by** $simp$

also have $\dots > 0$

using $ccdfX-0-equiv\ distrX-RD.cdf-nonneg\ x-lt-psi$ **by** $(smt\ (verit)\ linorder-not-le)$
finally show $?thesis$.

qed

lemma *PTx-pos[simp]*: $\mathcal{P}(\xi \text{ in } \mathfrak{M}. T x \xi > 0) > 0$
apply (*rewrite alive-T[THEN sym]*)
unfolding *alive-def* **by** *simp*

interpretation *alivex-PS*: *prob-space* $\mathfrak{M} \downarrow \text{alive } x$
by (*rule MM-PS.cond-prob-space-correct, simp-all add: alive-def*)

interpretation *distrTx-RD*: *real-distribution* *distr* ($\mathfrak{M} \downarrow \text{alive } x$) *borel* ($T x$) **by** *simp*

lemma *ccdfTx-cond-prob*:
 $\text{ccdf} (\text{distr } (\mathfrak{M} \downarrow \text{alive } x) \text{ borel } (T x)) t = \mathcal{P}(\xi \text{ in } \mathfrak{M}. T x \xi > t \mid T x \xi > 0)$ **for**
 $t :: \text{real}$
apply (*rewrite alivex-PS.ccdf-distr-P, simp*)
unfolding *alive-def*
apply (*rewrite MM-PS.cond-prob-space-cond-prob[THEN sym], simp-all add: pred-def*)
unfolding *futr-life-def* **by** *simp*

lemma *ccdfTx-0-1*: $\text{ccdf} (\text{distr } (\mathfrak{M} \downarrow \text{alive } x) \text{ borel } (T x)) 0 = 1$
apply (*rewrite ccdfTx-cond-prob*)
unfolding *futr-life-def cond-prob-def*
by (*smt (verit, best) Collect-cong PXX-pos divide-eq-1-iff*)

lemma *ccdfTx-nonpos-1*: $\text{ccdf} (\text{distr } (\mathfrak{M} \downarrow \text{alive } x) \text{ borel } (T x)) t = 1$ **if** $t \leq 0$ **for**
 $t :: \text{real}$
using *antisym ccdfTx-0-1 that*
by (*metis distrTx-RD.ccdf-bounded-prob distrTx-RD.ccdf-nonincreasing*)

lemma *ccdfTx-0-equiv*: $\text{ccdf} (\text{distr } (\mathfrak{M} \downarrow \text{alive } x) \text{ borel } (T x)) t = 0 \iff x+t \geq$
 $\$ \psi$ **for** $t :: \text{real}$

proof –

have $\text{ccdf} (\text{distr } (\mathfrak{M} \downarrow \text{alive } x) \text{ borel } (T x)) t =$
 $\mathcal{P}(\xi \text{ in } \mathfrak{M}. X \xi > x+t \wedge X \xi > x) / \mathcal{P}(\xi \text{ in } \mathfrak{M}. X \xi > x)$
apply (*rewrite ccdfTx-cond-prob*)
unfolding *cond-prob-def futr-life-def* **by** (*smt (verit) Collect-cong*)
hence $\text{ccdf} (\text{distr } (\mathfrak{M} \downarrow \text{alive } x) \text{ borel } (T x)) t = 0 \iff$
 $\mathcal{P}(\xi \text{ in } \mathfrak{M}. X \xi > x+t \wedge X \xi > x) / \mathcal{P}(\xi \text{ in } \mathfrak{M}. X \xi > x) = 0$
by *simp*
also have $\dots \iff \mathcal{P}(\xi \text{ in } \mathfrak{M}. X \xi > x+t \wedge X \xi > x) = 0$
using *x-lt-psi PXX-pos* **by** (*smt (verit) divide-eq-0-iff*)
also have $\dots \iff x+t \geq \$ \psi$
using *ccdfX-0-equiv MM-PS.ccdf-distr-P*
by (*smt (verit) Collect-cong X-RV le-ereal-le linorder-not-le x-lt-psi*)
finally show *?thesis* .

qed

lemma *ccdfTx-continuous-on-nonpos[simp]*:
 $\text{continuous-on } \{..0\} (\text{ccdf} (\text{distr } (\mathfrak{M} \downarrow \text{alive } x) \text{ borel } (T x)))$

by (*metis atMost-iff ccdfTx-nonpos-1 continuous-on-cong continuous-on-const*)

lemma *ccdfTx-differentiable-on-nonpos*[*simp*]:

(*ccdf (distr (M | alive x) borel (T x)) differentiable-on {..0}*)

by (*rewrite differentiable-on-cong[where f= λ -. 1]*; *simp add: ccdfTx-nonpos-1*)

lemma *ccdfTx-has-real-derivative-0-at-neg*:

(*ccdf (distr (M | alive x) borel (T x)) has-real-derivative 0*) (at *t*) **if** *t* < 0 **for** *t::real*

apply (*rewrite has-real-derivative-iff-has-vector-derivative*)

apply (*rule has-vector-derivative-transform-within-open[of λ -. 1 - - {..<0}]*)

using *ccdfTx-nonpos-1* **that** **by** *simp-all*

lemma *ccdfTx-integrable-Icc*:

set-integrable lborel {a..b} (*ccdf (distr (M | alive x) borel (T x))*) **for** *a b :: real*

proof –

have (\int^{+t} . *ennreal (indicat-real {a..b} t * ccdf (distr (M | alive x) borel (T x)) t) ∂ lborel*)

< \top

proof –

have (\int^{+t} . *ennreal (indicat-real {a..b} t * ccdf (distr (M | alive x) borel (T x)) t) ∂ lborel*)

\leq (\int^{+t} . *ennreal (indicat-real {a..b} t) ∂ lborel*)

apply (*rule nn-integral-mono*)

using *distrTx-RD.ccdf-bounded*

by (*simp add: distrTx-RD.ccdf-bounded-prob indicator-times-eq-if(1)*)

also have ... = *nn-integral lborel (indicator {a..b})* **by** (*meson ennreal-indicator*)

also have ... = *emeasure lborel {a..b}* **by** (*rewrite nn-integral-indicator; simp*)

also have ... < \top

using *emeasure-lborel-Icc-eq ennreal-less-top infinity-ennreal-def* **by** *presburger*
finally show *?thesis* .

qed

thus *?thesis*

unfolding *set-integrable-def*

apply (*intro integrableI-nonneg, simp-all*)

using *distrTx-RD.ccdf-nonneg* **by** (*intro always-eventually*) *auto*

qed

corollary *ccdfTx-integrable-on-Icc*:

ccdf (distr (M | alive x) borel (T x)) integrable-on {a..b} **for** *a b :: real*

using *set-borel-integral-eq-integral ccdfTx-integrable-Icc* **by** *force*

lemma *ccdfTx-PX*:

ccdf (distr (M | alive x) borel (T x)) t = $\mathcal{P}(\xi$ in \mathfrak{M} . $X \xi > x+t)$ / $\mathcal{P}(\xi$ in \mathfrak{M} . $X \xi > x)$

if *t* \geq 0 **for** *t::real*

apply (*rewrite ccdfTx-cond-prob*)

unfolding *cond-prob-def futr-life-def PXX-pos* **by** (*smt (verit) Collect-cong that*)

lemma *ccdfTx-ccdfX*: $ccdf (distr (\mathfrak{M} \downarrow \text{alive } x) \text{ borel } (T x)) t =$
 $ccdf (distr \mathfrak{M} \text{ borel } X) (x + t) / ccdf (distr \mathfrak{M} \text{ borel } X) x$
if $t \geq 0$ **for** $t::\text{real}$
using *ccdfTx-PX* **that** *MM-PS.ccdf-distr-P X-RV* **by** *presburger*

lemma *ccdfT0-eq-ccdfX*: $ccdf (distr (\mathfrak{M} \downarrow \text{alive } 0) \text{ borel } (T 0)) = ccdf (distr \mathfrak{M} \text{ borel } X)$
proof
fix x
show $ccdf (distr (\mathfrak{M} \downarrow \text{alive } 0) \text{ borel } (T 0)) x = ccdf (distr \mathfrak{M} \text{ borel } X) x$
proof (*cases* $\langle x \geq 0 \rangle$)
case *True*
thus *?thesis*
using *survival-model.ccdfTx-ccdfX[where x=0]* *ccdfX-0-1 survival-model-axioms*
by *fastforce*
next
case *False*
hence $x \leq 0$ **by** *simp*
thus *?thesis*
apply (*rewrite ccdfX-unborn-1, simp*)
by (*rewrite survival-model.ccdfTx-nonpos-1; simp add: survival-model-axioms*)
qed
qed

lemma *continuous-ccdfX-ccdfTx*:
 $continuous (at (x+t) \text{ within } \{x..\}) (ccdf (distr \mathfrak{M} \text{ borel } X)) \longleftrightarrow$
 $continuous (at t \text{ within } \{0..\}) (ccdf (distr (\mathfrak{M} \downarrow \text{alive } x) \text{ borel } (T x)))$
if $t \geq 0$ **for** $t::\text{real}$
proof –
let $?srvl = ccdf (distr \mathfrak{M} \text{ borel } X)$
have [*simp*]: $\mathcal{P}(\xi \text{ in } \mathfrak{M}. X \xi > x) \neq 0$ **using** *PXx-pos* **by** *force*
have $\star: \bigwedge u. u \geq 0 \implies ccdf (distr (\mathfrak{M} \downarrow \text{alive } x) \text{ borel } (T x)) u =$
 $?srvl (x + u) / \mathcal{P}(\xi \text{ in } \mathfrak{M}. X \xi > x)$
using *survive-def MM-PS.ccdf-distr-P ccdfTx-PX* **that** **by** *simp*
have $continuous (at t \text{ within } \{0..\}) (ccdf (distr (\mathfrak{M} \downarrow \text{alive } x) \text{ borel } (T x))) \longleftrightarrow$
 $continuous (at t \text{ within } \{0..\}) (\lambda u. ?srvl (x+u) / \mathcal{P}(\xi \text{ in } \mathfrak{M}. x < X \xi))$
proof –
have $\forall_F u \text{ in } at t \text{ within } \{0..\}. ccdf (distr (\mathfrak{M} \downarrow \text{alive } x) \text{ borel } (T x)) u =$
 $?srvl (x+u) / \mathcal{P}(\xi \text{ in } \mathfrak{M}. X \xi > x)$
using \star **by** (*rewrite eventually-at-topological, simp-all*) *blast*
thus *?thesis*
by (*intro continuous-at-within-cong, simp-all add: \star that*)
qed
also have $\dots \longleftrightarrow continuous (at t \text{ within } \{0..\}) (\lambda u. ?srvl (x+u))$
by (*rewrite at - = \sqcap continuous-cdivide-iff[of $\mathcal{P}(\xi \text{ in } \mathfrak{M}. X \xi > x)$], simp-all*)
also have $\dots \longleftrightarrow continuous (at (x+t) \text{ within } \{x..\}) ?srvl$
proof
let $?subx = \lambda v. v - x$
assume *LHS*: $continuous (at t \text{ within } \{0..\}) (\lambda u. ?srvl (x+u))$

hence *continuous* (at (*?subx* ($x+t$)) within *?subx* ‘ $\{x..\}$ ’) ($\lambda u.$ *?srvl* ($x+u$))
proof –
have *?subx* ‘ $\{x..\}$ ’ = $\{0..\}$
by (*metis* (*no-types*, *lifting*) *add.commute add-uminus-conv-diff diff-self image-add-atLeast image-cong*)
thus *?thesis* **using** *LHS* **by** *simp*
qed
moreover **have** *continuous* (at ($x+t$) within $\{x..\}$) *?subx* **by** (*simp add: continuous-diff*)
ultimately **have** *continuous* (at ($x+t$) within $\{x..\}$) ($\lambda u.$ *?srvl* ($x + (?subx u)$))
using *continuous-within-compose2* **by** *force*
thus *continuous* (at ($x+t$) within $\{x..\}$) *?srvl* **by** *simp*
next
assume *RHS: continuous* (at ($x+t$) within $\{x..\}$) *?srvl*
hence *continuous* (at (*plus* x) t) within (*plus* x) ‘ $\{0..\}$ ’) *?srvl* **by** *simp*
moreover **have** *continuous* (at t within $\{0..\}$) (*plus* x) **by** (*simp add: continuous-add*)
ultimately **show** *continuous* (at t within $\{0..\}$) ($\lambda u.$ *?srvl* ($x+u$))
using *continuous-within-compose2* **by** *force*
qed
finally **show** *?thesis* **by** *simp*
qed

lemma *isCont-ccdfX-ccdfTx*:

isCont (*ccdf* (*distr* \mathfrak{M} *borel* X)) ($x+t$) \longleftrightarrow
isCont (*ccdf* (*distr* ($\mathfrak{M} \downarrow$ *alive* x) *borel* ($T x$))) t
if $t > 0$ **for** $t::\text{real}$

proof –

have *isCont* (*ccdf* (*distr* \mathfrak{M} *borel* X)) ($x+t$) \longleftrightarrow
continuous (at ($x+t$) within $\{x<..\}$) (*ccdf* (*distr* \mathfrak{M} *borel* X))
by (*smt* (*verit*) *at-within-open greaterThan-iff open-greaterThan that*)
also **have** $\dots \longleftrightarrow$ *continuous* (at ($x+t$) within $\{x..\}$) (*ccdf* (*distr* \mathfrak{M} *borel* X))
by (*meson Ioi-le-Ico calculation continuous-within-subset top-greatest*)
also **have** $\dots \longleftrightarrow$ *continuous* (at t within $\{0..\}$) (*ccdf* (*distr* ($\mathfrak{M} \downarrow$ *alive* x) *borel* ($T x$)))
using *that continuous-ccdfX-ccdfTx* **by** *force*
also **have** $\dots \longleftrightarrow$ *continuous* (at t within $\{0<..\}$) (*ccdf* (*distr* ($\mathfrak{M} \downarrow$ *alive* x) *borel* ($T x$)))
by (*metis Ioi-le-Ico at-within-open continuous-at-imp-continuous-at-within continuous-within-subset greaterThan-iff open-greaterThan that*)
also **have** $\dots \longleftrightarrow$ *isCont* (*ccdf* (*distr* ($\mathfrak{M} \downarrow$ *alive* x) *borel* ($T x$))) t
by (*metis at-within-open greaterThan-iff open-greaterThan that*)
finally **show** *?thesis* .
qed

lemma *has-real-derivative-ccdfX-ccdfTx*:

((*ccdf* (*distr* \mathfrak{M} *borel* X)) *has-real-derivative* D) (at ($x+t$)) \longleftrightarrow
((*ccdf* (*distr* ($\mathfrak{M} \downarrow$ *alive* x) *borel* ($T x$))) *has-real-derivative* ($D / \mathcal{P}(\xi$ in $\mathfrak{M}. X \xi$

$> x$)) (at t)
if $t > 0$ **for** $t D :: \text{real}$
proof –
have ((ccdf (distr $\mathfrak{M} \downarrow \text{alive } x$) borel (T x))) has-real-derivative
 $(D / \mathcal{P}(\xi \text{ in } \mathfrak{M}. X \xi > x))$ (at t) \longleftrightarrow
 $((\lambda t. (\text{ccdf (distr } \mathfrak{M} \text{ borel } X)) (x+t) / \mathcal{P}(\xi \text{ in } \mathfrak{M}. X \xi > x))$ has-real-derivative
 $(D / \mathcal{P}(\xi \text{ in } \mathfrak{M}. X \xi > x))$ (at t)
proof –
let $?d = t/2$
{ **fix** $u :: \text{real}$ **assume** $\text{dist } u \ t < ?d$
hence $u > 0$ **by** (smt (verit) dist-real-def dist-triangle-half-r)
hence $\text{ccdf (distr } \mathfrak{M} \downarrow \text{alive } x \text{) borel (T } x \text{) } u =$
 $\text{ccdf (distr } \mathfrak{M} \text{ borel } X \text{) (} x + u \text{) / MM-PS.prob } \{\xi :: 'a \in \text{space } \mathfrak{M}. x < X \xi\}$
using survive-def MM-PS.ccdf-distr-P ccdfTx-PX **that** **by** simp }
moreover **have** $?d > 0$ **using** *that* **by** simp
ultimately **show** *?thesis*
apply –
apply (rule DERIV-cong-ev, simp)
apply (rewrite eventually-nhds-metric, blast)
by simp
qed
also **have** $\dots \longleftrightarrow ((\lambda t. (\text{ccdf (distr } \mathfrak{M} \text{ borel } X)) (x+t))$ has-real-derivative D)
(at t)
using PXx-pos **by** (rewrite DERIV-cdivide-iff[of $\mathcal{P}(\xi \text{ in } \mathfrak{M}. X \xi > x)$, THEN
sym]; force)
also **have** $\dots \longleftrightarrow (\text{ccdf (distr } \mathfrak{M} \text{ borel } X \text{) has-real-derivative } D)$ (at $(x+t)$)
by (simp add: DERIV-shift add commute)
finally **show** *?thesis* **by** simp
qed

lemma differentiable-ccdfX-ccdfTx:
 $(\text{ccdf (distr } \mathfrak{M} \text{ borel } X))$ differentiable at $(x+t) \longleftrightarrow$
 $(\text{ccdf (distr } (\mathfrak{M} \downarrow \text{alive } x) \text{ borel (T } x)))$ differentiable at t
if $t > 0$ **for** $t :: \text{real}$
apply (rewrite differentiable-eq-field-differentiable-real)+
unfolding field-differentiable-def **using** has-real-derivative-ccdfX-ccdfTx **that**
by (smt (verit, del-insts) PXx-pos nonzero-mult-div-cancel-left)

5.1.5 Properties of $\$p\{-t\&x\}$

lemma $p\text{-}0\text{-}1$: $\$p\{-0\&x\} = 1$
unfolding survive-def **using** ccdfTx-0-1 **by** simp

lemma $p\text{-nonneg}$ [simp]: $\$p\{-t\&x\} \geq 0$ **for** $t :: \text{real}$
unfolding survive-def **using** distrTx-RD.ccdf-nonneg **by** simp

lemma $p\text{-le-}1$ [simp]: $\$p\{-t\&x\} \leq 1$ **for** $t :: \text{real}$
unfolding survive-def **using** distrTx-RD.ccdf-bounded-prob **by** auto

lemma *p-0-equiv*: $\$p\{-t\&x\} = 0 \iff x+t \geq \ψ **for** $t::real$
unfolding *survive-def* **by** (rule *ccdfTx-0-equiv*)

lemma *p-PTx*: $\$p\{-t\&x\} = \mathcal{P}(\xi \text{ in } \mathfrak{M}. T x \xi > t \mid T x \xi > 0)$ **for** $t::real$
unfolding *survive-def* **using** *ccdfTx-cond-prob* **by** *simp*

lemma *p-PX*: $\$p\{-t\&x\} = \mathcal{P}(\xi \text{ in } \mathfrak{M}. X \xi > x+t) / \mathcal{P}(\xi \text{ in } \mathfrak{M}. X \xi > x)$ **if** $t \geq 0$ **for** $t::real$
unfolding *survive-def* **using** *ccdfTx-PX* **that** **by** *simp*

lemma *p-mult*: $\$p\{-t+t' \& x\} = \$p\{-t\&x\} * \$p\{-t' \& x+t\}$
if $t \geq 0$ $t' \geq 0$ $x+t < \$\psi$ **for** $t t' :: real$

proof –

have $\$p\{-t+t' \& x\} = \mathcal{P}(\xi \text{ in } \mathfrak{M}. X \xi > x+t+t') / \mathcal{P}(\xi \text{ in } \mathfrak{M}. X \xi > x)$

apply (rewrite *p-PX*; (*simp add: that*)?)

by (rule *disjI2*, *smt (verit, best) Collect-cong*)

also have $\dots = (\mathcal{P}(\xi \text{ in } \mathfrak{M}. X \xi > x+t+t') / \mathcal{P}(\xi \text{ in } \mathfrak{M}. X \xi > x+t)) * (\mathcal{P}(\xi \text{ in } \mathfrak{M}. X \xi > x+t) / \mathcal{P}(\xi \text{ in } \mathfrak{M}. X \xi > x))$

using *that survival-model.PXx-pos survival-model-axioms* **by** *fastforce*

also have $\dots = \$p\{-t\&x\} * \$p\{-t' \& x+t\}$

apply (rewrite *p-PX*, *simp add: that*)

by (rewrite *survival-model.p-PX*, *simp-all add: that survival-model-axioms*)

finally show *?thesis* .

qed

lemma *p-PTx-ge-ccdf-isCont*: $\$p\{-t\&x\} = \mathcal{P}(\xi \text{ in } \mathfrak{M}. T x \xi \geq t \mid T x \xi > 0)$
if *isCont (ccdf (distr \mathfrak{M} borel X)) (x+t) t > 0* **for** $t::real$
unfolding *survive-def* **using** *that isCont-ccdfX-ccdfTx*
apply (rewrite *aliveX-PS.ccdf-continuous-distr-P-ge*, *simp-all*)
by (rewrite *MM-PS.cond-prob-space-cond-prob*, *simp-all add: alive-T*)

end

5.1.6 Properties of Survival Function for X

lemma *ccdfX-continuous-unborn[*simp*]*: *continuous-on* $\{..0\}$ (*ccdf (distr \mathfrak{M} borel X)*)
using *ccdfTx-continuous-on-nonpos* **by** (*metis ccdfT0-eq-ccdfX psi-pos'*)

lemma *ccdfX-differentiable-unborn[*simp*]*: (*ccdf (distr \mathfrak{M} borel X)*) *differentiable-on* $\{..0\}$
using *ccdfTx-differentiable-on-nonpos* **by** (*metis ccdfT0-eq-ccdfX psi-pos'*)

lemma *ccdfX-has-real-derivative-0-unborn*:
(*ccdf (distr \mathfrak{M} borel X) has-real-derivative 0*) (at x) **if** $x < 0$ **for** $x::real$
using *ccdfTx-has-real-derivative-0-at-neg* **by** (*metis ccdfT0-eq-ccdfX psi-pos' that*)

lemma *ccdfX-integrable-Icc*:
set-integrable lborel $\{a..b\}$ (*ccdf (distr \mathfrak{M} borel X)*) **for** $a b :: real$

using *ccdfTx-integrable-Icc* **by** (*metis ccdfT0-eq-ccdfX psi-pos'*)

corollary *ccdfX-integrable-on-Icc*:

ccdf (distr \mathfrak{M} borel X) integrable-on {a..b} for a b :: real
using *set-borel-integral-eq-integral ccdfX-integrable-Icc* **by force**

lemma *ccdfX-p*: *ccdf (distr \mathfrak{M} borel X) x = $\mathbb{P}\{x \leq 0\}$ for $x :: \text{real}$*
by (*metis ccdfT0-eq-ccdfX survive-def psi-pos'*)

5.1.7 Introduction of Cumulative Distributive Function for X

lemma *cdfX-0-0*: *cdf (distr \mathfrak{M} borel X) 0 = 0*
using *ccdfX-0-1 distrX-RD.cdf-cdf distrX-RD.prob-space* **by fastforce**

lemma *cdfX-unborn-0*: *cdf (distr \mathfrak{M} borel X) x = 0 if $x \leq 0$*
using *ccdfX-unborn-1 cdfX-0-0 distrX-RD.cdf-ccdf* **that by fastforce**

lemma *cdfX-beyond-1*: *cdf (distr \mathfrak{M} borel X) x = 1 if $x > \psi$ for $x :: \text{real}$*
using *ccdfX-beyond-0 distrX-RD.cdf-ccdf* **that distrX-RD.prob-space by force**

lemma *cdfX-psi-1*: *cdf (distr \mathfrak{M} borel X) (real-of-ereal ψ) = 1 if $\psi < \infty$*
using *ccdfX-psi-0 distrX-RD.cdf-ccdf distrX-RD.prob-space* **that by fastforce**

lemma *cdfX-1-equiv*: *cdf (distr \mathfrak{M} borel X) x = 1 \longleftrightarrow $x \geq \psi$ for $x :: \text{real}$*
using *ccdfX-0-equiv distrX-RD.cdf-ccdf distrX-RD.prob-space* **by force**

5.1.8 Properties of Cumulative Distributive Function for $T(x)$

context

fixes $x :: \text{real}$
assumes $x \text{ lt } \psi$ [*simp*]: $x < \psi$

begin

interpretation *aliveX-PS*: *prob-space $\mathfrak{M} \mid \text{alive } x$*
by (*rule MM-PS.cond-prob-space-correct, simp-all add: alive-def*)

interpretation *distrTx-RD*: *real-distribution distr ($\mathfrak{M} \mid \text{alive } x$) borel (T x) by simp*

lemma *cdfTx-cond-prob*:

cdf (distr ($\mathfrak{M} \mid \text{alive } x$) borel (T x)) t = $\mathcal{P}(\xi \text{ in } \mathfrak{M}. T x \xi \leq t \mid T x \xi > 0)$ for $t :: \text{real}$

apply (*rewrite distrTx-RD.cdf-ccdf, rewrite distrTx-RD.prob-space*)

apply (*rewrite ccdfTx-cond-prob, simp*)

by (*rewrite not-less[THEN sym], rewrite MM-PS.cond-prob-neg; simp*)

lemma *cdfTx-0-0*: *cdf (distr ($\mathfrak{M} \mid \text{alive } x$) borel (T x)) 0 = 0*
using *ccdfTx-0-1 distrTx-RD.cdf-ccdf distrTx-RD.prob-space* **by force**

lemma *cdfTx-nonpos-0*: $\text{cdf} (\text{distr } (\mathfrak{M} \mid \text{alive } x) \text{ borel } (T x)) t = 0$ if $t \leq 0$ for $t :: \text{real}$

using *ccdfTx-nonpos-1 distrTx-RD.cdf-ccdf distrTx-RD.prob-space* that **by force**

lemma *cdfTx-1-equiv*: $\text{cdf} (\text{distr } (\mathfrak{M} \mid \text{alive } x) \text{ borel } (T x)) t = 1 \iff x+t \geq \psi$ for $t::\text{real}$

using *ccdfTx-0-equiv distrTx-RD.cdf-ccdf distrTx-RD.prob-space* **by force**

lemma *cdfTx-continuous-on-nonpos*[simp]:

continuous-on $\{..0\}$ ($\text{cdf} (\text{distr } (\mathfrak{M} \mid \text{alive } x) \text{ borel } (T x))$)

by (*rewrite continuous-on-cong*[**where** $g=\lambda t. 0$]) (*simp-all add: cdfTx-nonpos-0*)+

lemma *cdfTx-differentiable-on-nonpos*[simp]:

($\text{cdf} (\text{distr } (\mathfrak{M} \mid \text{alive } x) \text{ borel } (T x))$) *differentiable-on* $\{..0\}$

by (*rewrite differentiable-on-cong*[**where** $f=\lambda t. 0$]; *simp add: cdfTx-nonpos-0*)

lemma *cdfTx-has-real-derivative-0-at-neg*:

($\text{cdf} (\text{distr } (\mathfrak{M} \mid \text{alive } x) \text{ borel } (T x))$) *has-real-derivative* 0 (at t) if $t < 0$ for $t::\text{real}$

apply (*rewrite has-real-derivative-iff-has-vector-derivative*)

apply (*rule has-vector-derivative-transform-within-open*[of $\lambda-. 0 - - \{..<0\}$])

using *cdfTx-nonpos-0* that **by simp-all**

lemma *cdfTx-integrable-Icc*:

set-integrable *lborel* $\{a..b\}$ ($\text{cdf} (\text{distr } (\mathfrak{M} \mid \text{alive } x) \text{ borel } (T x))$) for $a b :: \text{real}$

proof –

have *set-integrable lborel* $\{a..b\}$ ($\lambda-. 1::\text{real}$)

unfolding *set-integrable-def*

using *emeasure-compact-finite* **by** (*simp, intro integrable-real-indicator; force*)

thus *?thesis*

apply (*rewrite distrTx-RD.cdf-ccdf, rewrite distrTx-RD.prob-space*)

using *ccdfTx-integrable-Icc* **by** (*rewrite set-integral-diff; simp*)

qed

corollary *cdfTx-integrable-on-Icc*:

$\text{cdf} (\text{distr } (\mathfrak{M} \mid \text{alive } x) \text{ borel } (T x))$ *integrable-on* $\{a..b\}$ for $a b :: \text{real}$

using *cdfTx-integrable-Icc set-borel-integral-eq-integral* **by force**

lemma *cdfTx-PX*:

$\text{cdf} (\text{distr } (\mathfrak{M} \mid \text{alive } x) \text{ borel } (T x)) t = \mathcal{P}(\xi \text{ in } \mathfrak{M}. x < X \xi \wedge X \xi \leq x+t) / \mathcal{P}(\xi \text{ in } \mathfrak{M}. X \xi > x)$

for $t::\text{real}$

apply (*rewrite cdfTx-cond-prob*)

unfolding *cond-prob-def futr-life-def PXX-pos* **by** (*smt (verit) Collect-cong*)

lemma *cdfT0-eq-cdfX*: $\text{cdf} (\text{distr } (\mathfrak{M} \mid \text{alive } 0) \text{ borel } (T 0)) = \text{cdf} (\text{distr } \mathfrak{M} \text{ borel } X)$

proof

interpret *alive0-PS*: *prob-space* $\mathfrak{M} \mid \text{alive } 0$

apply (rule *MM-PS.cond-prob-space-correct*, *simp*)
using *PXx-pos alive-def psi-pos'* **by** *presburger*
interpret *distrT0-RD: real-distribution distr* ($\mathfrak{M} \mid \text{alive } 0$) *borel* ($T\ 0$) **by** *simp*
show $\bigwedge x. \text{cdf} (\text{distr } (\mathfrak{M} \mid \text{alive } 0) \text{ borel } (T\ 0))\ x = \text{cdf} (\text{distr } \mathfrak{M} \text{ borel } X)\ x$
using *ccdfT0-eq-ccdfX distrX-RD.ccdf-cdf distrT0-RD.ccdf-cdf*
by (*smt (verit, best) distrT0-RD.prob-space distrX-RD.prob-space psi-pos'*)
qed

lemma *continuous-cdfX-cdfTx:*

$\text{continuous} (\text{at } (x+t) \text{ within } \{x..\}) (\text{cdf} (\text{distr } \mathfrak{M} \text{ borel } X)) \longleftrightarrow$
 $\text{continuous} (\text{at } t \text{ within } \{0..\}) (\text{cdf} (\text{distr } (\mathfrak{M} \mid \text{alive } x) \text{ borel } (T\ x)))$
if $t \geq 0$ **for** $t::\text{real}$

proof –

have $\text{continuous} (\text{at } (x+t) \text{ within } \{x..\}) (\text{cdf} (\text{distr } \mathfrak{M} \text{ borel } X)) \longleftrightarrow$
 $\text{continuous} (\text{at } (x+t) \text{ within } \{x..\}) (\text{ccdf} (\text{distr } \mathfrak{M} \text{ borel } X))$
by (rule *distrX-RD.continuous-cdf-ccdf*)
also have $\dots \longleftrightarrow \text{continuous} (\text{at } t \text{ within } \{0..\}) (\text{ccdf} (\text{distr } (\mathfrak{M} \mid \text{alive } x) \text{ borel } (T\ x)))$
using *continuous-ccdfX-ccdfTx that* **by** *simp*
also have $\dots \longleftrightarrow \text{continuous} (\text{at } t \text{ within } \{0..\}) (\text{cdf} (\text{distr } (\mathfrak{M} \mid \text{alive } x) \text{ borel } (T\ x)))$
using *distrTx-RD.continuous-cdf-ccdf* **by** *simp*
finally show *?thesis* .
qed

lemma *isCont-cdfX-cdfTx:*

$\text{isCont} (\text{cdf} (\text{distr } \mathfrak{M} \text{ borel } X))\ (x+t) \longleftrightarrow$
 $\text{isCont} (\text{cdf} (\text{distr } (\mathfrak{M} \mid \text{alive } x) \text{ borel } (T\ x)))\ t$
if $t > 0$ **for** $t::\text{real}$
apply (rewrite *distrX-RD.isCont-cdf-ccdf*)
apply (rewrite *isCont-ccdfX-ccdfTx, simp-all add: that*)
by (rule *distrTx-RD.isCont-cdf-ccdf[THEN sym]*)

lemma *has-real-derivative-cdfX-cdfTx:*

$((\text{cdf} (\text{distr } \mathfrak{M} \text{ borel } X)) \text{ has-real-derivative } D) (\text{at } (x+t)) \longleftrightarrow$
 $((\text{cdf} (\text{distr } (\mathfrak{M} \mid \text{alive } x) \text{ borel } (T\ x))) \text{ has-real-derivative } (D / \mathcal{P}(\xi \text{ in } \mathfrak{M}. X\ \xi > x))) (\text{at } t)$
if $t > 0$ **for** $t\ D :: \text{real}$
proof –
have $((\text{cdf} (\text{distr } \mathfrak{M} \text{ borel } X)) \text{ has-real-derivative } D) (\text{at } (x+t)) \longleftrightarrow$
 $(\text{ccdf} (\text{distr } \mathfrak{M} \text{ borel } X) \text{ has-real-derivative } -D) (\text{at } (x+t))$
using *distrX-RD.has-real-derivative-cdf-ccdf* **by** *force*
also have $\dots \longleftrightarrow$
 $((\text{ccdf} (\text{distr } (\mathfrak{M} \mid \text{alive } x) \text{ borel } (T\ x))) \text{ has-real-derivative } (-D / \mathcal{P}(\xi \text{ in } \mathfrak{M}. X\ \xi > x))) (\text{at } t)$
using *has-real-derivative-ccdfX-ccdfTx that* **by** *simp*
also have $\dots \longleftrightarrow$
 $((\text{cdf} (\text{distr } (\mathfrak{M} \mid \text{alive } x) \text{ borel } (T\ x))) \text{ has-real-derivative } (D / \mathcal{P}(\xi \text{ in } \mathfrak{M}. X\ \xi > x))) (\text{at } t)$

by (simp add: distrTx-RD.has-real-derivative-cdf-ccdf)
 finally show ?thesis .
 qed

lemma differentiable-cdfX-cdfTx:
 (cdf (distr \mathfrak{M} borel X)) differentiable at (x+t) \longleftrightarrow
 (cdf (distr ($\mathfrak{M} \setminus \text{alive } x$) borel (T x))) differentiable at t
 if t > 0 for t::real
apply (rewrite differentiable-eq-field-differentiable-real)+
unfolding field-differentiable-def **using** has-real-derivative-cdfX-cdfTx that
by (meson differentiable-ccdfX-ccdfTx distrTx-RD.finite-borel-measure-axioms
 distrX-RD.finite-borel-measure-axioms finite-borel-measure.differentiable-cdf-ccdf
 real-differentiable-def x-lt-psi)

5.1.9 Properties of $q\{t\&x\}$

lemma q-nonpos-0: $q\{t\&x\} = 0$ if $t \leq 0$ for t::real
unfolding die-def **using** that cdfTx-nonpos-0 **by** simp

corollary q-0-0: $q\{0\&x\} = 0$
using q-nonpos-0 **by** simp

lemma q-nonneg[simp]: $q\{t\&x\} \geq 0$ for t::real
unfolding die-def **using** distrTx-RD.cdf-nonneg **by** simp

lemma q-le-1[simp]: $q\{t\&x\} \leq 1$ for t::real
unfolding die-def **using** distrTx-RD.cdf-bounded-prob **by** force

lemma q-1-equiv: $q\{t\&x\} = 1 \longleftrightarrow x+t \geq \psi$ for t::real
unfolding die-def **using** cdfTx-1-equiv **by** simp

lemma q-PTx: $q\{t\&x\} = \mathcal{P}(\xi \text{ in } \mathfrak{M}. T x \xi \leq t \mid T x \xi > 0)$ for t::real
unfolding die-def **using** cdfTx-cond-prob **by** simp

lemma q-PX: $q\{t\&x\} = \mathcal{P}(\xi \text{ in } \mathfrak{M}. x < X \xi \wedge X \xi \leq x + t) / \mathcal{P}(\xi \text{ in } \mathfrak{M}. X \xi > x)$
unfolding die-def **using** cdfTx-PX **by** simp

lemma q-defer-0-q[simp]: $q\{0\mid t\&x\} = q\{t\&x\}$ for t::real
unfolding die-defer-def **using** q-0-0 **by** simp

lemma q-defer-0-0: $q\{f\mid 0\&x\} = 0$ for f::real
unfolding die-defer-def **by** simp

lemma q-defer-nonneg[simp]: $q\{f\mid t\&x\} \geq 0$ for f t :: real
unfolding die-defer-def **by** simp

lemma q-defer-q: $q\{f\mid t\&x\} = q\{f+t \& x\} - q\{f\&x\}$ if $t \geq 0$ for f t :: real
unfolding die-defer-def die-def **using** distrTx-RD.cdf-nondecreasing that **by**

simp

corollary $q\text{-defer-le-1}[simp]: \$q\{-f|t\&x\} \leq 1$ **if** $t \geq 0$ **for** $f\ t :: \text{real}$
by (*smt (verit, ccfv-SIG) q-defer-q q-le-1 q-nonneg that*)

lemma $q\text{-defer-PTx}: \$q\{-f|t\&x\} = \mathcal{P}(\xi \text{ in } \mathfrak{M}. f < T\ x\ \xi \wedge T\ x\ \xi \leq f + t \mid T\ x\ \xi > 0)$

if $t \geq 0$ **for** $f\ t :: \text{real}$

proof –

have $\$q\{-f|t\&x\} = \$q\{-f+t\ \&\ x\} - \$q\{-f\&x\}$ **using** $q\text{-defer-q}$ **that** **by** *simp*

also have $\dots = \mathcal{P}(\xi \text{ in } \mathfrak{M}. T\ x\ \xi \leq f + t \mid T\ x\ \xi > 0) - \mathcal{P}(\xi \text{ in } \mathfrak{M}. T\ x\ \xi \leq f \mid T\ x\ \xi > 0)$

using $q\text{-PTx}$ **by** *simp*

also have $\dots = \mathcal{P}(\xi \text{ in } (\mathfrak{M} \mid \text{alive } x). T\ x\ \xi \leq f + t) - \mathcal{P}(\xi \text{ in } (\mathfrak{M} \mid \text{alive } x). T\ x\ \xi \leq f)$

using $MM\text{-PS.cond-prob-space-cond-prob alive-T}$ **by** *simp*

also have $\dots = \mathcal{P}(\xi \text{ in } (\mathfrak{M} \mid \text{alive } x). f < T\ x\ \xi \wedge T\ x\ \xi \leq f + t)$

proof –

have $\{\xi \in \text{space } (\mathfrak{M} \mid \text{alive } x). T\ x\ \xi \leq f + t\} - \{\xi \in \text{space } (\mathfrak{M} \mid \text{alive } x). T\ x\ \xi \leq f\} =$

$\{\xi \in \text{space } (\mathfrak{M} \mid \text{alive } x). f < T\ x\ \xi \wedge T\ x\ \xi \leq f + t\}$

using *that by force*

hence alive-x-PS.prob

$(\{\xi \in \text{space } (\mathfrak{M} \mid \text{alive } x). T\ x\ \xi \leq f + t\} - \{\xi \in \text{space } (\mathfrak{M} \mid \text{alive } x). T\ x\ \xi \leq f\}) =$

$\mathcal{P}(\xi \text{ in } (\mathfrak{M} \mid \text{alive } x). f < T\ x\ \xi \wedge T\ x\ \xi \leq f + t)$

by *simp*

moreover have $\{\xi \in \text{space } (\mathfrak{M} \mid \text{alive } x). T\ x\ \xi \leq f\} \subseteq \{\xi \in \text{space } (\mathfrak{M} \mid \text{alive } x). T\ x\ \xi \leq f + t\}$

using *that by force*

ultimately show *?thesis by (rewrite alive-x-PS.finite-measure-Diff[THEN sym]; simp)*

qed

also have $\dots = \mathcal{P}(\xi \text{ in } \mathfrak{M}. f < T\ x\ \xi \wedge T\ x\ \xi \leq f + t \mid T\ x\ \xi > 0)$

using $MM\text{-PS.cond-prob-space-cond-prob alive-T}$ **by** *simp*

finally show *?thesis .*

qed

lemma $q\text{-defer-PX}: \$q\{-f|t\&x\} = \mathcal{P}(\xi \text{ in } \mathfrak{M}. x + f < X\ \xi \wedge X\ \xi \leq x + f + t) / \mathcal{P}(\xi \text{ in } \mathfrak{M}. X\ \xi > x)$

if $f \geq 0\ t \geq 0$ **for** $f\ t :: \text{real}$

proof –

have $\$q\{-f|t\&x\} = \mathcal{P}(\xi \text{ in } \mathfrak{M}. f < T\ x\ \xi \wedge T\ x\ \xi \leq f + t \wedge T\ x\ \xi > 0) / \mathcal{P}(\xi \text{ in } \mathfrak{M}. T\ x\ \xi > 0)$

apply (*rewrite q-defer-PTx; (simp add: that)?*)

unfolding cond-prob-def **by** *simp*

also have $\dots = \mathcal{P}(\xi \text{ in } \mathfrak{M}. f < T\ x\ \xi \wedge T\ x\ \xi \leq f + t) / \mathcal{P}(\xi \text{ in } \mathfrak{M}. T\ x\ \xi > 0)$

proof –

have $\bigwedge \xi. \xi \in \text{space } \mathfrak{M} \implies f < T x \xi \wedge T x \xi \leq f + t \wedge T x \xi > 0 \iff f < T x \xi \wedge T x \xi \leq f + t$
using *that by auto*
hence $\{\xi \in \text{space } \mathfrak{M}. f < T x \xi \wedge T x \xi \leq f + t \wedge T x \xi > 0\} = \{\xi \in \text{space } \mathfrak{M}. f < T x \xi \wedge T x \xi \leq f + t\}$ **by** *blast*
thus *?thesis by simp*
qed
also have $\dots = \mathcal{P}(\xi \text{ in } \mathfrak{M}. x + f < X \xi \wedge X \xi \leq x + f + t) / \mathcal{P}(\xi \text{ in } \mathfrak{M}. X \xi > x)$
unfolding *futr-life-def by (smt (verit) Collect-cong)*
finally show *?thesis .*
qed

lemma *q-defer-old-0*: $\$q\{-f|t\&x\} = 0$ **if** $x+f \geq \psi$ $t \geq 0$ **for** $f t :: \text{real}$

proof –

have $\$q\{-f|t\&x\} = \$q\{-f+t \& x\} - \$q\{-f\&x\}$ **using** *q-defer-q that by simp*
moreover have $\$q\{-f+t \& x\} = 1$ **using** *q-1-equiv that le-ereal-le by auto*
moreover have $\$q\{-f\&x\} = 1$ **using** *q-1-equiv that by simp*
ultimately show *?thesis by simp*

qed

end

5.1.10 Properties of Cumulative Distributive Function for X

lemma *cdfX-continuous-unborn[simp]*: *continuous-on* $\{..0\}$ (*cdf (distr* \mathfrak{M} *borel* X) **using** *cdfTx-continuous-on-nonpos by (metis cdfT0-eq-cdfX psi-pos')*)

lemma *cdfX-differentiable-unborn[simp]*: (*cdf (distr* \mathfrak{M} *borel* X) *differentiable-on* $\{..0\}$ **using** *cdfTx-differentiable-on-nonpos by (metis cdfT0-eq-cdfX psi-pos')*)

lemma *cdfX-has-real-derivative-0-unborn*:
(cdf (distr \mathfrak{M} *borel* X) *has-real-derivative* 0) (at x) **if** $x < 0$ **for** $x :: \text{real}$
using *cdfTx-has-real-derivative-0-at-neg by (metis cdfT0-eq-cdfX psi-pos' that)*

lemma *cdfX-integrable-Icc*:
set-integrable lborel $\{a..b\}$ (*cdf (distr* \mathfrak{M} *borel* X) **for** $a b :: \text{real}$
using *cdfTx-integrable-Icc by (metis cdfT0-eq-cdfX psi-pos')*)

corollary *cdfX-integrable-on-Icc*:
cdf (distr \mathfrak{M} *borel* X) *integrable-on* $\{a..b\}$ **for** $a b :: \text{real}$
using *cdfX-integrable-Icc set-borel-integral-eq-integral by force*

lemma *cdfX-q*: *cdf (distr* \mathfrak{M} *borel* X) $x = \$q\{-x\&0\}$ **if** $x \geq 0$ **for** $x :: \text{real}$
by (*metis cdfT0-eq-cdfX die-def psi-pos')*

5.1.11 Relations between $\$p\{-t\&x\}$ and $\$q\{-t\&x\}$

context

fixes $x::real$
assumes $x\text{-lt-psi}[simp]: x < \psi$
begin

interpretation $alive\text{-}PS$: *prob-space* $\mathfrak{M} \mid \text{alive } x$
by (*rule* $MM\text{-}PS.\text{cond-prob-space-correct}$, *simp-all add: alive-def*)

interpretation $distrTx\text{-}RD$: *real-distribution* $distr (\mathfrak{M} \mid \text{alive } x)$ *borel* $(T x)$ **by**
 $simp$

lemma $p\text{-}q\text{-}1$: $\$p\text{-}\{t\&x\} + \$q\text{-}\{t\&x\} = 1$ **for** $t::real$
unfolding $survive\text{-}def$ $die\text{-}def$ **using** $distrTx\text{-}RD.add\text{-}cdf\text{-}ccdf$
by (*smt* (*verit*) $distrTx\text{-}RD.\text{prob-space } x\text{-lt-psi}$)

lemma $q\text{-}defer\text{-}p$: $\$q\text{-}\{f|t\&x\} = \$p\text{-}\{f\&x\} - \$p\text{-}\{f+t \& x\}$ **if** $t \geq 0$ **for** $f t :: real$
using $q\text{-}defer\text{-}q$ $p\text{-}q\text{-}1$ *that* $x\text{-lt-psi}$ **by** *smt*

lemma $q\text{-}defer\text{-}p\text{-}q\text{-}defer$: $\$p\text{-}\{f\&x\} * \$q\text{-}\{f'|t \& x+f\} = \$q\text{-}\{f+f'|t \& x\}$
if $x+f < \psi$ $f \geq 0$ $f' \geq 0$ $t \geq 0$ **for** $f f' t :: real$

proof –

have $\$p\text{-}\{f\&x\} * \$q\text{-}\{f'|t \& x+f\} =$
 $\mathcal{P}(\xi \text{ in } \mathfrak{M}. X \xi > x+f) / \mathcal{P}(\xi \text{ in } \mathfrak{M}. X \xi > x) *$
 $\mathcal{P}(\xi \text{ in } \mathfrak{M}. x+f+f' < X \xi \wedge X \xi \leq x+f+f'+t) / \mathcal{P}(\xi \text{ in } \mathfrak{M}. X \xi > x+f)$
apply (*rewrite* $p\text{-}PX$, (*simp-all add: that*)[2])
by (*rewrite* $survival\text{-}model.q\text{-}defer\text{-}PX$, *simp-all add: that survival-model-axioms*)
also have $\dots = \mathcal{P}(\xi \text{ in } \mathfrak{M}. x+f+f' < X \xi \wedge X \xi \leq x+f+f'+t) / \mathcal{P}(\xi \text{ in } \mathfrak{M}. X$
 $\xi > x)$
using $survival\text{-}model.PXx\text{-}pos$ [*of* $\mathfrak{M} X x+f$] *nonzero-mult-div-cancel-left that*
by (*smt* (*verit*, $ccfv\text{-}SIG$) *survival-model-axioms times-divide-eq-left times-divide-eq-right*)
also have $\dots = \$q\text{-}\{f+f'|t \& x\}$
by (*rewrite* $q\text{-}defer\text{-}PX$; *simp add: that group-cancel.add1*)
finally show *?thesis* .

qed

lemma $q\text{-}defer\text{-}pq$: $\$q\text{-}\{f|t\&x\} = \$p\text{-}\{f\&x\} * \$q\text{-}\{t \& x+f\}$
if $x+f < \psi$ $t \geq 0$ $f \geq 0$ **for** $f t :: real$
using $q\text{-}defer\text{-}p\text{-}q\text{-}defer$ [**where** $f'=0$] *that*
by (*simp add: survival-model.q-defer-0-q survival-model-axioms*)

5.1.12 Properties of Life Expectation

lemma $e\text{-nonneg}$: $\$e'\circ\text{-}x \geq 0$
unfolding $life\text{-}expect\text{-}def$
by (*rule* $Bochner\text{-}Integration.\text{integral-nonneg}$, *simp add: less-eq-real-def*)

lemma $e\text{-}P$: $\$e'\circ\text{-}x =$
 $MM\text{-}PS.\text{expectation } (\lambda\xi. \text{indicator } (\text{alive } x) \xi * T x \xi) / \mathcal{P}(\xi \text{ in } \mathfrak{M}. T x \xi > 0)$
unfolding $life\text{-}expect\text{-}def$
by (*rewrite* $MM\text{-}PS.\text{integral-cond-prob-space-nn}$, *auto simp add: alive-T*)

proposition *nn-integral-T-p*:

$(\int^{+\xi}. \text{ennreal } (T x \xi) \partial(\mathfrak{M} \downarrow \text{alive } x)) = (\int^{+t \in \{0..\}}. \text{ennreal } (\$p\text{-}\{t\&x\}) \partial\text{lborel})$
apply (*rewrite alive-x-PS.expectation-nonneg-tail, simp-all add: less-imp-le*)
apply (*rule nn-integral-cong*)
unfolding *survive-def* **using** *distrTx-RD.prob-space distrTx-RD.ccdf-cdf* **by** *presburger*

lemma *nn-integral-T-pos*: $(\int^{+\xi}. \text{ennreal } (T x \xi) \partial(\mathfrak{M} \downarrow \text{alive } x)) > 0$

proof –

let $?f = \lambda t. - \text{ccdf } (\text{distr } (\mathfrak{M} \downarrow \text{alive } x) \text{ borel } (T x)) t$
have $\wedge t u. t \leq u \implies ?f t \leq ?f u$ **using** *distrTx-RD.ccdf-nonincreasing* **by** *simp*
moreover **have** *continuous (at-right 0)* $?f$
using *distrTx-RD.ccdf-is-right-cont* **by** (*intro continuous-intros*)
ultimately **have** $\forall e > 0. \exists d > 0. ?f (0 + d) - ?f 0 < e$
using *continuous-at-right-real-increasing* **by** *simp*
hence $\exists d > 0. ?f (0 + d) - ?f 0 < 1/2$ **by** (*smt (verit, del-Insts) field-sum-of-halves*)
from this **obtain** d **where** *d-pos: d > 0* **and** $\$p\text{-}\{d\&x\} \geq 1/2$
using *p-0-1* **unfolding** *survive-def* **by** *auto*
hence $\wedge t. t \in \{0..d\} \implies \$p\text{-}\{t\&x\} \geq 1/2$
unfolding *survive-def* **using** *distrTx-RD.ccdf-nonincreasing* **by** *force*
hence $(\int^{+t \in \{0..d\}}. \text{ennreal } (\$p\text{-}\{t\&x\}) \partial\text{lborel}) \geq (\int^{+t \in \{0..d\}}. \text{ennreal } (1/2) \partial\text{lborel})$
apply (*intro nn-set-integral-mono, simp-all*)
unfolding *survive-def* **using** *Tx-alive-x-measurable* **apply** *force*
by (*rule AE-I2*) (*smt (verit) ennreal-half ennreal-leI half-bounded-equal*)
moreover **have** $(\int^{+t \in \{0..\}}. \text{ennreal } (\$p\text{-}\{t\&x\}) \partial\text{lborel}) \geq (\int^{+t \in \{0..d\}}. \text{ennreal } (\$p\text{-}\{t\&x\}) \partial\text{lborel})$
by (*rule nn-set-integral-set-mono*) *simp*
moreover **have** $(\int^{+t \in \{0..d\}}. \text{ennreal } (1/2) \partial\text{lborel}) > 0$
apply (*rewrite nn-integral-cmult-indicator, simp-all*)
using *d-pos emeasure-lborel-Icc ennreal-zero-less-mult-iff* **by** *fastforce*
ultimately **show** *?thesis* **using** *nn-integral-T-p* **by** *simp*
qed

lemma *e-pos-Tx*: $\$e\text{'o-}x > 0$ **if** *integrable* $(\mathfrak{M} \downarrow \text{alive } x) (T x)$

unfolding *life-expect-def*
apply (*rewrite integral-eq-nn-integral, simp-all*)
apply (*smt (verit, ccfv-SIG) AE-I2 alive-x-Tx-pos*)
using *nn-integral-T-pos* **that**
by (*smt (verit) AE-I2 alive-x-Tx-pos enn2real-ennreal ennreal-less-zero-iff nn-integral-cong nn-integral-eq-integral*)

proposition *e-LBINT-p*: $\$e\text{'o-}x = (\text{LBINT } t:\{0..\}. \$p\text{-}\{t\&x\})$

– Note that $0 = 0$ holds when the integral diverges.

unfolding *life-expect-def* **apply** (*rewrite integral-eq-nn-integral, simp-all add: less-imp-le*)
unfolding *set-lebesgue-integral-def* **apply** (*rewrite integral-eq-nn-integral, simp-all*)
apply (*measurable, simp add: survive-def*)

by (rewrite nn-integral-T-p) (simp add: indicator-mult-ennreal mult.commute)

corollary *e-integral-p*: $\$e^{\circ}x = \text{integral } \{0..\} (\lambda t. \$p\{t\&x\})$

— Note that $0 = 0$ holds when the integral diverges.

proof –

have $\$e^{\circ}x = (\text{LBINT } t:\{0..\}. \$p\{t\&x\})$ using *e-LBINT-p* by *simp*

also have $\dots = \text{integral } \{0..\} (\lambda t. \$p\{t\&x\})$

apply (rule *set-borel-integral-eq-integral-nonneg*, *simp-all*)

unfolding *survive-def* by *simp*

finally show *?thesis* .

qed

lemma *e-pos*: $\$e^{\circ}x > 0$ if *set-integrable lborel* $\{0..\} (\lambda t. \$p\{t\&x\})$

proof –

have $(\int^{+} t \in \{0..\}. \text{ennreal } (\$p\{t\&x\}) \partial \text{lborel}) = \text{ennreal } (\int t \in \{0..\}. \$p\{t\&x\} \partial \text{lborel})$

by (*intro set-nn-integral-eq-set-integral*; *simp add: that*)

also have $\dots < \infty$ using *that* by *simp*

finally have $(\int^{+} \xi. \text{ennreal } (T x \xi) \partial (\mathfrak{M} \downarrow \text{alive } x)) < \infty$ using *nn-integral-T-p* by *simp*

hence *integrable* $(\mathfrak{M} \downarrow \text{alive } x) (T x)$

by (*smt (verit) alive-x-Tx-pos integrableI-bounded nn-integral-cong real-norm-def survival-model.Tx-alive-x-measurable survival-model-axioms*)

thus *?thesis* by (rule *e-pos-Tx*)

qed

corollary *e-pos'*: $\$e^{\circ}x > 0$ if $(\lambda t. \$p\{t\&x\})$ *integrable-on* $\{0..\}$

apply (rule *e-pos*)

using *that* apply (rewrite *integrable-on-iff-set-integrable-nonneg*; *simp*)

unfolding *survive-def* by *simp*

lemma *e-LBINT-p-Icc*: $\$e^{\circ}x = (\text{LBINT } t:\{0..n\}. \$p\{t\&x\})$ if $x+n \geq \$\psi$ for $n::\text{real}$

proof –

have [*simp*]: $\{0..n\} \cap \{n<..\} = \{\}$ using *ivl-disj-int-one(7)* by *blast*

have [*simp*]: $\{0..n\} \cup \{n<..\} = \{0..\}$

by (*smt (verit) ereal-less-le ivl-disj-un-one(7) leD that x-lt-psi*)

have [*simp*]: $\bigwedge t. n < t \implies 0 \leq t$ using *that x-lt-psi* by (*smt (verit) ereal-less-le leD*)

have [*simp*]: $\bigwedge t. n < t \implies \$\psi \leq \text{ereal } (x+t)$ using *that* by (*simp add: le-ereal-le*)

have *gt-n-0*: *has-bochner-integral lborel* $(\lambda t. \text{indicat-real } \{n<..\} t * \$p\{t\&x\}) 0$

apply (rewrite *has-bochner-integral-cong*[**where** $N=\text{lborel}$ **and** $g=\lambda t.0$ **and** $y=0$], *simp-all*)

using *p-0-equiv that x-lt-psi*

apply (*smt (verit, ccfv-SIG) greaterThan-iff indicator-simps le-ereal-le linorder-not-le*)

by (rule *has-bochner-integral-zero*)

hence *gt-n*: *set-integrable lborel* $\{n<..\} (\lambda t. \$p\{t\&x\})$

unfolding *set-integrable-def* using *integrable.simps* by *auto*

moreover have *le-n*: *set-integrable lborel* $\{0..n\} (\lambda t. \$p\{t\&x\})$

unfolding *survive-def* **by** (*intro cdfTx-integrable-Icc*) *simp*
ultimately have *set-integrable lborel* ($\{0..n\} \cup \{n<..\}$) ($\lambda t. \$p-\{t\&x\}$)
using *set-integrable-Un* **by** *force*
hence *set-integrable lborel* $\{0..\}$ ($\lambda t. \$p-\{t\&x\}$) **by** *force*
thus *?thesis*
apply (*rewrite e-LBINT-p, simp*)
apply (*rewrite set-integral-Un*[of $\{0..n\} \{n<..\}$, *simplified*], *simp-all add: gt-n le-n*)
unfolding *set-lebesgue-integral-def* **using** *gt-n-0 has-bochner-integral-integral-eq*
by *fastforce*
qed

lemma *e-integral-p-Icc*: $\$e^{\circ-x} = \text{integral } \{0..n\}$ ($\lambda t. \$p-\{t\&x\}$) **if** $x+n \geq \$\psi$ **for** $n::\text{real}$

using *that* **apply** (*rewrite e-LBINT-p-Icc, simp-all*)
using *cdfTx-integrable-Icc* **unfolding** *survive-def*
by (*rewrite set-borel-integral-eq-integral; simp*)

lemma *temp-e-le-n*: $\$e^{\circ-\{x:n\}} \leq n$ **if** $n \geq 0$ **for** $n::\text{real}$

proof –

have *nni-n*: $(\int^{+}. \text{ennreal } n \partial(\mathfrak{M} \downarrow \text{alive } x)) = \text{ennreal } n$
by (*rewrite nn-integral-const, rewrite alivex-PS.emeasure-space-1*) *simp*
hence *hbi-n*: *has-bochner-integral* $(\mathfrak{M} \downarrow \text{alive } x)$ $(\lambda-. n)$ n
by (*intro has-bochner-integral-nn-integral; simp add: that*)
hence *integrable* $(\mathfrak{M} \downarrow \text{alive } x)$ $(\lambda-. n)$ **by** *simp*
moreover have *integrable* $(\mathfrak{M} \downarrow \text{alive } x)$ $(\lambda\xi. \min (T x \xi) n)$

proof –

have $(\int^{+\xi}. \text{ennreal } (\text{norm } (\min (T x \xi) n)) \partial(\mathfrak{M} \downarrow \text{alive } x)) \leq \int^{+}. \text{ennreal } n \partial(\mathfrak{M} \downarrow \text{alive } x)$

apply (*rule nn-integral-mono, rule ennreal-leI*)

apply (*rewrite real-norm-def, rewrite abs-of-nonneg; simp add: that*)

by (*smt (verit) alivex-Tx-pos*)

also have $\dots < \infty$ **using** *nni-n* **by** *simp*

finally have $(\int^{+\xi}. \text{ennreal } (\text{norm } (\min (T x \xi) n)) \partial(\mathfrak{M} \downarrow \text{alive } x)) < \infty$.

thus *?thesis* **by** (*intro integrableI-bounded; simp*)

qed

ultimately have $\$e^{\circ-\{x:n\}} \leq \text{integral}^L (\mathfrak{M} \downarrow \text{alive } x)$ $(\lambda-. n)$

unfolding *temp-life-expect-def* **by** (*intro integral-mono; simp*)

also have $\dots = n$ **using** *hbi-n has-bochner-integral-iff* **by** *blast*

finally show *?thesis* .

qed

lemma *temp-e-P*: $\$e^{\circ-\{x:n\}} =$

$MM\text{-PS.expectation } (\lambda\xi. \text{indicator } (\text{alive } x) \xi * \min (T x \xi) n) / \mathcal{P}(\xi \text{ in } \mathfrak{M}. T x \xi > 0)$

if $n \geq 0$ **for** $n::\text{real}$

unfolding *temp-life-expect-def*

by (*rewrite MM-PS.integral-cond-prob-space-nn; simp add: alive-T that*)

lemma *temp-e-LBINT-p*: $\$e^{\circ}\{-x:n\} = (LBINT\ t:\{0..n\}. \$p\{-t&x\})$ if $n \geq 0$ for $n::real$

proof –

let $?minTxn = \lambda\xi. \min (T\ x\ \xi)\ n$

let $?F = cdf\ (distr\ (\mathfrak{M}\ \downarrow\ alive\ x)\ borel\ (T\ x))$

let $?Fn = cdf\ (distr\ (\mathfrak{M}\ \downarrow\ alive\ x)\ borel\ ?minTxn)$

interpret *distrTxn-RD*: *real-distribution* $distr\ (\mathfrak{M}\ \downarrow\ alive\ x)\ borel\ ?minTxn$ **by** (*simp* *add*: *that*)

have [*simp*]: $\bigwedge\xi. \xi \in space\ (\mathfrak{M}\ \downarrow\ alive\ x) \implies 0 \leq T\ x\ \xi$ **by** (*smt* (*verit*) *alivex-Tx-pos*)

have $(\int^{+\xi}. ennreal\ (\min\ (T\ x\ \xi)\ n)\ \partial(\mathfrak{M}\ \downarrow\ alive\ x)) = (\int^{+t \in \{0..n\}}. ennreal\ (1 - ?Fn\ t)\ \partial lborel)$

by (*rewrite* *alivex-PS.expectation-nonneg-tail*; *simp* *add*: *that*)

also have $\dots = (\int^{+t \in \{0..n\}}. (ennreal\ (1 - ?F\ t) * indicator\ \{..<n\}\ t)\ \partial lborel)$

apply (*rule* *nn-integral-cong*)

by (*rewrite* *alivex-PS.cdf-distr-min*; *simp* *add*: *indicator-mult-ennreal\ mult.\ commute*)

also have $\dots = (\int^{+t \in \{0..<n\}}. ennreal\ (1 - ?F\ t)\ \partial lborel)$

apply (*rule* *nn-integral-cong*) **using** *nn-integral-set-ennreal*

by (*smt* (*verit*, *best*) *Int-def* *atLeastLessThan-def* *ennreal-mult-right-cong* *indicator-simps* *mem-Collect-eq\ mult.\ commute\ mult-1*)

also have $\dots = (\int^{+t \in \{0..n\}}. ennreal\ (1 - ?F\ t)\ \partial lborel)$

proof –

have *sym-diff* $\{0..<n\}\ \{0..n\} = \{n\}$ **using** *that* **by** *force*

thus *?thesis* **by** (*intro* *nn-integral-null-delta*; *force*)

qed

also have $\dots = ennreal\ (LBINT\ t:\{0..n\}. \$p\{-t&x\})$

proof –

have *set-integrable* *lborel* $\{0..n\}$ $(\lambda t. \$p\{-t&x\})$

unfolding *survive-def* **by** (*intro* *ccdfTx-integrable-Icc*) *simp*

thus *?thesis*

unfolding *set-lebesgue-integral-def* **unfolding** *set-integrable-def*

apply (*rewrite* *nn-integral-eq-integral*[*THEN* *sym*]; *simp*)

apply (*rule* *nn-integral-cong*, *simp*)

unfolding *survive-def* **using** *distrTx-RD.ccdf-cdf* *distrTx-RD.prob-space* *nn-integral-set-ennreal*

by (*simp* *add*: *indicator-mult-ennreal\ mult.\ commute*)

qed

finally have $(\int^{+\xi}. ennreal\ (\min\ (T\ x\ \xi)\ n)\ \partial(\mathfrak{M}\ \downarrow\ alive\ x)) = ennreal\ (LBINT\ t:\{0..n\}. \$p\{-t&x\})$.

thus *?thesis*

unfolding *temp-life-expect-def* **by** (*rewrite* *integral-eq-nn-integral*; *simp* *add*: *that*)

qed

lemma *temp-e-integral-p*: $\$e^{\circ}\{-x:n\} = integral\ \{0..n\}\ (\lambda t. \$p\{-t&x\})$ if $n \geq 0$ for $n::real$

using *that* **apply** (*rewrite* *temp-e-LBINT-p*, *simp-all*)

using *ccdfTx-integrable-Icc* **unfolding** *survive-def*

by (*rewrite* *set-borel-integral-eq-integral*; *simp*)

lemma *e-eq-temp*: $\$e^{\circ}x = \$e^{\circ}\{x:n\}$ if $n \geq 0$ $x+n \geq \$\psi$ for $n::real$
 using that *e-LBINT-p-Icc temp-e-LBINT-p* by *simp*

lemma *curt-e-P*: $\$e-x =$
 $MM-PS.expectation (\lambda\xi. indicator (alive x) \xi * \lfloor T x \xi \rfloor) / \mathcal{P}(\xi \text{ in } \mathfrak{M}. T x \xi >$
 $0)$
unfolding *curt-life-expect-def*
apply (*rewrite MM-PS.integral-cond-prob-space-nn; simp add: alive-T*)
by (*metis (no-types, lifting) Bochner-Integration.integral-cong indicator-simps*
of-int-0 of-int-1)

lemma *curt-e-sum-P*: $\$e-x = (\sum k. \mathcal{P}(\xi \text{ in } \mathfrak{M}. T x \xi \geq k + 1 \mid T x \xi > 0))$
 if *summable* $(\lambda k. \mathcal{P}(\xi \text{ in } \mathfrak{M}. T x \xi \geq k + 1 \mid T x \xi > 0))$

proof –

let $?F-flrTx = cdf (distr (\mathfrak{M} \downarrow alive x) borel (\lambda\xi. \lfloor T x \xi \rfloor))$
have [*simp*]: $\bigwedge\xi. \xi \in space (\mathfrak{M} \downarrow alive x) \implies 0 \leq T x \xi$ by (*smt (verit)*
alivex-Tx-pos)

have $integral^N (\mathfrak{M} \downarrow alive x) (\lambda\xi. ennreal \lfloor T x \xi \rfloor) =$
 $(\int^{+t \in \{0..\}. ennreal (1 - ?F-flrTx t) \partial lborel}$

by (*rewrite alivex-PS.expectation-nonneg-tail; simp*)

also have $\dots = (\int^{+t \in \{0::real..\}. ennreal \mathcal{P}(\xi \text{ in } \mathfrak{M}. T x \xi \geq \lfloor t \rfloor + 1 \mid T x \xi$
 $> 0) \partial lborel}$

proof –

{ **fix** $t::real$ **assume** $t \geq 0$

hence $1 - ?F-flrTx t = \mathcal{P}(\xi \text{ in } \mathfrak{M}. T x \xi \geq real-of-int \lfloor t \rfloor + 1 \mid T x \xi > 0)$

proof –

have $1 - ?F-flrTx t = 1 - \mathcal{P}(\xi \text{ in } (\mathfrak{M} \downarrow alive x). T x \xi < real-of-int \lfloor t \rfloor$
 $+ 1)$

by (*rewrite alivex-PS.cdf-distr-floor-P; simp*)

also have $\dots = 1 - \mathcal{P}(\xi \text{ in } \mathfrak{M}. T x \xi < real-of-int \lfloor t \rfloor + 1 \mid T x \xi > 0)$

using *alive-T* by (*rewrite MM-PS.cond-prob-space-cond-prob; simp*)

also have $\dots = \mathcal{P}(\xi \text{ in } \mathfrak{M}. T x \xi \geq real-of-int \lfloor t \rfloor + 1 \mid T x \xi > 0)$

by (*rewrite not-le[THEN sym], rewrite MM-PS.cond-prob-neg; simp*)

finally show *?thesis* .

qed }

thus *?thesis*

apply –

by (*rule nn-set-integral-cong2, rule AE-I2*) *simp*

qed

also have $\dots = (\sum k. \int^{+t \in \{k..<k+1\}. ennreal \mathcal{P}(\xi \text{ in } \mathfrak{M}. T x \xi \geq \lfloor t \rfloor + 1 \mid$
 $T x \xi > 0) \partial lborel}$

apply (*rewrite nn-integral-disjoint-family[THEN sym]; simp*)

apply (*rewrite add.commute, rule Ico-nat-disjoint*)

by (*rewrite Ico-nat-union[THEN sym], simp add: add.commute*)

also have $\dots = (\sum k. \int^{+t \in \{k..<k+1::nat\}. ennreal \mathcal{P}(\xi \text{ in } \mathfrak{M}. T x \xi \geq k + 1$
 $\mid T x \xi > 0) \partial lborel}$

proof –

{ **fix** $k::nat$ **and** $t::real$

assume $real\ k \leq t$ **and** $t < 1 + real\ k$
hence $real-of-int\ \lfloor t \rfloor = real\ k$
by (*metis add.commute floor-eq2 of-int-of-nat-eq*)
hence $\mathcal{P}(\xi \text{ in } \mathfrak{M}. T\ x\ \xi \geq real-of-int\ \lfloor t \rfloor + 1 \mid T\ x\ \xi > 0) =$
 $\mathcal{P}(\xi \text{ in } \mathfrak{M}. T\ x\ \xi \geq 1 + real\ k \mid T\ x\ \xi > 0)$
by (*simp add: add.commute*) }
thus *?thesis*
apply –
apply (*rule suminf-cong, rule nn-set-integral-cong2, rule AE-I2*)
by (*rule impI*) *simp*
qed
also have $\dots = (\sum k. ennreal\ \mathcal{P}(\xi \text{ in } \mathfrak{M}. T\ x\ \xi \geq k + 1 \mid T\ x\ \xi > 0))$
by (*rewrite nn-integral-cmult-indicator; simp add: add.commute*)
also have $\dots = ennreal\ (\sum k. \mathcal{P}(\xi \text{ in } \mathfrak{M}. T\ x\ \xi \geq k + 1 \mid T\ x\ \xi > 0))$
by (*rewrite suminf-ennreal2; simp add: that*)
finally have $integral^N\ (\mathfrak{M} \mid alive\ x)\ (\lambda\xi. ennreal\ \lfloor T\ x\ \xi \rfloor) =$
 $ennreal\ (\sum k. \mathcal{P}(\xi \text{ in } \mathfrak{M}. T\ x\ \xi \geq k + 1 \mid T\ x\ \xi > 0))$.
hence $integral^L\ (\mathfrak{M} \mid alive\ x)\ (\lambda\xi. \lfloor T\ x\ \xi \rfloor) = (\sum k. \mathcal{P}(\xi \text{ in } \mathfrak{M}. T\ x\ \xi \geq k + 1$
 $\mid T\ x\ \xi > 0))$
apply (*rewrite integral-eq-nn-integral; simp*)
apply (*rewrite enn2real-ennreal; simp add: add.commute*)
apply (*rule suminf-nonneg; simp?*)
by (*rewrite add.commute, simp add: that*)
thus *?thesis unfolding curt-life-expect-def by (simp add: add.commute)*
qed

lemma *curt-e-sum-P-finite*: $\$e-x = (\sum k < n. \mathcal{P}(\xi \text{ in } \mathfrak{M}. T\ x\ \xi \geq k + 1 \mid T\ x\ \xi > 0))$
if $x+n+1 > \$\psi$ **for** $n::nat$
proof –
from *that have psi-fin: $\psi < \infty$ by force*
let $?P = \lambda k::nat. \mathcal{P}(\xi \text{ in } \mathfrak{M}. T\ x\ \xi \geq k + 1 \mid T\ x\ \xi > 0)$
let $?P-fin = \lambda k::nat. \text{if } k \in \{..<n\} \text{ then } ?P\ k \text{ else } 0$
have $\bigwedge k. ?P\ k = ?P-fin\ k$
proof –
fix k
show $?P\ k = ?P-fin\ k$
proof (*cases $\langle k \in \{..<n\} \rangle$*)
case *True*
thus *?thesis by simp*
next
case *False*
hence $\neg k < n$ **by** *simp*
hence $x + k + 1 > real-of-ereal\ \ψ
using *that psi-nonneg real-of-ereal-ord-simps(4) by fastforce*
hence $\{\xi \in space\ \mathfrak{M}. T\ x\ \xi \geq k + 1 \wedge T\ x\ \xi > 0\} \subseteq \{\xi \in space\ \mathfrak{M}. X\ \xi >$
 $real-of-ereal\ \$\psi\}$
unfolding *futr-life-def using that less-ereal-le of-nat-1 of-nat-add by force*
hence $\mathcal{P}(\xi \text{ in } \mathfrak{M}. T\ x\ \xi \geq k + 1 \wedge T\ x\ \xi > 0) \leq \mathcal{P}(\xi \text{ in } \mathfrak{M}. X\ \xi > real-of-ereal$

$\$ \psi$)
by (*intro MM-PS.finite-measure-mono, simp-all*)
also have $\dots = 0$ **using** *MM-PS.ccdf-distr-P X-RV ccdfX-psi-0 psi-fin* **by**
presburger
finally have $\mathcal{P}(\xi \text{ in } \mathfrak{M}. T x \xi \geq k + 1 \wedge T x \xi > 0) = 0$ **using** *measure-le-0-iff*
by *blast*
hence $?P k = 0$ **unfolding** *cond-prob-def* **by** (*simp add: add.commute*)
thus *?thesis* **by** *simp*
qed
qed
moreover have $?P\text{-fin sums } (\sum_{k < n}. ?P k)$ **using** *sums-If-finite-set* **by** *force*
ultimately have $\star: ?P \text{ sums } (\sum_{k < n}. ?P k)$ **using** *sums-cong* **by** *simp*
moreover hence *summable ?P* **using** *sums-summable* **by** *blast*
ultimately have $?P \text{ sums } \$e\text{-}x$ **using** *curt-e-sum-P* **by** *force*
hence $\$e\text{-}x = (\sum_{k < n}. ?P k)$ **by** (*rewrite sums-unique2[of ?P]; simp add: \star*)
thus *?thesis* **by** (*simp add: add.commute*)
qed

lemma *curt-e-sum-p*: $\$e\text{-}x = (\sum k. \$p\text{-}\{k+1\&x\})$
if *summable* $(\lambda k. \$p\text{-}\{k+1\&x\}) \wedge k::\text{nat. isCont } (\lambda t. \$p\text{-}\{t\&x\}) (k+1)$
proof –
have $\wedge k::\text{nat. } \$p\text{-}\{k+1\&x\} = \mathcal{P}(\xi \text{ in } \mathfrak{M}. T x \xi \geq k + 1 \mid T x \xi > 0)$
apply (*rewrite p-PTx-ge-ccdf-isCont, simp-all*)
using *that(2) isCont-ccdfX-ccdfTx* **unfolding** *survive-def* **by** *simp*
thus *?thesis* **using** *that p-PTx-ge-ccdf-isCont curt-e-sum-P* **by** *presburger*
qed

lemma *curt-e-rec*: $\$e\text{-}x = \$p\text{-}x * (1 + \$e\text{-}(x+1))$
if *summable* $(\lambda k. \$p\text{-}\{k+1\&x\}) \wedge k::\text{nat. isCont } (\lambda t. \$p\text{-}\{t\&x\}) (\text{real } k + 1) x+1 < \$ \psi$
proof –
have *px-neq-0[simp]*: $\$p\text{-}x \neq 0$ **using** *p-0-equiv that* **by** *auto*
have $(\lambda k. \$p\text{-}\{k+1\&x\}) \text{ sums } \$e\text{-}x$
using *that* **apply** (*rewrite curt-e-sum-p, simp-all add: add.commute*)
by (*rule summable-sums, simp add: that*)
hence $(\lambda k. \$p\text{-}x * \$p\text{-}\{k\&x+1\}) \text{ sums } \$e\text{-}x$
apply (*rewrite sums-cong[where g= $\lambda k. \$p\text{-}\{k+1\&x\}$]; simp?*)
using *p-mult* **by** (*smt (verit) of-nat-0-le-iff that(3) x-lt-psi*)
hence $(\lambda k. \$p\text{-}\{k\&x+1\}) \text{ sums } (\$e\text{-}x / \$p\text{-}x)$
using *sums-mult-D that* **by** (*smt (verit, best) linorder-not-le p-0-equiv sums-cong x-lt-psi*)
hence *p-e-p*: $(\lambda k. \$p\text{-}\{Suc k \& x+1\}) \text{ sums } (\$e\text{-}x / \$p\text{-}x - \$p\text{-}\{0\&x+1\})$
using *sums-split-initial-segment[where n=1]* **by** *force*
moreover have $(\lambda k. \$p\text{-}\{Suc k \& x+1\}) \text{ sums } \$e\text{-}(x+1)$
proof –
have [*simp*]: *summable* $(\lambda k::\text{nat. } \$p\text{-}\{\text{real } k + 1 \& x + 1\})$
apply (*intro sums-summable[where l= $\$e\text{-}x / \$p\text{-}x - \$p\text{-}\{0\&x+1\}$]*)
using *p-e-p* **by** (*simp add: add.commute*)
have [*simp*]: $\wedge k::\text{nat. isCont } (\lambda t. \$p\text{-}\{t\&x+1\}) (\text{real } k + 1)$

proof –
fix $k::nat$
have $isCont (\lambda t. \$p-x * \$p-\{t-1\&x+1\}) (real\ k + 2)$
proof –
let $?S=\{real\ k + 1 <..< real\ k + 3\}$
have open $?S$ **by** $simp$
moreover have $real\ k + 2 \in ?S$ **by** $simp$
moreover have $\bigwedge t. t \in ?S \implies \$p-x * \$p-\{t-1\&x+1\} = \$p-\{t\&x\}$
using $p-mult$
by ($smt\ (verit, del-insts)\ greaterThanLessThan-iff\ of-nat-0-le-iff\ that(3)$
 $x-lt-psi$)
ultimately show $?thesis$
apply ($rewrite\ isCont-cong[\mathbf{where}\ g=\lambda t. \$p-\{t\&x\}]$)
apply ($rewrite\ eventually-nhds, blast$)
using that by ($smt\ (verit)\ of-nat-1\ of-nat-add$)
qed
hence $isCont (\lambda t. \$p-x * \$p-\{t-1\&x+1\} / \$p-x) (real\ k + 2)$
by ($intro\ isCont-divide[\mathbf{where}\ g=\lambda t. \$p-x], auto$)
hence $isCont ((\lambda t. \$p-\{t-1\&x+1\}) \circ (\lambda t. t+1)) (real\ k + 1)$
by $simp\ (rule\ continuous-at-compose, simp-all\ add: add.commute)$
thus $isCont (\lambda t. \$p-\{t\&x+1\}) (real\ k + 1)$ **unfolding comp-def by** $simp$
qed
show $?thesis$
apply ($rewrite\ survival-model.curt-e-sum-p; simp\ add: survival-model-axioms$
 $that$)
using $summable-sums$ **by** ($rewrite\ add.commute$) $force$
qed
ultimately have $\$e-x / \$p-x - \$p-\{0\&x+1\} = \$e-(x+1)$ **by** ($rule\ sums-unique2$)
thus $?thesis$
using $p-0-1$ $that$
by ($smt\ (verit)\ px-neq-0\ divide-mult-cancel\ mult.commute\ mult-cancel-left2$
 $p-mult\ that(3)$)
qed

lemma $curt-e-sum-p-finite: \$e-x = (\sum k < n. \$p-\{k+1\&x\})$
if $\bigwedge k::nat. k < n \implies isCont (\lambda t. \$p-\{t\&x\}) (real\ k + 1)$ $x+n+1 > \$\psi$ **for**
 $n::nat$
proof –
have $\bigwedge k::nat. k < n \implies \$p-\{k+1\&x\} = \mathcal{P}(\xi\ in\ \mathfrak{M}. T\ x\ \xi \geq k + 1 \mid T\ x\ \xi >$
 $0)$
apply ($rewrite\ p-PTx-ge-ccdf-isCont, simp-all$)
using that $isCont-ccdfX-ccdfTx$ **unfolding survive-def by** ($smt\ (verit)\ of-nat-0-le-iff$
 $x-lt-psi$)
thus $?thesis$ **using that** $p-PTx-ge-ccdf-isCont\ curt-e-sum-P-finite$ **by** $simp$
qed

lemma $temp-curt-e-P: \$e-\{x:n\} =$
 $MM-PS.expectation (\lambda \xi. indicator\ (alive\ x)\ \xi * \lfloor min\ (T\ x\ \xi)\ n \rfloor) / \mathcal{P}(\xi\ in\ \mathfrak{M}. T$
 $x\ \xi > 0)$

```

if  $n \geq 0$  for  $n::\text{real}$ 
unfolding temp-curt-life-expect-def
apply (rewrite MM-PS.integral-cond-prob-space-nn; simp add: alive-T that)
apply (rule disjI2, rule Bochner-Integration.integral-cong; simp)
using indicator-simps of-int-0 of-int-1 by smt

lemma temp-curt-e-sum-P:  $\$e\{-x:n\} = (\sum_{k < n}. \mathcal{P}(\xi \text{ in } \mathfrak{M}. T x \xi \geq k + 1 \mid T x \xi > 0))$  for  $n::\text{nat}$ 
proof –
  let  $?F\text{-flrmin}Tx = \text{cdf} (\text{distr } (\mathfrak{M} \mid \text{alive } x) \text{ borel } (\lambda\xi. \lfloor \min (T x \xi) n \rfloor))$ 
  have [simp]:  $\bigwedge \xi. \xi \in \text{space } (\mathfrak{M} \mid \text{alive } x) \implies 0 \leq T x \xi$  by (smt (verit) alive-x-Tx-pos)
  have  $\text{integral}^N (\mathfrak{M} \mid \text{alive } x) (\lambda\xi. \text{ennreal } \lfloor \min (T x \xi) n \rfloor) =$ 
     $(\int^{+t \in \{0..\}. \text{ennreal } (1 - ?F\text{-flrmin}Tx t) \partial \text{lborel}}$ 
    by (rewrite alive-x-PS.expectation-nonneg-tail; simp)
  also have  $\dots = (\int^{+t \in \{0::\text{real}..\}. \text{ennreal}}$ 
     $((1 - \mathcal{P}(\xi \text{ in } (\mathfrak{M} \mid \text{alive } x). T x \xi < \lfloor t \rfloor + 1)) * \text{of-bool } (\lfloor t \rfloor + 1 \leq n)) \partial \text{lborel})$ 
  proof –
    have  $\bigwedge t. t \geq 0 \implies \text{ennreal } (1 - ?F\text{-flrmin}Tx t) =$ 
       $\text{ennreal } ((1 - \mathcal{P}(\xi \text{ in } (\mathfrak{M} \mid \text{alive } x). T x \xi < \lfloor t \rfloor + 1)) * \text{of-bool } (\lfloor t \rfloor + 1 \leq n))$ 
  proof –
    fix  $t::\text{real}$  assume  $t \geq 0$ 
    thus  $\text{ennreal } (1 - ?F\text{-flrmin}Tx t) =$ 
       $\text{ennreal } ((1 - \mathcal{P}(\xi \text{ in } (\mathfrak{M} \mid \text{alive } x). T x \xi < \lfloor t \rfloor + 1)) * \text{of-bool } (\lfloor t \rfloor + 1 \leq n))$ 
  proof (cases  $\langle \lfloor t \rfloor + 1 \leq n \rangle$ )
    case True
    thus ?thesis
    apply (rewrite alive-x-PS.cdf-distr-floor-P; simp)
    using min-less-iff-disj
    by (smt (verit, ccfv-SIG) Collect-cong
      floor-eq floor-less-cancel floor-of-nat of-int-floor-le)
  next
  case False
  thus ?thesis
  apply (rewrite alive-x-PS.cdf-distr-floor-P; simp)
  using min-less-iff-disj
  by (smt (verit, ccfv-SIG) Collect-cong MM-PS.space-cond-prob-space alive-T
alive-event
alive-x-PS.prob-space mem-Collect-eq of-int-of-nat-eq of-nat-less-of-int-iff)
  qed
qed
thus ?thesis
by (intro nn-set-integral-cong2, intro AE-I2) auto
qed
also have  $\dots = (\int^{+t \in \{0..<n\}. \text{ennreal } (1 - \mathcal{P}(\xi \text{ in } (\mathfrak{M} \mid \text{alive } x). T x \xi < \lfloor t \rfloor + 1)) \partial \text{lborel}}$ 
proof –

```

```

{ fix t::real
  have ( $\lfloor t \rfloor + 1 \leq n$ ) = ( $t < n$ ) by linarith
  hence  $\bigwedge r::real.$ 
    ennreal ( $r * \text{of-bool} (\lfloor t \rfloor + 1 \leq n)$ ) * indicator {0..} t = ennreal r * indicator
{0..<n} t
  unfolding atLeastLessThan-def by (simp add: indicator-def) }
  thus ?thesis by simp
qed
also have ... = ( $\int^{+t \in \{0..<n\}}.$  ennreal  $\mathcal{P}(\xi \text{ in } (\mathfrak{M} \downarrow \text{alive } x). T x \xi \geq \lfloor t \rfloor + 1)$   $\partial \text{lborel}$ )
  by (rewrite alive-PS.prob-neg[THEN sym]; simp add: not-less)
also have ... = ( $\sum k < n.$   $\int^{+t \in \{k..<k+1\}}.$  ennreal  $\mathcal{P}(\xi \text{ in } (\mathfrak{M} \downarrow \text{alive } x). T x \xi \geq \lfloor t \rfloor + 1)$   $\partial \text{lborel}$ )
  apply (rewrite Ico-nat-union-finite[of n, THEN sym])
  apply (rewrite nn-integral-disjoint-family-on-finite; simp add: add commute)
  apply (rule disjoint-family-on-mono[of - {0..}]; simp)
  by (rewrite add commute, rule Ico-nat-disjoint)
also have ... = ( $\sum k < n.$  ennreal  $\mathcal{P}(\xi \text{ in } \mathfrak{M}. T x \xi \geq k + 1 \mid T x \xi > 0)$ )
proof -
  { fix k::nat
    assume  $k < n$ 
    hence ( $\int^{+t \in \{k..<(1 + \text{real } k)\}}.$ 
      ennreal  $\mathcal{P}(\xi \text{ in } (\mathfrak{M} \downarrow \text{alive } x). T x \xi \geq \text{real-of-int } \lfloor t \rfloor + 1)$   $\partial \text{lborel}$ ) =
       $\mathcal{P}(\xi \text{ in } \mathfrak{M}. T x \xi \geq 1 + \text{real } k \mid T x \xi > 0)$  (is ?LHS = ?RHS)
    proof -
      have ?LHS = ( $\int^{+t \in \{k..<(1 + \text{real } k)\}}.$  ennreal  $\mathcal{P}(\xi \text{ in } (\mathfrak{M} \downarrow \text{alive } x). T x \xi \geq k + 1)$   $\partial \text{lborel}$ )
    proof -
      { fix t::real
        assume  $k \leq t < 1 + \text{real } k$ 
        hence  $\text{real-of-int } \lfloor t \rfloor = \text{real } k$  by (smt (verit) floor-eq2 of-int-of-nat-eq)
        hence  $\mathcal{P}(\xi \text{ in } (\mathfrak{M} \downarrow \text{alive } x). T x \xi \geq \text{real-of-int } \lfloor t \rfloor + 1) =$ 
           $\mathcal{P}(\xi \text{ in } (\mathfrak{M} \downarrow \text{alive } x). T x \xi \geq 1 + \text{real } k)$ 
        by (simp add: add commute) }
      thus ?thesis by (intro nn-set-integral-cong2, intro AE-I2) auto
    qed
    also have ... = ennreal  $\mathcal{P}(\xi \text{ in } (\mathfrak{M} \downarrow \text{alive } x). T x \xi \geq k + 1)$ 
      by (rewrite nn-integral-cmult-indicator; simp)
    also have ... = ?RHS
      by (rewrite MM-PS.cond-prob-space-cond-prob, simp-all add: alive-T)
    finally show ?thesis .
  }
qed }
  thus ?thesis by (intro sum.cong) auto
qed
also have ... = ennreal ( $\sum k < n.$   $\mathcal{P}(\xi \text{ in } \mathfrak{M}. T x \xi \geq k + 1 \mid T x \xi > 0)$ ) by
simp
finally have  $\text{integral}^N (\mathfrak{M} \downarrow \text{alive } x) (\lambda \xi. \text{ennreal } \lfloor \min (T x \xi) n \rfloor) =$ 
  ennreal ( $\sum k < n.$   $\mathcal{P}(\xi \text{ in } \mathfrak{M}. T x \xi \geq k + 1 \mid T x \xi > 0)$ ) .
hence  $\text{integral}^L (\mathfrak{M} \downarrow \text{alive } x) (\lambda \xi. \lfloor \min (T x \xi) n \rfloor) =$ 

```

$(\sum k < n. \mathcal{P}(\xi \text{ in } \mathfrak{M}. T x \xi \geq k + 1 \mid T x \xi > 0))$
apply (*intro nn-integral-eq-integrable*[*THEN iffD1, THEN conjunct2*]; *simp*)
using *MM-PS.cond-prob-nonneg* **by** (*meson sum-nonneg*)
thus *?thesis unfolding temp-curt-life-expect-def* **by** *simp*
qed

corollary *curt-e-eq-temp*: $\$e\text{-}x = \$e\text{-}\{x:n\}$ **if** $x+n+1 > \$\psi$ **for** $n::\text{nat}$
using *curt-e-sum-P-finite temp-curt-e-sum-P* **that** **by** *simp*

lemma *temp-curt-e-sum-p*: $\$e\text{-}\{x:n\} = (\sum k < n. \$p\text{-}\{k+1\&x\})$
if $\bigwedge k::\text{nat}. k < n \implies \text{isCont } (\lambda t. \$p\text{-}\{t\&x\}) (\text{real } k + 1)$ **for** $n::\text{nat}$
proof –
have $\bigwedge k::\text{nat}. k < n \implies \$p\text{-}\{k+1\&x\} = \mathcal{P}(\xi \text{ in } \mathfrak{M}. T x \xi \geq k + 1 \mid T x \xi > 0)$
apply (*rewrite p-PTx-ge-ccdf-isCont, simp-all*)
using *that isCont-ccdfX-ccdfTx unfolding survive-def* **by** (*smt (verit) of-nat-0-le-iff x-lt-psi*)
thus *?thesis*
apply (*rewrite temp-curt-e-sum-P*)
by (*rule sum.cong*) *simp-all*
qed

lemma *temp-curt-e-rec*: $\$e\text{-}\{x:n\} = \$p\text{-}x * (1 + \$e\text{-}\{x+1:n-1\})$
if $\bigwedge k::\text{nat}. k < n \implies \text{isCont } (\lambda t. \$p\text{-}\{t\&x\}) (\text{real } k + 1)$ $x+1 < \$\psi$ $n \neq 0$ **for** $n::\text{nat}$
proof –
have [*simp*]: $\$p\text{-}x \neq 0$ **using** *p-0-equiv* **that** **by** *auto*
have [*simp*]: $\bigwedge k. k < n-1 \implies \text{isCont } (\lambda t. \$p\text{-}\{t\&x+1\}) (\text{real } k + 1)$
proof –
fix $k::\text{nat}$ **assume** $k-n: k < n-1$
have $\text{isCont } (\lambda t. \$p\text{-}x * \$p\text{-}\{t-1\&x+1\}) (\text{real } k + 2)$
proof –
let $?S = \{\text{real } k + 1 <..< \text{real } k + 3\}$
have *open* $?S$ **by** *simp*
moreover **have** $\text{real } k + 2 \in ?S$ **by** *simp*
moreover **have** $\bigwedge t. t \in ?S \implies \$p\text{-}x * \$p\text{-}\{t-1\&x+1\} = \$p\text{-}\{t\&x\}$
using *p-mult*
by (*smt (verit, del-insts) greaterThanLessThan-iff of-nat-0-le-iff that(2) x-lt-psi*)
ultimately show *?thesis*
apply (*rewrite isCont-cong*[**where** $g = \lambda t. \$p\text{-}\{t\&x\}$])
apply (*rewrite eventually-nhds, blast*)
using *that k-n* **by** (*smt (verit) less-diff-conv of-nat-1 of-nat-add*)
qed
hence $\text{isCont } (\lambda t. \$p\text{-}x * \$p\text{-}\{t-1\&x+1\} / \$p\text{-}x)$ ($\text{real } k + 2$)
by (*intro isCont-divide*[**where** $g = \lambda t. \$p\text{-}x$], *auto*)
hence $\text{isCont } ((\lambda t. \$p\text{-}\{t-1\&x+1\}) \circ (\lambda t. t+1))$ ($\text{real } k + 1$)
by *simp* (*rule continuous-at-compose, simp-all add: add commute*)
thus $\text{isCont } (\lambda t. \$p\text{-}\{t\&x+1\})$ ($\text{real } k + 1$) **unfolding** *comp-def* **by** *simp*

qed
have $\$p\text{-}x * (1 + \$e\text{-}\{x+1:n-1\}) = \$p\text{-}x * (1 + (\sum k < (n-1). \$p\text{-}\{k+1\&x+1\}))$
by (*rewrite survival-model.temp-curt-e-sum-p; simp add: survival-model-axioms that*)
also have $\dots = \$p\text{-}x + (\sum k < (n-1). \$p\text{-}x * \$p\text{-}\{k+1\&x+1\})$
apply (*rewrite distrib-left, simp*)
by (*rewrite vector-space-over-itself.scale-sum-right, simp*)
also have $\dots = \$p\text{-}x + (\sum k < (n-1). \$p\text{-}\{k+2\&x\})$
apply (*rewrite sum.cong, simp-all add: add commute*)
using *p-mult* **by** (*smt (verit) of-nat-0-le-iff that(2) x-lt-psi*)
also have $\dots = (\sum k < \text{Suc}(n-1). \$p\text{-}\{k+1\&x\})$
apply (*rewrite lessThan-atLeast0+*)
by (*rewrite sum.atLeast0-lessThan-Suc-shift*) *simp*
also have $\dots = \$e\text{-}\{x:n\}$ **using** *that* **by** (*rewrite temp-curt-e-sum-p; simp*)
finally show *?thesis* **by** *simp*
qed

end

lemma *p-set-integrable-shift*:

set-integrable lborel $\{0..\}$ $(\lambda t. \$p\text{-}\{t\&0\}) \longleftrightarrow \text{set-integrable lborel } \{0..\} (\lambda t. \$p\text{-}\{t\&x\})$

if $x < \$\psi$ **for** $x :: \text{real}$

proof –

have *set-integrable lborel* $\{0..\}$ $(\lambda t. \$p\text{-}\{t\&0\}) \longleftrightarrow \text{set-integrable lborel } \{x..\} (\lambda t. \$p\text{-}\{t\&0\})$

by (*rule set-integrable-Ici-equiv*)

(*metis (no-types, lifting) cdfX-integrable-Icc cdfX-p set-integrable-cong*)

also have $\dots \longleftrightarrow \text{set-integrable lborel } \{0..\} (\lambda t. \$p\text{-}\{x+t\&0\})$

using *set-integrable-Ici-shift*[of x] **by** *force*

also have $\dots \longleftrightarrow \text{set-integrable lborel } \{0..\} (\lambda t. \$p\text{-}\{x+t\&0\} / \$p\text{-}\{x\&0\})$

using *that p-0-equiv* **by** (*rewrite set-integrable-mult-divide-iff; simp*)

also have $\dots \longleftrightarrow \text{set-integrable lborel } \{0..\} (\lambda t. \$p\text{-}\{t\&x\})$

by (*rule set-integrable-cong; simp*) (*simp add: cdfTx-cdfX cdfX-p survive-def that*)

finally show *?thesis* .

qed

lemma *e-p-e*: $\$e\text{'o-}x = \$e\text{'o-}\{x:n\} + \$p\text{-}\{n\&x\} * \$e\text{'o-}(x+n)$

if *set-integrable lborel* $\{0..\}$ $(\lambda t. \$p\text{-}\{t\&x\})$ $n \geq 0$ $x+n < \$\psi$ **for** $x n :: \text{real}$

proof –

have [*simp*]: *ereal* $x < \$\psi$ **using** *that* **by** (*simp add: ereal-less-le*)

hence $\$e\text{'o-}x = (\text{LBINT } t:\{0..\}. \$p\text{-}\{t\&x\})$ **by** (*simp add: e-LBINT-p*)

also have $\dots = (\text{LBINT } t:\{0..n\}. \$p\text{-}\{t\&x\}) + (\text{LBINT } t:\{n..\}. \$p\text{-}\{t\&x\})$

proof –

have *AE* t *in lborel*. $\neg (t \in \{0..n\} \wedge t \in \{n..\})$ **using** *AE-lborel-singleton* **by** *force*

moreover have $\{0..\} = \{0..n\} \cup \{n..\}$ **using** *that* **by** *auto*

moreover have *set-integrable lborel* $\{0..n\}$ $(\lambda t. \$p\text{-}\{t\&x\})$

using *that*

by (metis Icc-subset-Ici-iff atLeastAtMost-borel order.refl set-integrable-subset sets-lborel)

moreover have set-integrable lborel {n..} ($\lambda t. p\{t\&x\}$)

using that by (metis atLeast-borel atLeast-subset-iff set-integrable-subset sets-lborel)

ultimately show ?thesis

using set-integral-Un-AE

by (metis (no-types, lifting) AE-cong atLeastAtMost-borel atLeast-borel sets-lborel)

qed

also have ... = (LBINT t:{0..n}. $p\{t\&x\}$) + $p\{n\&x\} * (LBINT t:\{0..\}. p\{t\&x+n\})$

proof -

have (LBINT t:{n..}. $p\{t\&x\}$) = (LBINT t:\{0..\}. $p\{n+t\&x\}$)

using lborel-set-integral-Ici-shift[of n - n, simplified] by force

also have ... = (LBINT t:\{0..\}. $p\{n\&x\} * p\{t\&x+n\}$)

apply (rule set-lebesgue-integral-cong; simp)

using that p-mult by force

finally show ?thesis by simp

qed

also have ... = $e'_{\circ}\{x:n\} + p\{n\&x\} * e'_{\circ}(x+n)$

apply (rewrite temp-e-LBINT-p, (simp-all add: that)[2])

by (rewrite e-LBINT-p; simp add: that)

finally show ?thesis .

qed

proposition x-ex-mono: $x + e'_{\circ}x \leq y + e'_{\circ}y$ if $x \leq y$ $y < \psi$ for $x y :: \text{real}$

proof -

have x-lt-psi[simp]: $e_{\text{real}} x < \psi$ using that $e_{\text{real}}\text{-less-le}$ by simp

show ?thesis

proof (cases $\langle \text{set-integrable lborel } \{0..\} (\lambda t. p\{t\&x\}) \rangle$)

case True

hence $e'_{\circ}x = e'_{\circ}\{x:y-x\} + p\{y-x\&x\} * e'_{\circ}y$ by (rewrite e-p-e[$of\ x\ y-x$]; simp add: that)

also have ... $\leq y - x + e'_{\circ}y$

proof -

have $e'_{\circ}\{x:y-x\} \leq y - x$ using that by (intro temp-e-le-n; simp)

moreover have $p\{y-x\&x\} * e'_{\circ}y \leq e'_{\circ}y$

using p-le-1 x-lt-psi that

by (smt (verit, ccfv-threshold) e-nonneg mult-less-cancel-right1)

ultimately show ?thesis by simp

qed

finally show ?thesis by simp

next

case False

hence $e'_{\circ}x = 0$

using e-LBINT-p not-integrable-integral-eq

unfolding set-integrable-def set-lebesgue-integral-def

by simp

moreover have $e'_{\circ}y = 0$

proof –
have \neg *set-integrable lborel* $\{0..\}$ $(\lambda t. \$p\{t\&y\})$
using *that False*
apply (*rewrite p-set-integrable-shift*[*THEN sym*], *simp*)
by (*rewrite p-set-integrable-shift*[*of x*]; *simp*)
thus *?thesis*
using *e-LBINT-p not-integrable-integral-eq that*
unfolding *set-integrable-def set-lebesgue-integral-def*
by *simp*
qed
ultimately show *?thesis using that by simp*
qed
qed

proposition *x-ex-const-equiv*: $x + \$e' \circ x = y + \$e' \circ y \iff \$q\{y-x\&x\} = 0$
if *set-integrable lborel* $\{0..\}$ $(\lambda t. \$p\{t\&0\})$ $x \leq y < \$\psi$ **for** $x y :: \text{real}$
proof –
have *ey*: *set-integrable lborel* $\{0..\}$ $(\lambda t. \$p\{t\&y\})$ **using** *that p-set-integrable-shift*
by *blast*
have *x-lt-psi*[*simp*]: *ereal* $x < \$\psi$ **using** *that ereal-less-le* **by** *simp*
hence *ex*: *set-integrable lborel* $\{0..\}$ $(\lambda t. \$p\{t\&x\})$ **using** *that p-set-integrable-shift*
by *blast*
show *?thesis*
proof
assume *const*: $x + \$e' \circ x = y + \$e' \circ y$
hence $0 = y - x - \$e' \circ x + \$e' \circ y$ **by** *simp*
also have $\dots = y - x - \$e' \circ \{x:y-x\} - \$p\{y-x\&x\} * \$e' \circ y + \$e' \circ y$
using *e-p-e*[*of x y-x*] *ex that by simp*
also have $\dots = (y - x - \$e' \circ \{x:y-x\}) + (1 - \$p\{y-x\&x\}) * \$e' \circ y$
by (*simp add: left-diff-distrib*)
finally have $0 = (y - x - \$e' \circ \{x:y-x\}) + (1 - \$p\{y-x\&x\}) * \$e' \circ y$.
moreover have $y - x - \$e' \circ \{x:y-x\} \geq 0$ **using** *temp-e-le-n* **that by simp**
ultimately have $(1 - \$p\{y-x\&x\}) * \$e' \circ y = 0$
by (*smt (verit, ccfv-threshold) e-nonneg mult-nonneg-nonneg p-le-1 that*
x-lt-psi)
moreover have $\$e' \circ y > 0$ **using** *that e-pos ey* **by** *simp*
ultimately have $1 - \$p\{y-x\&x\} = 0$ **by** *simp*
thus $\$q\{y-x\&x\} = 0$ **by** (*smt (verit) p-q-1 x-lt-psi*)
next
interpret *alive-x-PS*: *prob-space* $\mathfrak{M} \mid \text{alive } x$
by (*rule MM-PS.cond-prob-space-correct, simp-all add: alive-def*)
interpret *distrTx-RD*: *real-distribution* *distr* $(\mathfrak{M} \mid \text{alive } x)$ *borel* $(T x)$ **by** *simp*
assume $\$q\{y-x\&x\} = 0$
hence *p1*: $\$p\{y-x\&x\} = 1$ **using** *p-q-1* **by** (*metis add.right-neutral x-lt-psi*)
hence $\bigwedge t. t \in \{0..y-x\} \implies \$p\{t\&x\} = 1$
unfolding *survive-def* **using** *distrTx-RD.ccdf-nonincreasing*
by *simp (smt (verit) distrTx-RD.ccdf-bounded-prob)*
hence $\$e' \circ \{x:y-x\} = y - x$
using *that apply (rewrite temp-e-LBINT-p; simp)*

by (rewrite set-lebesgue-integral-cong[where $g=\lambda-. 1$]; simp)
 moreover have $e'_{\circ-x} = e'_{\circ-\{x:y-x\}} + p-\{y-x&x\} * e'_{\circ-y}$
 by (rewrite e-p-e[of $x y-x$]; simp add: that ex)
 ultimately show $x + e'_{\circ-x} = y + e'_{\circ-y}$ using p1 by simp
 qed
 qed
 end

5.2 Piecewise Differentiable Survival Function

locale smooth-survival-function = survival-model +
 assumes cdfX-piecewise-differentiable[simp]:
 (cdf (distr \mathfrak{M} borel X)) piecewise-differentiable-on UNIV
 begin

interpretation distrX-RD: real-distribution distr \mathfrak{M} borel X
 using MM-PS.real-distribution-distr by simp

5.2.1 Properties of Survival Function for X

lemma cdfX-continuous[simp]: continuous-on UNIV (cdf (distr \mathfrak{M} borel X))
 using cdfX-piecewise-differentiable piecewise-differentiable-on-imp-continuous-on
 by fastforce

corollary cdfX-borel-measurable[measurable]: cdf (distr \mathfrak{M} borel X) \in borel-measurable borel
 by (rule borel-measurable-continuous-onI) simp

lemma cdfX-nondifferentiable-finite-set[simp]:
 finite $\{x. \neg \text{cdf (distr } \mathfrak{M} \text{ borel X) differentiable at } x\}$

proof –

obtain S where

fin: finite S and $\bigwedge x. x \notin S \implies (\text{cdf (distr } \mathfrak{M} \text{ borel X)})$ differentiable at x

using cdfX-piecewise-differentiable unfolding piecewise-differentiable-on-def

by blast

hence $\{x. \neg \text{cdf (distr } \mathfrak{M} \text{ borel X) differentiable at } x\} \subseteq S$ by blast

thus ?thesis using finite-subset fin by blast

qed

lemma cdfX-differentiable-open-set: open $\{x. \text{cdf (distr } \mathfrak{M} \text{ borel X) differentiable at } x\}$

using cdfX-nondifferentiable-finite-set finite-imp-closed

by (metis (mono-tags, lifting) Collect-cong open-Collect-neg)

lemma cdfX-differentiable-borel-set[measurable, simp]:

$\{x. \text{cdf (distr } \mathfrak{M} \text{ borel X) differentiable at } x\} \in \text{sets borel}$

using cdfX-differentiable-open-set by simp

lemma cdfX-differentiable-AE:

AE x in lborel. (ccdf (distr \mathfrak{M} borel X)) differentiable at x
apply (rule AE-I'[of {x. \neg ccdf (distr \mathfrak{M} borel X) differentiable at x}], simp-all)
using ccdfX-nondifferentiable-finite-set **by** (simp add: finite-imp-null-set-lborel)

lemma deriv-ccdfX-measurable[measurable]: deriv (ccdf (distr \mathfrak{M} borel X)) \in borel-measurable borel

proof –

have set-borel-measurable borel UNIV (deriv (ccdf (distr \mathfrak{M} borel X)))
by (rule piecewise-differentiable-on-deriv-measurable-real; simp)
thus ?thesis **unfolding** set-borel-measurable-def **by** simp
qed

5.2.2 Properties of Cumulative Distributive Function for X

lemma cdfX-piecewise-differentiable[simp]:
(cdf (distr \mathfrak{M} borel X)) piecewise-differentiable-on UNIV
by (rewrite distrX-RD.cdf-ccdf) (rule piecewise-differentiable-diff; simp)

lemma cdfX-continuous[simp]: continuous-on UNIV (cdf (distr \mathfrak{M} borel X))
using cdfX-piecewise-differentiable piecewise-differentiable-on-imp-continuous-on
by fastforce

corollary cdfX-borel-measurable[measurable]: cdf (distr \mathfrak{M} borel X) \in borel-measurable borel
by (rule borel-measurable-continuous-onI) simp

lemma cdfX-nondifferentiable-finite-set[simp]:
finite {x. \neg cdf (distr \mathfrak{M} borel X) differentiable at x}
using distrX-RD.differentiable-cdf-ccdf cdfX-nondifferentiable-finite-set **by** simp

lemma cdfX-differentiable-open-set: open {x. cdf (distr \mathfrak{M} borel X) differentiable at x}
using distrX-RD.differentiable-cdf-ccdf cdfX-differentiable-open-set **by** simp

lemma cdfX-differentiable-borel-set[measurable, simp]:
{x. cdf (distr \mathfrak{M} borel X) differentiable at x} \in sets borel
using cdfX-differentiable-open-set **by** simp

lemma cdfX-differentiable-AE:
AE x in lborel. (cdf (distr \mathfrak{M} borel X)) differentiable at x
using cdfX-differentiable-AE distrX-RD.differentiable-cdf-ccdf AE-iffI **by** simp

lemma deriv-cdfX-measurable[measurable]: deriv (cdf (distr \mathfrak{M} borel X)) \in borel-measurable borel

proof –

have set-borel-measurable borel UNIV (deriv (cdf (distr \mathfrak{M} borel X)))
by (rule piecewise-differentiable-on-deriv-measurable-real; simp)
thus ?thesis **unfolding** set-borel-measurable-def **by** simp
qed

5.2.3 Introduction of Probability Density Functions of X and $T(x)$

definition $pdfX :: real \Rightarrow real$

where $pdfX\ x \equiv$ if $cdf\ (distr\ \mathfrak{M}\ borel\ X)$ differentiable at x
then $deriv\ (cdf\ (distr\ \mathfrak{M}\ borel\ X))\ x$ else 0

— This function is defined to be always nonnegative for future application.

definition $pdfT :: real \Rightarrow real \Rightarrow real$

where $pdfT\ x\ t \equiv$ if $cdf\ (distr\ (\mathfrak{M}\ \downarrow\ alive\ x)\ borel\ (T\ x))$ differentiable at t
then $deriv\ (cdf\ (distr\ (\mathfrak{M}\ \downarrow\ alive\ x)\ borel\ (T\ x)))\ t$ else 0

— This function is defined to be always nonnegative for future application.

lemma $pdfX$ -measurable[measurable]: $pdfX \in borel$ -measurable borel

proof –

let $?cdfX = cdf\ (distr\ \mathfrak{M}\ borel\ X)$

have countable $\{x. \neg ?cdfX\ \text{differentiable at } x\}$

using $cdfX$ -nondifferentiable-finite-set uncountable-infinite **by** force

thus $?thesis$

unfolding $pdfX$ -def

apply –

by (rule measurable-discrete-difference

[**where** $X = \{x. \neg ?cdfX\ \text{differentiable at } x\}$ **and** $f = deriv\ ?cdfX$]; simp)

qed

lemma distributed- $pdfX$: distributed \mathfrak{M} lborel X $pdfX$

proof –

let $?cdfX = cdf\ (distr\ \mathfrak{M}\ borel\ X)$

obtain S **where** fin : finite S **and** $diff$: $\bigwedge t. t \notin S \implies ?cdfX$ differentiable at t

using $cdfX$ -piecewise-differentiable **unfolding** piecewise-differentiable-on-def

by blast

{ **fix** $t::real$ **assume** t -notin: $t \notin S$

have $?cdfX$ differentiable at t **using** $diff$ t -notin **by** simp

hence ($?cdfX$ has-real-derivative $pdfX\ t$) (at t)

unfolding $pdfX$ -def **using** $DERIV$ -deriv-iff-real-differentiable **by** auto }

thus $?thesis$

by (intro MM - PS .distributed-deriv-cdf[**where** $S=S$]; simp add: fin)

qed

lemma $pdfT0$ - X : $pdfT\ 0 = pdfX$

unfolding $pdfT$ -def $pdfX$ -def **using** $cdfT0$ -eq- $cdfX$ psi -pos' **by** fastforce

5.2.4 Properties of Survival Function for $T(x)$

context

fixes $x::real$

assumes x -lt- psi [simp]: $x < \psi$

begin

interpretation $alive$ - PS : prob-space $\mathfrak{M} \downarrow alive\ x$

by (rule MM-PS.cond-prob-space-correct; simp add: alive-def)

interpretation *distrTx-RD*: real-distribution $\text{distr } (\mathfrak{M} \downarrow \text{alive } x) \text{ borel } (T \ x)$ by *simp*

lemma *ccdfTx-continuous-on-nonneg*[*simp*]:
continuous-on $\{0..\}$ (ccdf (distr $(\mathfrak{M} \downarrow \text{alive } x) \text{ borel } (T \ x)$))
apply (rewrite continuous-on-eq-continuous-within, auto)
apply (rewrite continuous-ccdfX-ccdfTx[THEN sym], auto)
by (metis UNIV-I ccdfX-continuous continuous-at-imp-continuous-at-within
continuous-on-eq-continuous-within)

lemma *ccdfTx-continuous*[*simp*]: continuous-on UNIV (ccdf (distr $(\mathfrak{M} \downarrow \text{alive } x) \text{ borel } (T \ x)$))

proof –

have [*simp*]: $\{..0::\text{real}\} \cup \{0..\} = \text{UNIV}$ by *auto*

have continuous-on $\{..0\}$ (ccdf (distr $(\mathfrak{M} \downarrow \text{alive } x) \text{ borel } (T \ x)$))

by (rule ccdfTx-continuous-on-nonpos) *simp*

moreover have continuous-on $\{0..\}$ (ccdf (distr $(\mathfrak{M} \downarrow \text{alive } x) \text{ borel } (T \ x)$)) by *simp*

ultimately have continuous-on $(\{..0\} \cup \{0..\})$ (ccdf (distr $(\mathfrak{M} \downarrow \text{alive } x) \text{ borel } (T \ x)$))

by (intro continuous-on-closed-Un) *simp-all*

thus ?thesis by *simp*

qed

corollary *ccdfTx-borel-measurable*[*measurable*]:

ccdf (distr $(\mathfrak{M} \downarrow \text{alive } x) \text{ borel } (T \ x)$) \in borel-measurable borel

by (rule borel-measurable-continuous-onI) *simp*

lemma *ccdfTx-nondifferentiable-finite-set*[*simp*]:

finite $\{t. \neg \text{ccdf } (\text{distr } (\mathfrak{M} \downarrow \text{alive } x) \text{ borel } (T \ x)) \text{ differentiable at } t\}$

proof –

let $?P = \lambda t. \text{ccdf } (\text{distr } (\mathfrak{M} \downarrow \text{alive } x) \text{ borel } (T \ x)) \text{ differentiable at } t$

have $\{t. t < 0 \wedge \neg ?P \ t\} = \{\}$

proof (rule equals0I)

fix t assume *asm*: $t \in \{t. t < 0 \wedge \neg ?P \ t\}$

hence $?P \ t$ using *ccdfTx-has-real-derivative-0-at-neg real-differentiable-def x-lt-psi* by *blast*

with *asm* show *False* by *simp*

qed

hence $\{t. \neg ?P \ t\} \subseteq \text{insert } 0 \ \{t. t > 0 \wedge \neg ?P \ t\}$ by *force*

moreover have finite $\{t. t > 0 \wedge \neg ?P \ t\}$

proof –

have $\{t. \neg \text{ccdf } (\text{distr } \mathfrak{M} \text{ borel } X) \text{ differentiable at } (x+t)\} =$

plus $(-x)$ ‘ $\{s. \neg \text{ccdf } (\text{distr } \mathfrak{M} \text{ borel } X) \text{ differentiable at } s\}$

by *force*

hence finite $\{t. \neg \text{ccdf } (\text{distr } \mathfrak{M} \text{ borel } X) \text{ differentiable at } (x+t)\}$

using *ccdfX-nondifferentiable-finite-set* by *simp*

thus *?thesis*
using *finite-subset differentiable-ccdfX-ccdfTx*
by (*metis (no-types, lifting) mem-Collect-eq subsetI x-lt-psi*)
qed
ultimately show *?thesis using finite-subset by auto*
qed

lemma *ccdfTx-differentiable-open-set*:
open {t. ccdf (distr (M | alive x) borel (T x)) differentiable at t}
using *ccdfTx-nondifferentiable-finite-set finite-imp-closed*
by (*metis (mono-tags, lifting) Collect-cong open-Collect-neg*)

lemma *ccdfTx-differentiable-borel-set[measurable, simp]*:
{t. ccdf (distr (M | alive x) borel (T x)) differentiable at t} ∈ sets borel
using *ccdfTx-differentiable-open-set by simp*

lemma *ccdfTx-differentiable-AE*:
AE t in lborel. (ccdf (distr (M | alive x) borel (T x))) differentiable at t
apply (*rule AE-I'*[of *{t. ¬ ccdf (distr (M | alive x) borel (T x)) differentiable at t}*]; *simp?*)
using *ccdfTx-nondifferentiable-finite-set by (simp add: finite-imp-null-set-lborel)*

lemma *ccdfTx-piecewise-differentiable[simp]*:
(ccdf (distr (M | alive x) borel (T x))) piecewise-differentiable-on UNIV
proof –
have $\forall t \in UNIV - \{t. \neg \text{ccdf (distr (M | alive x) borel (T x)) differentiable at t}\}.$
ccdf (distr (M | alive x) borel (T x)) differentiable at t
by *auto*
thus *?thesis*
unfolding *piecewise-differentiable-on-def*
using *ccdfTx-continuous ccdfTx-nondifferentiable-finite-set by blast*
qed

lemma *deriv-ccdfTx-measurable[measurable]*:
deriv (ccdf (distr (M | alive x) borel (T x))) ∈ borel-measurable borel
proof –
have *set-borel-measurable borel UNIV (deriv (ccdf (distr (M | alive x) borel (T x))))*
by (*rule piecewise-differentiable-on-deriv-measurable-real; simp*)
thus *?thesis unfolding set-borel-measurable-def by simp*
qed

5.2.5 Properties of Cumulative Distributive Function for $T(x)$

lemma *cdfTx-continuous[simp]*:
continuous-on UNIV (cdf (distr (M | alive x) borel (T x)))
using *distrTx-RD.cdf-ccdf ccdfTx-continuous by (simp add: continuous-on-eq-continuous-within)*

corollary *cdfTx-borel-measurable[measurable]*:

cdf (distr (M | alive x) borel (T x)) ∈ borel-measurable borel
by (rule borel-measurable-continuous-onI) simp

lemma *cdfTx-nondifferentiable-finite-set[simp]:*
finite {t. ¬ cdf (distr (M | alive x) borel (T x)) differentiable at t}
using *distrTx-RD.differentiable-cdf-ccdf* **by** simp

lemma *cdfTx-differentiable-open-set:*
open {t. cdf (distr (M | alive x) borel (T x)) differentiable at t}
using *distrTx-RD.differentiable-cdf-ccdf cdfTx-differentiable-open-set* **by** simp

lemma *cdfTx-differentiable-borel-set[measurable, simp]:*
{t. cdf (distr (M | alive x) borel (T x)) differentiable at t} ∈ sets borel
using *cdfTx-differentiable-open-set* **by** simp

lemma *cdfTx-differentiable-AE:*
AE t in lborel. (cdf (distr (M | alive x) borel (T x))) differentiable at t
by (rewrite *distrTx-RD.differentiable-cdf-ccdf*, simp add: *cdfTx-differentiable-AE*)

lemma *cdfTx-piecewise-differentiable[simp]:*
(cdf (distr (M | alive x) borel (T x))) piecewise-differentiable-on UNIV
using *piecewise-differentiable-diff piecewise-differentiable-const cdfTx-piecewise-differentiable*
by (rewrite *distrTx-RD.cdf-ccdf*) blast

lemma *deriv-cdfTx-measurable[measurable]:*
deriv (cdf (distr (M | alive x) borel (T x))) ∈ borel-measurable borel
proof –
have *set-borel-measurable borel UNIV (deriv (cdf (distr (M | alive x) borel (T x))))*
by (rule *piecewise-differentiable-on-deriv-measurable-real*; simp)
thus *?thesis unfolding set-borel-measurable-def* **by** simp
qed

5.2.6 Properties of Probability Density Function of $T(x)$

lemma *pdfTx-nonneg: pdfT x t ≥ 0 for t::real*
proof –
fix t
show *pdfT x t ≥ 0*
proof (cases *⟨cdf (distr (M | alive x) borel (T x)) differentiable at t⟩*)
case True
have *mono-on UNIV (cdf (distr (M | alive x) borel (T x)))*
unfolding *mono-on-def* **using** *distrTx-RD.cdf-nondecreasing* **by** blast
moreover **have** *(cdf (distr (M | alive x) borel (T x))) has-real-derivative pdfT x t (at t)*
unfolding *pdfT-def* **using** True *DERIV-deriv-iff-real-differentiable* **by** presburger
ultimately **show** *?thesis* **by** (rule *mono-on-imp-deriv-nonneg*) simp
next

```

    case False
    thus ?thesis unfolding pdfT-def by simp
  qed
qed

lemma pdfTx-neg-0: pdfT x t = 0 if t < 0 for t::real
proof -
  let ?e = -t/2
  have (cdf (distr (M | alive x) borel (T x)) has-real-derivative 0) (at t)
    apply (rewrite DERIV-cong-ev[of t t - λ-. 0 0 0], simp-all add: that)
    apply (rewrite eventually-nhds)
    apply (rule exI[of - ball t ?e], auto simp add: that)
    by (rule cdfTx-nonpos-0; simp add: dist-real-def)
  thus ?thesis unfolding pdfT-def by (meson DERIV-imp-deriv)
qed

lemma pdfTx-0-0: pdfT x 0 = 0
proof (cases ⟨cdf (distr (M | alive x) borel (T x)) differentiable at 0⟩)
  case True
  let ?cdfx = cdf (distr (M | alive x) borel (T x))
  have (?cdfx has-real-derivative 0) (at 0)
  proof -
    from True obtain c where cdfx-deriv: (?cdfx has-real-derivative c) (at 0)
    unfolding differentiable-def using has-real-derivative by blast
    hence (?cdfx has-real-derivative c) (at 0 within {..0})
      by (rule has-field-derivative-at-within)
    moreover have (?cdfx has-real-derivative 0) (at 0 within {..0})
  proof -
    have ∀F t in at 0 within {..0}. ?cdfx t = 0
  proof -
    { fix t assume t ∈ {..0::real} t ≠ 0 dist t 0 < 1
      hence ?cdfx t = 0 using cdfTx-nonpos-0 x-lt-psi by blast }
    hence ∃ d > 0::real. ∀ t ∈ {..0}. t ≠ 0 ∧ dist t 0 < d ⟶ ?cdfx t = 0 by (smt
(verit))
    thus ?thesis by (rewrite eventually-at) simp
  qed
  moreover have ?cdfx 0 = 0
  proof -
    have continuous (at 0 within {..0}) ?cdfx
      using True differentiable-imp-continuous-within differentiable-subset by
blast
    thus ?thesis by (simp add: cdfTx-nonpos-0)
  qed
  ultimately show ?thesis
    by (rewrite has-field-derivative-cong-eventually[of - λ-. 0::real 0 {..0} 0];
simp)
  qed
  ultimately have c = 0
    using has-real-derivative-iff-has-vector-derivative

```



```

apply (intro vector-derivative-unique-within[of 0 {..0} ?cdfx c 0]; blast?)
by (rewrite at-within-eq-bot-iff)
      (metis closure-lessThan islimpt-in-closure limpt-of-closure
        trivial-limit-at-left-real trivial-limit-within)
thus (?cdfx has-real-derivative 0) (at 0) using cdfx-deriv by simp
qed
thus ?thesis unfolding pdfT-def by (meson DERIV-imp-deriv)
next
  case False
  thus ?thesis unfolding pdfT-def by simp
qed

lemma pdfTx-nonpos-0: pdfT x t = 0 if t ≤ 0 for t::real
  using pdfTx-neg-0 pdfTx-0-0 that by force

lemma pdfTx-beyond-0: pdfT x t = 0 if x+t ≥ $ψ for t::real
proof (cases ‹t ≤ 0›)
  case True
  thus ?thesis using pdfTx-nonpos-0 by simp
next
  let ?cdfTx = cdf (distr (ℳ | alive x) borel (T x))
  case False
  hence t-pos: t > 0 by simp
  thus ?thesis
  proof –
    have (?cdfTx has-real-derivative 0) (at-right t)
    apply (rewrite has-field-derivative-cong-eventually[where g=λ-. 1], simp-all)
    apply (rewrite eventually-at-right-field)
    using that cdfTx-1-equiv
    by (intro exI[of -t+1], auto simp add: le-ereal-less less-eq-ereal-def)
    thus ?thesis unfolding pdfT-def
    by (meson has-real-derivative-iff-has-vector-derivative has-vector-derivative-at-within
      DERIV-deriv-iff-real-differentiable trivial-limit-at-right-real
      vector-derivative-unique-within)
  qed
qed

lemma pdfTx-pdfX: pdfT x t = pdfX (x+t) / P(ξ in ℳ. X ξ > x) if t > 0 for
t::real
proof (cases ‹cdf (distr ℳ borel X) differentiable at (x+t)›)
  case True
  let ?cdfX = cdf (distr ℳ borel X)
  let ?cdfTx = cdf (distr (ℳ | alive x) borel (T x))
  have [simp]: ?cdfTx differentiable at t using differentiable-cdfX-cdfTx that True
by simp
  have pdfT x t = deriv ?cdfTx t unfolding pdfT-def using that differentiable-cdfX-cdfTx
by simp
  hence (?cdfTx has-field-derivative (pdfT x t)) (at t)
    using True DERIV-deriv-iff-real-differentiable by simp

```

moreover have $\bigwedge u. \text{dist } u \ t < t \implies$
 $?cdfX (x+u) / \mathcal{P}(\xi \text{ in } \mathfrak{M}. X \ \xi > x) - (1 / \mathcal{P}(\xi \text{ in } \mathfrak{M}. X \ \xi > x) - 1) = ?cdfTx$
 u
proof –
fix $u::\text{real}$
assume $\text{dist } u \ t < t$
hence $[simp]: u > 0$ **using** dist-real-def **by** fastforce
have $?cdfX (x+u) / \mathcal{P}(\xi \text{ in } \mathfrak{M}. X \ \xi > x) - (1 / \mathcal{P}(\xi \text{ in } \mathfrak{M}. X \ \xi > x) - 1) =$
 $(1 - \mathcal{P}(\xi \text{ in } \mathfrak{M}. X \ \xi > x+u)) / \mathcal{P}(\xi \text{ in } \mathfrak{M}. X \ \xi > x) - (1 / \mathcal{P}(\xi \text{ in } \mathfrak{M}. X \ \xi$
 $> x) - 1)$
using $MM-PS.ccdf-distr-P \ X-RV \ distrX-RD.cdf-ccdf \ distrX-RD.prob-space$ **by**
 presburger
also have $\dots = 1 - \mathcal{P}(\xi \text{ in } \mathfrak{M}. X \ \xi > x+u) / \mathcal{P}(\xi \text{ in } \mathfrak{M}. X \ \xi > x)$
by $(\text{simp add: diff-divide-distrib})$
also have $\dots = ?cdfTx \ u$
apply $(\text{rewrite ccdfTx-PX}[THEN \ \text{sym}], \ \text{simp-all add: less-eq-real-def})$
using $\text{distrTx-RD.cdf-ccdf} \ \text{distrTx-RD.prob-space}$ **by** presburger
finally show $?cdfX (x+u) / \mathcal{P}(\xi \text{ in } \mathfrak{M}. X \ \xi > x) - (1 / \mathcal{P}(\xi \text{ in } \mathfrak{M}. X \ \xi > x)$
 $- 1) = ?cdfTx \ u .$
qed
ultimately have $((\lambda u. ?cdfX (x+u) / \mathcal{P}(\xi \text{ in } \mathfrak{M}. X \ \xi > x) - (1 / \mathcal{P}(\xi \text{ in } \mathfrak{M}. X \ \xi > x) - 1))$
 $X \ \xi > x) - 1))$
 $\text{has-field-derivative} (\text{pdfT } x \ t) (at \ t)$
apply –
by $(\text{rule has-field-derivative-transform-within}[\mathbf{where} \ d=t]; \ \text{simp add: that})$
hence $((\lambda u. ?cdfX (x+u) / \mathcal{P}(\xi \text{ in } \mathfrak{M}. X \ \xi > x)) \text{has-field-derivative} (\text{pdfT } x \ t))$
 $(at \ t)$
unfolding $\text{has-field-derivative-def}$
using $\text{has-derivative-add-const}[\mathbf{where} \ c=1 / \mathcal{P}(\xi \text{ in } \mathfrak{M}. X \ \xi > x) - 1]$ **by**
 force
hence $((\lambda u. ?cdfX (x+u) / \mathcal{P}(\xi \text{ in } \mathfrak{M}. X \ \xi > x) * \mathcal{P}(\xi \text{ in } \mathfrak{M}. X \ \xi > x))$
 $\text{has-field-derivative}$
 $\text{pdfT } x \ t * \mathcal{P}(\xi \text{ in } \mathfrak{M}. X \ \xi > x)) (at \ t)$
using $\text{DERIV-cmult-right}[\mathbf{where} \ c=\mathcal{P}(\xi \text{ in } \mathfrak{M}. X \ \xi > x)]$ **by** force
hence $((\lambda u. ?cdfX (x+u)) \text{has-field-derivative} \text{pdfT } x \ t * \mathcal{P}(\xi \text{ in } \mathfrak{M}. X \ \xi > x))$
 $(at \ t)$
unfolding $\text{has-field-derivative-def}$ **using** $\text{has-derivative-transform} \ \text{PXX-pos } x\text{-lt-psi}$
by $(\text{smt} (\text{verit}) \ \text{Collect-cong} \ \text{UNIV-I} \ \text{nonzero-mult-div-cancel-right} \ \text{times-divide-eq-left})$
hence $(?cdfX \ \text{has-field-derivative} \ \text{pdfT } x \ t * \mathcal{P}(\xi \text{ in } \mathfrak{M}. X \ \xi > x)) (at \ (x+t))$
using $\text{DERIV-at-within-shift}$ **by** force
moreover have $(?cdfX \ \text{has-field-derivative} \ \text{deriv} \ ?cdfX \ (x+t)) (at \ (x+t))$
using $\text{True} \ \text{DERIV-deriv-iff-real-differentiable}$ **by** blast
ultimately have $\text{pdfT } x \ t * \mathcal{P}(\xi \text{ in } \mathfrak{M}. X \ \xi > x) = \text{deriv} \ ?cdfX \ (x+t)$
by $(\text{simp add: DERIV-imp-deriv})$
thus $\text{pdfT } x \ t = \text{pdfX} \ (x+t) / \mathcal{P}(\xi \text{ in } \mathfrak{M}. X \ \xi > x)$
unfolding pdfX-def **using** True
by $\text{simp} (\text{metis} \ \text{PXX-pos} \ \text{nonzero-mult-div-cancel-right} \ \text{x-lt-psi} \ \text{zero-less-measure-iff})$
next
case False

hence [simp]: \neg cdf (distr (\mathfrak{M} | alive x) borel ($T x$)) differentiable at t
 using differentiable-cdfX-cdfTx that by simp
 thus pdfT $x t = \text{pdfX } (x+t) / \mathcal{P}(\xi \text{ in } \mathfrak{M}. X \xi > x)$ unfolding pdfT-def pdfX-def
 using False by simp
 qed

lemma pdfTx-measurable[measurable]: pdfT $x \in$ borel-measurable borel
proof –
 let ?cdfTx = cdf (distr (\mathfrak{M} | alive x) borel ($T x$))
 have countable $\{x. \neg ?cdfTx \text{ differentiable at } x\}$
 using cdfX-nondifferentiable-finite-set uncountable-infinite by force
 thus ?thesis
 unfolding pdfT-def
 apply –
 by (rule measurable-discrete-difference
 [where $X = \{x. \neg ?cdfTx \text{ differentiable at } x\}$ and $f = \text{deriv } ?cdfTx$]; simp)
 qed

lemma distributed-pdfTx: distributed (\mathfrak{M} | alive x) lborel ($T x$) (pdfT x)
proof –
 let ?cdfTx = cdf (distr (\mathfrak{M} | alive x) borel ($T x$))
 obtain S where fin: finite S and diff: $\bigwedge t. t \notin S \implies ?cdfTx \text{ differentiable at } t$
 using cdfTx-piecewise-differentiable unfolding piecewise-differentiable-on-def
 by blast
 { fix $t::\text{real}$ assume t-notin: $t \notin S$
 have ?cdfTx differentiable at t using diff t-notin by simp
 hence (?cdfTx has-real-derivative pdfT $x t$) (at t)
 unfolding pdfT-def using DERIV-deriv-iff-real-differentiable by force }
 thus ?thesis
 by (intro aliveX-PS.distributed-deriv-cdf[where $S=S$]; simp add: fin)
 qed

lemma nn-integral-pdfTx-1: $(\int^+ s. \text{pdfT } x s \partial \text{lborel}) = 1$
proof –
 have $(\int^+ s. \text{pdfT } x s \partial \text{lborel}) = \text{emeasure } (\text{density lborel } (\text{pdfT } x)) \text{ UNIV}$
 by (rewrite emeasure-density) simp-all
 also have $\dots = \text{emeasure } (\text{distr } (\mathfrak{M} | \text{alive } x) \text{ lborel } (T x)) \text{ UNIV}$
 using distributed-pdfTx unfolding distributed-def by simp
 also have $\dots = 1$
 by (metis distrTx-RD.emeasure-space-1 distrTx-RD.space-eq-univ distr-cong sets-lborel)
 finally show ?thesis .
 qed

corollary has-bochner-integral-pdfTx-1: has-bochner-integral lborel (pdfT x) 1
 by (rule has-bochner-integral-nn-integral; simp add: pdfTx-nonneg nn-integral-pdfTx-1)

corollary LBINT-pdfTx-1: $(\text{LBINT } s. \text{pdfT } x s) = 1$
 using has-bochner-integral-pdfTx-1 by (simp add: has-bochner-integral-integral-eq)

corollary *pdfTx-has-integral-1*: $(\text{pdfT } x \text{ has-integral } 1) \text{ UNIV}$
by (*rule nn-integral-has-integral*; *simp add: pdfTx-nonneg nn-integral-pdfTx-1*)

lemma *set-nn-integral-pdfTx-1*: $(\int^+ s \in \{0..\}, \text{pdfT } x \text{ s } \partial \text{lborel}) = 1$
apply (*rewrite nn-integral-pdfTx-1 [THEN sym]*)
apply (*rule nn-integral-cong*)
using *pdfTx-nonpos-0*
by (*metis atLeast-iff ennreal-0 indicator-simps(1) linorder-le-cases*
mult.commute mult-1 mult-zero-left)

corollary *has-bochner-integral-pdfTx-1-nonpos*:
has-bochner-integral lborel $(\lambda s. \text{pdfT } x \text{ s} * \text{indicator } \{0..\} \text{ s}) = 1$
apply (*rule has-bochner-integral-nn-integral, simp-all*)
using *pdfTx-nonneg* **apply** *simp*
using *set-nn-integral-pdfTx-1* **by** (*simp add: nn-integral-set-ennreal*)

corollary *set-LBINT-pdfTx-1*: $(\text{LBINT } s : \{0..\}, \text{pdfT } x \text{ s}) = 1$
unfolding *set-lebesgue-integral-def*
using *has-bochner-integral-pdfTx-1-nonpos has-bochner-integral-integral-eq*
apply (*simp, rewrite mult.commute*)
by (*smt (verit) Bochner-Integration.integral-cong has-bochner-integral-integral-eq*)

corollary *pdfTx-has-integral-1-nonpos*: $(\text{pdfT } x \text{ has-integral } 1) \{0..\}$
proof –

have *set-integrable lebesgue* $\{0..\}$ $(\text{pdfT } x)$
unfolding *set-integrable-def*
apply (*rewrite integrable-completion, simp-all*)
using *has-bochner-integral-pdfTx-1-nonpos* **by** (*rewrite mult.commute, rule integrable.intros*)
moreover **have** $(\text{LINT } s : \{0..\} | \text{lebesgue. pdfT } x \text{ s}) = 1$
using *set-LBINT-pdfTx-1* **unfolding** *set-lebesgue-integral-def*
by (*rewrite integral-completion; simp*)
ultimately show *?thesis* **using** *has-integral-set-lebesgue* **by force**
qed

lemma *set-nn-integral-pdfTx-PTx*: $(\int^+ s \in A. \text{pdfT } x \text{ s } \partial \text{lborel}) = \mathcal{P}(\xi \text{ in } \mathfrak{M}. T x \xi \in A \mid T x \xi > 0)$

if $A \in \text{sets lborel}$ **for** $A :: \text{real set}$

proof –

have [*simp*]: $A \in \text{sets borel}$ **using** *that* **by** *simp*
have $(\int^+ s \in A. \text{pdfT } x \text{ s } \partial \text{lborel}) = \text{emeasure}(\text{density lborel}(\text{pdfT } x)) A$
using *that* **by** (*rewrite emeasure-density; force*)
also **have** $\dots = \text{emeasure}(\text{distr } (\mathfrak{M} \mid \text{alive } x) \text{ lborel } (T x)) A$
using *distributed-pdfTx* **unfolding** *distributed-def* **by** *simp*
also **have** $\dots = \text{ennreal } \mathcal{P}(\xi \text{ in } (\mathfrak{M} \mid \text{alive } x). T x \xi \in A)$
apply (*rewrite emeasure-distr, simp-all*)
apply (*rewrite finite-measure.emeasure-eq-measure, simp*)
by (*smt (verit) Collect-cong vimage-eq Int-def*)

also have $\dots = \text{ennreal } \mathcal{P}(\xi \text{ in } \mathfrak{M}. T x \xi \in A \mid T x \xi > 0)$
unfolding *alive-def*
apply (*rewrite MM-PS.cond-prob-space-cond-prob[THEN sym]*, *simp-all add:*
pred-def futr-life-def)
using *borel-measurable-diff X-RV that by measurable*
finally show *?thesis* .
qed

lemma *pdfTx-set-integrable: set-integrable lborel A (pdfT x) if A ∈ sets lborel*
unfolding *set-integrable-def*
using *that pdfTx-nonneg apply (intro integrableI-nonneg, simp-all)*
apply (*rewrite mult.commute*)
using *set-nn-integral-pdfTx-PTx that*
by (*metis (no-types, lifting) ennreal-indicator ennreal-less-top ennreal-mult' nn-integral-cong*)

lemma *set-integral-pdfTx-PTx: (LBINT s:A. pdfT x s) = $\mathcal{P}(\xi \text{ in } \mathfrak{M}. T x \xi \in A \mid T x \xi > 0)$*
if *A ∈ sets lborel for A :: real set*
unfolding *set-lebesgue-integral-def*
apply (*rewrite integral-eq-nn-integral*)
using *that apply (simp-all add: pdfTx-nonneg)*
apply (*rewrite indicator-mult-ennreal[THEN sym]*, *rewrite mult.commute*)
by (*rewrite set-nn-integral-pdfTx-PTx; simp*)

lemma *pdfTx-has-integral-PTx: (pdfT x has-integral $\mathcal{P}(\xi \text{ in } \mathfrak{M}. T x \xi \in A \mid T x \xi > 0)$) A*
if *A ∈ sets lborel for A :: real set*
proof –
have ($(\lambda s. \text{pdfT } x \ s * \text{indicat-real } A \ s)$ *has-integral $\mathcal{P}(\xi \text{ in } \mathfrak{M}. T x \xi \in A \mid T x \xi > 0)$*) *UNIV*
using *that pdfTx-nonneg apply (intro nn-integral-has-integral, simp-all)*
using *set-nn-integral-pdfTx-PTx that by (simp add: nn-integral-set-ennreal)*
thus *?thesis*
by (*smt (verit) has-integral-cong has-integral-restrict-UNIV*
indicator-eq-0-iff indicator-simps(1) mult-cancel-left2 mult-eq-0-iff)
qed

corollary *pdfTx-has-integral-PTx-Icc:*
(pdfT x has-integral $\mathcal{P}(\xi \text{ in } \mathfrak{M}. a \leq T x \xi \wedge T x \xi \leq b \mid T x \xi > 0)$) {a..b} for
a b :: real
using *pdfTx-has-integral-PTx[of {a..b}] by simp*

corollary *pdfTx-integrable-on-Icc: pdfT x integrable-on {a..b} for a b :: real*
using *pdfTx-has-integral-PTx-Icc by auto*

end

5.2.7 Properties of Probability Density Function of X

lemma *pdfX-nonneg*: $\text{pdfX } x \geq 0$ for $x::\text{real}$
using *pdfTx-nonneg pdfT0-X psi-pos'* by *smt*

lemma *pdfX-nonpos-0*: $\text{pdfX } x = 0$ if $x \leq 0$ for $x::\text{real}$
using *pdfTx-nonpos-0 pdfT0-X psi-pos'* that by *smt*

lemma *pdfX-beyond-0*: $\text{pdfX } x = 0$ if $x \geq \psi$ for $x::\text{real}$
using *pdfTx-beyond-0 pdfT0-X psi-pos'* that by *smt*

lemma *nn-integral-pdfX-1*: $\text{integral}^N \text{ lborel pdfX} = 1$
using *nn-integral-pdfTx-1 pdfT0-X psi-pos'* by *metis*

corollary *has-bochner-integral-pdfX-1*: *has-bochner-integral lborel pdfX 1*
by (*rule has-bochner-integral-nn-integral; simp add: pdfX-nonneg nn-integral-pdfX-1*)

corollary *LBINT-pdfX-1*: $(\text{LBINT } s. \text{pdfX } s) = 1$
using *has-bochner-integral-pdfX-1* by (*simp add: has-bochner-integral-integral-eq*)

corollary *pdfX-has-integral-1*: $(\text{pdfX has-integral } 1)$ UNIV
by (*rule nn-integral-has-integral; simp add: pdfX-nonneg nn-integral-pdfX-1*)

lemma *set-nn-integral-pdfX-PX*: $\text{set-nn-integral lborel } A \text{ pdfX} = \mathcal{P}(\xi \text{ in } \mathfrak{M}. X \xi \in A)$
if $A \in \text{sets lborel}$ for $A :: \text{real set}$
using *PT0-eq-PX-lborel* that
by (*rewrite pdfT0-X[THEN sym], rewrite set-nn-integral-pdfTx-PTx; simp*)

lemma *pdfX-set-integrable*: $\text{set-integrable lborel } A \text{ pdfX}$ if $A \in \text{sets lborel}$ for $A :: \text{real set}$
using *pdfTx-set-integrable pdfT0-X psi-pos'* that by *smt*

lemma *set-integral-pdfX-PX*: $(\text{LBINT } s:A. \text{pdfX } s) = \mathcal{P}(\xi \text{ in } \mathfrak{M}. X \xi \in A)$
if $A \in \text{sets lborel}$ for $A :: \text{real set}$
using *PT0-eq-PX-lborel* that by (*rewrite pdfT0-X[THEN sym], rewrite set-integral-pdfTx-PTx; simp*)

lemma *pdfX-has-integral-PX*: $(\text{pdfX has-integral } \mathcal{P}(\xi \text{ in } \mathfrak{M}. X \xi \in A)) A$
if $A \in \text{sets lborel}$ for $A :: \text{real set}$
using that **apply** (*rewrite pdfT0-X[THEN sym], rewrite PT0-eq-PX-lborel[THEN sym], simp*)
by (*rule pdfTx-has-integral-PTx; simp*)

corollary *pdfX-has-integral-PX-Icc*: $(\text{pdfX has-integral } \mathcal{P}(\xi \text{ in } \mathfrak{M}. a \leq X \xi \wedge X \xi \leq b)) \{a..b\}$
for $a b :: \text{real}$
using *pdfX-has-integral-PX[of {a..b}]* by *simp*

corollary *pdfX-integrable-on-Icc*: $\text{pdfX integrable-on } \{a..b\}$ for $a b :: \text{real}$

using pdfX-has-integral-PX-Icc by auto

5.2.8 Relations between Life Expectation and Probability Density Function

context

fixes $x::real$
 assumes $x\text{-lt-psi}[simp]: x < \psi$
begin

interpretation *alivex-PS: prob-space* $\mathfrak{M} \mid \text{alive } x$

by (rule *MM-PS.cond-prob-space-correct*; simp add: *alive-def*)

interpretation *distrTx-RD: real-distribution* $\text{distr } (\mathfrak{M} \mid \text{alive } x) \text{ borel } (T \ x)$ by *simp*

proposition *nn-integral-T-pdfT:*

$(\int^{+\xi}. \text{ennreal } (g \ (T \ x \ \xi)) \ \partial(\mathfrak{M} \mid \text{alive } x)) = (\int^{+s \in \{0..\}}. \text{ennreal } (\text{pdfT } x \ s * g \ s) \ \partial \text{lborel})$

if $g \in \text{borel-measurable lborel}$ for $g :: real \Rightarrow real$

proof –

have $(\int^{+\xi}. \text{ennreal } (g \ (T \ x \ \xi)) \ \partial(\mathfrak{M} \mid \text{alive } x)) = \int^{+s}. \text{ennreal } (\text{pdfT } x \ s) * \text{ennreal } (g \ s) \ \partial \text{lborel}$

proof –

have *distributed* $(\mathfrak{M} \mid \text{alive } x) \text{ lborel } (T \ x) (\lambda s. \text{ennreal } (\text{pdfT } x \ s))$

by (*intro distributed-pdfTx*) *simp*

moreover have $(\lambda s. \text{ennreal } (g \ s)) \in \text{borel-measurable borel}$ using *that* by *measurable*

ultimately show *?thesis* by (*rewrite distributed-nn-integral*; *simp*)

qed

also have $\dots = \int^{+s}. \text{ennreal } (\text{pdfT } x \ s * g \ s) \ \partial \text{lborel}$ using *ennreal-mult'* *pdfTx-nonneg* by *force*

also have $\dots = (\int^{+s \in \{0..\}}. \text{ennreal } (\text{pdfT } x \ s * g \ s) \ \partial \text{lborel})$

apply (*rule nn-integral-cong*, *simp*)

by (*metis atLeast-iff ennreal-0 indicator-simps linorder-not-le mult-1 mult-commute-abs mult-zero-left pdfTx-neg-0 x-lt-psi*)

finally show *?thesis* .

qed

lemma *expectation-LBINT-pdfT-nonneg:*

alivex-PS.expectation $(\lambda \xi. g \ (T \ x \ \xi)) = (\text{LBINT } s:\{0..\}. \text{pdfT } x \ s * g \ s)$

if $\bigwedge s. s \geq 0 \implies g \ s \geq 0$ $g \in \text{borel-measurable lborel}$ for $g :: real \Rightarrow real$

— Note that $0 = 0$ holds when the integral diverges.

using *that* apply (*rewrite integral-eq-nn-integral*, *simp*)

apply (*rule AE-I2*, *metis alivex-Tx-pos less-imp-le*)

unfolding *set-lebesgue-integral-def* apply (*rewrite integral-eq-nn-integral*, *simp-all*)

apply (*rule AE-I2*,

metis indicator-pos-le pdfTx-nonpos-0 x-lt-psi zero-le-mult-iff zero-le-square pdfTx-nonneg)

by (rewrite nn-integral-T-pdfT) (simp-all add: indicator-mult-ennreal mult.commute)

corollary *expectation-integral-pdfT-nonneg*:

alivex-PS.expectation $(\lambda\xi. g (T x \xi)) = \text{integral } \{0..\} (\lambda s. \text{pdfT } x s * g s)$

if $\bigwedge s. s \geq 0 \implies g s \geq 0$ $g \in \text{borel-measurable lborel}$ **for** $g :: \text{real} \Rightarrow \text{real}$

— Note that $0 = 0$ holds when the integral diverges.

proof —

have *alivex-PS.expectation* $(\lambda\xi. g (T x \xi)) = (\text{LBINT } s:\{0..\}. \text{pdfT } x s * g s)$

using *expectation-LBINT-pdfT-nonneg* **that** **by** *simp*

also have $\dots = \text{integral } \{0..\} (\lambda s. \text{pdfT } x s * g s)$

using *that pdfTx-nonneg* **by** (*intro set-borel-integral-eq-integral-nonneg; simp*)

finally show *?thesis* .

qed

proposition *expectation-LBINT-pdfT*:

alivex-PS.expectation $(\lambda\xi. g (T x \xi)) = (\text{LBINT } s:\{0..\}. \text{pdfT } x s * g s)$

if *set-integrable lborel* $\{0..\} (\lambda s. \text{pdfT } x s * g s)$ $g \in \text{borel-measurable lborel}$

for $g :: \text{real} \Rightarrow \text{real}$

proof —

have *pg-fin*: $(\int^{+\xi}. \text{ennreal } (g (T x \xi)) \partial(\mathfrak{M} \downarrow \text{alive } x)) \neq \infty$

using *that apply* (*rewrite nn-integral-T-pdfT, simp*)

using *that unfolding set-integrable-def apply* (*rewrite in asm real-integrable-def, simp*)

by (*simp add: indicator-mult-ennreal mult.commute*)

moreover have *mg-fin*: $(\int^{+\xi}. \text{ennreal } (- g (T x \xi)) \partial(\mathfrak{M} \downarrow \text{alive } x)) \neq \infty$

using *that apply* (*rewrite nn-integral-T-pdfT[of $\lambda s. - g s$], simp*)

using *that unfolding set-integrable-def apply* (*rewrite in asm real-integrable-def, simp*)

by (*simp add: indicator-mult-ennreal mult.commute*)

ultimately have [*simp*]: *integrable* $(\mathfrak{M} \downarrow \text{alive } x) (\lambda\xi. g (T x \xi))$

using *that by* (*rewrite real-integrable-def; simp*)

have $(\int^{+s \in \{0..\}. \text{ennreal } (\text{pdfT } x s * \max 0 (g s)) \partial \text{lborel}} =$

$(\int^{+s \in \{0..\}. \text{ennreal } (\text{pdfT } x s * g s) \partial \text{lborel}})$

using *SPMF.ennreal-max-0 ennreal-mult' pdfTx-nonneg x-lt-psi* **by** *presburger*

also have $\dots < \infty$

using *pg-fin nn-integral-T-pdfT that top.not-eq-extremum* **by** *auto*

finally have $(\int^{+s \in \{0..\}. \text{ennreal } (\text{pdfT } x s * \max 0 (g s)) \partial \text{lborel}} < \infty$.

hence [*simp*]: *set-integrable lborel* $\{0..\} (\lambda s. \text{pdfT } x s * \max 0 (g s))$

unfolding *set-integrable-def* **using** *that apply* (*intro integrableI-nonneg, simp-all*)

using *pdfTx-nonneg apply* (*intro AE-I2, force*)

by (*metis (no-types, lifting) indicator-mult-ennreal mult.commute nn-integral-cong*)

have $(\int^{+s \in \{0..\}. \text{ennreal } (\text{pdfT } x s * \max 0 (- g s)) \partial \text{lborel}} =$

$(\int^{+s \in \{0..\}. \text{ennreal } (\text{pdfT } x s * - g s) \partial \text{lborel}})$

using *SPMF.ennreal-max-0 ennreal-mult' pdfTx-nonneg x-lt-psi* **by** *presburger*

also have $\dots < \infty$

using *mg-fin nn-integral-T-pdfT[of $\lambda s. - g s$] that top.not-eq-extremum* **by** *force*

finally have $(\int^{+s \in \{0..\}. \text{ennreal } (\text{pdfT } x s * \max 0 (- g s)) \partial \text{lborel}} < \infty$.

hence [*simp*]: *set-integrable lborel* $\{0..\} (\lambda s. \text{pdfT } x s * \max 0 (- g s))$

unfolding *set-integrable-def* **using** *that apply* (*intro integrableI-nonneg, simp-all*)
using *pdfTx-nonneg apply* (*intro AE-I2, force*)
by (*metis (no-types, lifting) indicator-mult-ennreal mult commute nn-integral-cong*)
have *alivex-PS.expectation* $(\lambda\xi. g (T x \xi)) =$
alivex-PS.expectation $(\lambda\xi. \max 0 (g (T x \xi))) - \text{alivex-PS.expectation } (\lambda\xi. \max$
 $0 (-g (T x \xi)))$
by (*rewrite Bochner-Integration.integral-cong*
 $[\text{where } N = \mathfrak{M} \downarrow \text{alivex } x \text{ and } g = \lambda\xi. \max 0 (g (T x \xi)) - \max 0 (-g (T x$
 $\xi))]; \text{simp}$)
moreover have *alivex-PS.expectation* $(\lambda\xi. \max 0 (g (T x \xi))) =$
 $(\text{LBINT } s:\{0..\}. \text{pdfT } x s * \max 0 (g s))$
using that by (*rewrite expectation-LBINT-pdfT-nonneg* $[\text{where } g = \lambda s. \max 0$
 $(g s)]; \text{simp-all}$)
moreover have *alivex-PS.expectation* $(\lambda\xi. \max 0 (-g (T x \xi))) =$
 $(\text{LBINT } s:\{0..\}. \text{pdfT } x s * \max 0 (-g s))$
using that by (*rewrite expectation-LBINT-pdfT-nonneg* $[\text{where } g = \lambda s. \max 0$
 $(-g s)]; \text{simp-all}$)
ultimately have *alivex-PS.expectation* $(\lambda\xi. g (T x \xi)) =$
 $(\text{LBINT } s:\{0..\}. \text{pdfT } x s * \max 0 (g s)) - (\text{LBINT } s:\{0..\}. \text{pdfT } x s * \max 0$
 $(-g s))$ **by** *simp*
also have $\dots = (\text{LBINT } s:\{0..\}. (\text{pdfT } x s * \max 0 (g s) - \text{pdfT } x s * \max 0$
 $(-g s)))$
by (*rewrite set-integral-diff; simp*)
also have $\dots = (\text{LBINT } s:\{0..\}. \text{pdfT } x s * (\max 0 (g s) - \max 0 (-g s)))$
by (*simp add: right-diff-distrib*)
also have $\dots = (\text{LBINT } s:\{0..\}. \text{pdfT } x s * g s)$
using *minus-max-eq-min*
by (*metis (no-types, opaque-lifting) diff-zero max-def min-def minus-diff-eq*)
finally show *?thesis* .
qed

corollary *expectation-integral-pdfT*:

alivex-PS.expectation $(\lambda\xi. g (T x \xi)) = \text{integral } \{0..\} (\lambda s. \text{pdfT } x s * g s)$
if $(\lambda s. \text{pdfT } x s * g s)$ *absolutely-integrable-on* $\{0..\}$ $g \in \text{borel-measurable lborel}$
for $g :: \text{real} \Rightarrow \text{real}$

proof –

have $[\text{simp}]: \text{set-integrable lborel } \{0..\} (\lambda s. \text{pdfT } x s * g s)$
using that by (*rewrite absolutely-integrable-on-iff-set-integrable; simp*)
show *?thesis*
apply (*rewrite set-borel-integral-eq-integral(2)[THEN sym], simp*)
using that by (*rewrite expectation-LBINT-pdfT; simp*)

qed

corollary *e-LBINT-pdfT*: $\$e' \circ -x = (\text{LBINT } s:\{0..\}. \text{pdfT } x s * s)$

– Note that $0 = 0$ holds when the life expectation diverges.

unfolding *life-expect-def* **using** *expectation-LBINT-pdfT-nonneg* **by** *force*

corollary *e-integral-pdfT*: $\$e' \circ -x = \text{integral } \{0..\} (\lambda s. \text{pdfT } x s * s)$

– Note that $0 = 0$ holds when the life expectation diverges.

unfolding *life-expect-def* **using** *expectation-integral-pdfT-nonneg* **by** *force*
end

corollary *e-LBINT-pdfX*: $\$e' \circ - 0 = (LBINT\ x:\{0..\}. pdfX\ x * x)$
— Note that $0 = 0$ holds when the life expectation diverges.
using *e-LBINT-pdfT pdfT0-X psi-pos'* **by** *presburger*

corollary *e-integral-pdfX*: $\$e' \circ - 0 = integral\ \{0..\}\ (\lambda x. pdfX\ x * x)$
— Note that $0 = 0$ holds when the life expectation diverges.
using *e-integral-pdfT pdfT0-X psi-pos'* **by** *presburger*

5.2.9 Introduction of Force of Mortality

definition *force-mortal* :: *real* \Rightarrow *real* ($\$ \mu' - [101]\ 200$)
where $\$ \mu - x \equiv MM-PS.hazard-rate\ X\ x$

lemma *mu-pdfX*: $\$ \mu - x = pdfX\ x / cdf\ (distr\ \mathfrak{M}\ borel\ X)\ x$
if $(cdf\ (distr\ \mathfrak{M}\ borel\ X))$ *differentiable at x* **for** $x::real$
unfolding *force-mortal-def pdfX-def*
by (*rewrite MM-PS.hazard-rate-deriv-cdf, simp-all add: that*)

lemma *mu-unborn-0*: $\$ \mu - x = 0$ **if** $x < 0$ **for** $x::real$
apply (*rewrite mu-pdfX*)
using *cdfX-has-real-derivative-0-unborn real-differentiable-def that* **apply** *blast*
using *pdfX-nonpos-0 that* **by** *auto*

lemma *mu-beyond-0*: $\$ \mu - x = 0$ **if** $x \geq \$ \psi$ **for** $x::real$
— Note that division by 0 is defined as 0 in Isabelle/HOL.
unfolding *force-mortal-def* **using** *MM-PS.hazard-rate-0-cdf-0 cdfX-0-equiv that*
by *simp*

lemma *mu-nonneg-differentiable*: $\$ \mu - x \geq 0$
if $(cdf\ (distr\ \mathfrak{M}\ borel\ X))$ *differentiable at x* **for** $x::real$
apply (*rewrite mu-pdfX, simp add: that*)
using *pdfX-nonneg distrX-RD.cdf-nonneg* **by** *simp*

lemma *mu-nonneg-AE*: *AE x in lborel. $\$ \mu - x \geq 0$*
using *cdfX-differentiable-AE mu-nonneg-differentiable* **by** *auto*

lemma *mu-measurable[measurable]*: $(\lambda x. \$ \mu - x) \in borel-measurable\ borel$

proof —

obtain *S* **where**

finite S **and** $\bigwedge x. x \notin S \implies (cdf\ (distr\ \mathfrak{M}\ borel\ X))$ *differentiable at x*

using *cdfX-piecewise-differentiable* **unfolding** *piecewise-differentiable-on-def*

by *blast*

thus *?thesis*

apply (*rewrite measurable-discrete-difference*

[**where** $f = \lambda x. pdfX\ x / cdf\ (distr\ \mathfrak{M}\ borel\ X)\ x$ **and** $X = S$], *simp-all*)

by (simp-all add: MM-PS.ccdf-distr-measurable borel-measurable-divide count-able-finite mu-pdfX)

qed

lemma mu-deriv-ccdf: $\$ \mu - x = - \text{deriv } (\text{ccdf } (\text{distr } \mathfrak{M} \text{ borel } X)) x / \text{ccdf } (\text{distr } \mathfrak{M} \text{ borel } X) x$

if (ccdf (distr \mathfrak{M} borel X)) differentiable at x $x < \$\psi$ for $x::\text{real}$

unfolding force-mortal-def

by (rewrite MM-PS.hazard-rate-deriv-ccdf; simp add: that)

lemma mu-deriv-ln: $\$ \mu - x = - \text{deriv } (\lambda x. \ln (\text{ccdf } (\text{distr } \mathfrak{M} \text{ borel } X) x)) x$

if (ccdf (distr \mathfrak{M} borel X)) differentiable at x $x < \$\psi$ for $x::\text{real}$

unfolding force-mortal-def

apply (rewrite MM-PS.hazard-rate-deriv-ln-ccdf, simp-all add: that)

using ccdfX-0-equiv that by force

lemma p-exp-integral-mu: $\$ p - \{t \& x\} = \exp (- \text{integral } \{x..x+t\} (\lambda y. \$ \mu - y))$

if $x \geq 0$ $t \geq 0$ $x+t < \$\psi$ for $x t :: \text{real}$

proof -

have [simp]: $x < \$\psi$ using that by (simp add: ereal-less-le)

have $\$ p - \{t \& x\} = (\text{ccdf } (\text{distr } \mathfrak{M} \text{ borel } X) (x+t)) / (\text{ccdf } (\text{distr } \mathfrak{M} \text{ borel } X) x)$

apply (rewrite p-PX, simp-all add: that)

by (rewrite MM-PS.ccdf-distr-P, simp)+ simp

also have $\dots = \exp (- \text{integral } \{x..x+t\} (\text{MM-PS.hazard-rate } X))$

apply (rule MM-PS.ccdf-exp-cumulative-hazard, simp-all add: that)

using ccdfX-piecewise-differentiable piecewise-differentiable-on-subset apply

blast

using ccdfX-continuous continuous-on-subset apply blast

using ccdfX-0-equiv that ereal-less-le linorder-not-le by force

also have $\dots = \exp (- \text{integral } \{x..x+t\} (\lambda y. \$ \mu - y))$ unfolding force-mortal-def

by simp

finally show ?thesis .

qed

corollary ccdfX-exp-integral-mu: $\text{ccdf } (\text{distr } \mathfrak{M} \text{ borel } X) x = \exp (- \text{integral } \{0..x\} (\lambda y. \$ \mu - y))$

if $0 \leq x \wedge x < \$\psi$ for $x::\text{real}$

by (rewrite p-exp-integral-mu[where $t=x$ and $x=0$, simplified, THEN sym]; simp add: that ccdfX-p)

5.2.10 Properties of Force of Mortality

context

fixes $x::\text{real}$

assumes $x - lt - psi$ [simp]: $x < \$\psi$

begin

interpretation *aliveX-PS*: prob-space $\mathfrak{M} \mid \text{alive } x$

by (rule MM-PS.cond-prob-space-correct; simp add: alive-def)

interpretation *distrTx-RD*: real-distribution $\text{distr } (\mathfrak{M} \downarrow \text{alive } x) \text{ borel } (T x)$ by *simp*

lemma *hazard-rate-Tx-mu*: $\text{alivex-PS.hazard-rate } (T x) t = \$\mu\text{-}(x+t)$

if $t \geq 0$ $x+t < \$\psi$ for $t::\text{real}$

proof –

have $[\text{simp}]$: $\mathcal{P}(\xi \text{ in } \mathfrak{M}. X \xi > x) > 0$ by *force*

moreover have $[\text{simp}]$: $\mathcal{P}(\xi \text{ in } \mathfrak{M}. X \xi > x + t) > 0$ using *that* by *force*

moreover have $[\text{simp}]$: $\{\xi \in \text{space } (\mathfrak{M} \downarrow \text{alive } x). X \xi > x + t\} = \{\xi \in \text{space } \mathfrak{M}. X \xi > x + t\}$

unfolding *alive-def* using *that* by (rewrite *MM-PS.space-cond-prob-space*, *simp-all*, *force*)

hence $[\text{simp}]$: $\{\xi \in \text{space } (\mathfrak{M} \downarrow \text{alive } x). X \xi > x + t\} \in \text{MM-PS.events}$ by *simp*

ultimately have $\star[\text{simp}]$: $\mathcal{P}(\xi \text{ in } (\mathfrak{M} \downarrow \text{alive } x). X \xi > x + t) > 0$

unfolding *alive-def*

apply (rewrite *MM-PS.cond-prob-space-cond-prob[THEN sym]*, (*simp-all add: pred-def*)[3])

unfolding *cond-prob-def* by (*smt (verit) Collect-cong divide-pos-pos*)

have $\text{alivex-PS.hazard-rate } (T x) t =$

$\text{Lim } (\text{at-right } 0) (\lambda dt. \mathcal{P}(\xi \text{ in } (\mathfrak{M} \downarrow \text{alive } x). t < T x \xi \wedge T x \xi \leq t + dt \mid T x \xi > t) / dt)$

unfolding *alivex-PS.hazard-rate-def* by *simp*

also have $\dots = \text{Lim } (\text{at-right } 0)$

$(\lambda dt. \mathcal{P}(\xi \text{ in } (\mathfrak{M} \downarrow \text{alive } x). x + t < X \xi \wedge X \xi \leq x + t + dt \mid X \xi > x + t) / dt)$

apply (*rule Lim-cong*, *rule eventually-at-rightI[of 0 1]*, *simp-all*)

unfolding *futr-life-def cond-prob-def* using *Collect-cong* by *smt*

also have $\dots =$

$\text{Lim } (\text{at-right } 0) (\lambda dt. \mathcal{P}(\xi \text{ in } \mathfrak{M}. x + t < X \xi \wedge X \xi \leq x + t + dt \mid X \xi > x + t) / dt)$

proof –

have $\bigwedge dt. \mathcal{P}(\xi \text{ in } (\mathfrak{M} \downarrow \text{alive } x). x + t < X \xi \wedge X \xi \leq x + t + dt \mid X \xi > x + t) =$

$\mathcal{P}(\xi \text{ in } \mathfrak{M}. x + t < X \xi \wedge X \xi \leq x + t + dt \mid X \xi > x + t)$

proof –

fix dt

let $?rngX = \lambda \xi. x + t < X \xi \wedge X \xi \leq x + t + dt$

have $\mathcal{P}(\xi \text{ in } (\mathfrak{M} \downarrow \text{alive } x). ?rngX \xi \mid X \xi > x + t) =$

$\mathcal{P}(\xi \text{ in } ((\mathfrak{M} \downarrow \text{alive } x) \downarrow \{\eta \in \text{space } (\mathfrak{M} \downarrow \text{alive } x). X \eta > x + t\}). ?rngX \xi)$

using \star by (rewrite *alivex-PS.cond-prob-space-cond-prob*) *simp-all*

also have $\dots = \mathcal{P}(\xi \text{ in } (\mathfrak{M} \downarrow \{\eta \in \text{space } \mathfrak{M}. X \eta > x + t\}). ?rngX \xi)$

proof –

have $(\mathfrak{M} \downarrow \text{alive } x) \downarrow \{\eta \in \text{space } (\mathfrak{M} \downarrow \text{alive } x). X \eta > x + t\} =$

$\mathfrak{M} \downarrow \{\eta \in \text{space } (\mathfrak{M} \downarrow \text{alive } x). X \eta > x + t\}$

apply (rewrite *MM-PS.cond-cond-prob-space*, *simp-all*)

unfolding *alive-def* using *that* by *force*

also have $\dots = \mathfrak{M} \downarrow \{\eta \in \text{space } \mathfrak{M}. X \eta > x + t\}$ by *simp*

finally have $(\mathfrak{M} \downarrow \text{alive } x) \downarrow \{\eta \in \text{space } (\mathfrak{M} \downarrow \text{alive } x). X \eta > x + t\} =$

$\mathfrak{M} \mid \{\eta \in \text{space } \mathfrak{M}. X \eta > x + t\} .$
thus *?thesis* **by simp**
qed
also have $\dots = \mathcal{P}(\xi \text{ in } \mathfrak{M}. x + t < X \xi \wedge X \xi \leq x + t + dt \mid X \xi > x + t)$
by (*rewrite MM-PS.cond-prob-space-cond-prob, simp-all add: pred-def sets.sets-Collect-conj*)
finally show $\mathcal{P}(\xi \text{ in } (\mathfrak{M} \mid \text{alive } x). ?rng X \xi \mid X \xi > x + t) =$
 $\mathcal{P}(\xi \text{ in } \mathfrak{M}. x + t < X \xi \wedge X \xi \leq x + t + dt \mid X \xi > x + t) .$
qed
thus *?thesis*
apply –
by (*rule Lim-cong, rule eventually-at-rightI[of 0 1]; simp*)
qed
also have $\dots = \$\mu_{-(x+t)}$ **unfolding** *force-mortal-def MM-PS.hazard-rate-def*
by simp
finally show *?thesis* .
qed

lemma *pdfTx-p-mu*: $\text{pdfT } x \ t = \$p_{-}\{t\&x\} * \$\mu_{-(x+t)}$
if (*cdf (distr ($\mathfrak{M} \mid \text{alive } x$) borel (T x)) differentiable at t t > 0 for t::real*)
proof (*cases (x+t < \$ψ)*)
case True
hence [*simp*]: $t \geq 0$ **and** (*ccdf (distr ($\mathfrak{M} \mid \text{alive } x$) borel (T x)) t ≠ 0*)
using *that p-0-equiv unfolding survive-def by auto*
have *deriv (cdf (distr ($\mathfrak{M} \mid \text{alive } x$) borel (T x)) t =*
*ccdf (distr ($\mathfrak{M} \mid \text{alive } x$) borel (T x)) t * alivex-PS.hazard-rate (T x) t*
by (*rule alivex-PS.deriv-cdf-hazard-rate; simp add: that*)
thus *?thesis unfolding survive-def pdfT-def using hazard-rate-Tx-mu True that*
by simp
next
case False
hence $x+t \geq \$\psi$ **by simp**
thus *?thesis using pdfTx-beyond-0 mu-beyond-0 by simp*
qed

lemma *deriv-t-p-mu*: $\text{deriv } (\lambda s. \$p_{-}\{s\&x\}) \ t = - \$p_{-}\{t\&x\} * \$\mu_{-(x+t)}$
if ($\lambda s. \$p_{-}\{s\&x\}$ *differentiable at t t > 0 for t::real*)
proof –
let $?cdfTx = \text{cdf } (\text{distr } (\mathfrak{M} \mid \text{alive } x) \text{ borel } (T \ x))$
have *diff: ?cdfTx differentiable at t*
using *that distrTx-RD.differentiable-cdf-ccdf unfolding survive-def by blast*
hence *deriv ?cdfTx t = \$p_{-}\{t\&x\} * \$μ_{-(x+t)}* **using** *that pdfTx-p-mu pdfT-def*
by simp
moreover have *deriv ?cdfTx t = - deriv (λs. \$p_{-}\{s\&x\}) t*
using *that diff distrTx-RD.deriv-cdf-ccdf unfolding survive-def by presburger*
ultimately show *?thesis by simp*
qed

lemma *pdfTx-p-mu-AE*: *AE s in lborel. s > 0 → pdfT x s = \$p_{-}\{s\&x\} * \$μ_{-(x+s)}*

proof –

let $?cdfX = cdf (distr \mathfrak{M} \text{ borel } X)$
let $?cdfTx = cdf (distr (\mathfrak{M} \downarrow \text{alive } x) \text{ borel } (T x))$
from $cdfX\text{-differentiable-AE}$ **obtain** N
where $diff: \bigwedge r. r \in \text{space } lborel - N \implies ?cdfX \text{ differentiable at } r$
and $null: N \in \text{null-sets } lborel$
using $AE-E3$ **by** $blast$
let $?N' = \{s. x+s \in N\}$
have $?N' \in \text{null-sets } lborel$
using $\text{null-sets-translation}[of N - x, \text{simplified}] \text{ null}$ **by** $(simp \text{ add: add.commute})$
moreover **have** $\{s \in \text{space } lborel. \neg (s > 0 \longrightarrow pdfT x s = \int p\{s \& x\} * \mu_{-(x+s)})\} \subseteq ?N'$
proof $(\text{rewrite Compl-subset-Compl-iff}[THEN sym], \text{safe})$
fix $s::real$
assume $s \in \text{space } lborel$ **and** $x+s \notin N$ **and** $s > 0$
thus $pdfT x s = \int p\{s \& x\} * \mu_{-(x+s)}$
apply $(\text{intro } pdfTx\text{-p-mu}, \text{simp-all})$
by $(\text{rewrite differentiable-cdfX-cdfTx}[THEN sym]; \text{simp add: diff})$
qed
ultimately show $?thesis$ **using** $AE-I'$ $[of ?N']$ **by** $simp$
qed

lemma $LBINT\text{-p-mu-q-defer}: (LBINT s:\{f<..f+t\}. \int p\{s \& x\} * \mu_{-(x+s)}) = \int q\{f|t \& x\}$
if $t \geq 0$ **for** $t f :: real$
proof –
have $(LBINT s:\{f<..f+t\}. \int p\{s \& x\} * \mu_{-(x+s)}) = (LBINT s:\{f<..f+t\}. pdfT x s)$
apply $(\text{rule set-lebesgue-integral-cong-AE}; \text{simp})$
apply $(\text{simp add: survive-def})$
using $pdfTx\text{-p-mu-AE}$ **apply** (rule AE-mp)
using that **by** $(\text{intro always-eventually}; \text{simp add: ereal-less-le})$
also **have** $\dots = \text{enn2real} (\int^{+s \in \{f<..f+t\}. \text{ennreal } (pdfT x s) \partial lborel})$
proof –
have $(\int^{+s \in \{f<..f+t\}. \text{ennreal } (pdfT x s) \partial lborel}) < \top$
proof –
have $(\int^{+s \in \{f<..f+t\}. \text{ennreal } (pdfT x s) \partial lborel}) \leq \int^{+s. \text{ennreal } (pdfT x s) \partial lborel}$
by $(\text{smt (verit) indicator-simps le-zero-eq linorder-le-cases mult.right-neutral mult-zero-right nn-integral-mono})$
also **have** $\dots < \top$ **using** $\text{nn-integral-pdfTx-1}$ **by** $simp$
finally show $?thesis .$
qed
thus $?thesis$
unfolding $\text{set-lebesgue-integral-def}$
by $(\text{rewrite enn2real-nn-integral-eq-integral}[\text{where } g = \lambda s. pdfT x s * \text{indicator } \{f<..f+t\} s])$
 $(\text{simp-all add: pdfTx-nonneg mult.commute ennreal-indicator ennreal-mult}')$
qed
also **have** $\dots = \text{measure } (density \text{ borel } (pdfT x)) \{f<..f+t\}$

unfolding *measure-def* **by** (*rewrite emeasure-density; simp*)
also have $\dots = \text{measure } (\text{distr } (\mathfrak{M} \mid \text{alive } x) \text{ lborel } (T \ x)) \{f < ..f+t\}$
using *distributed-pdfTx unfolding distributed-def by simp*
also have $\dots =$
 $\text{cdf } (\text{distr } (\mathfrak{M} \mid \text{alive } x) \text{ lborel } (T \ x)) (f+t) - \text{cdf } (\text{distr } (\mathfrak{M} \mid \text{alive } x) \text{ lborel } (T \ x)) f$
using *that finite-borel-measure.cdf-diff-eq2*
by (*smt (verit) distrTx-RD.finite-borel-measure-axioms distr-cong sets-lborel*)
also have $\dots = \$q\{f|t\&x\}$
using *q-defer-q die-def that by (metis distr-cong sets-lborel x-lt-psi)*
finally show *?thesis* .
qed

corollary *LBINT-p-mu-q*: (*LBINT* $s:\{0 < ..t\}$. $\$p\{s\&x\} * \$\mu\text{-}(x+s) = \$q\{t\&x\}$
if $t \geq 0$ **for** $t :: \text{real}$
using *LBINT-p-mu-q-defer that by force*

lemma *set-integrable-p-mu*: *set-integrable lborel* $\{f < ..f+t\}$ ($\lambda s. \$p\{s\&x\} * \$\mu\text{-}(x+s)$)
if $t \geq 0$ **for** $t f :: \text{real}$
proof –
have $(\lambda s. \$p\{s\&x\}) \in \text{borel-measurable borel}$ **unfolding** *survive-def by simp*
moreover have *AE* s *in lborel*. $f < s \wedge s \leq f + t \longrightarrow \$p\{s\&x\} * \$\mu\text{-}(x+s) =$
 $\text{pdfT } x \ s$
using *pdfTx-p-mu-AE apply (rule AE-mp)*
using *that by (intro always-eventually; simp add: ereal-less-le)*
moreover have *set-integrable lborel* $\{f < ..f+t\}$ (*pdfT* x) **using** *pdfTx-set-integrable*
by *simp*
ultimately show *?thesis by (rewrite set-integrable-cong-AE[where g=pdfT x]; simp)*
qed

lemma *p-mu-has-integral-q-defer-Ioc*:
 $((\lambda s. \$p\{s\&x\} * \$\mu\text{-}(x+s)) \text{ has-integral } \$q\{f|t\&x\}) \{f < ..f+t\}$
if $t \geq 0$ **for** $t f :: \text{real}$
apply (*rewrite LBINT-p-mu-q-defer[THEN sym], simp-all add: that*)
apply (*rewrite set-borel-integral-eq-integral, simp add: set-integrable-p-mu that*)
by (*rewrite has-integral-integral[THEN sym]; simp add: set-borel-integral-eq-integral set-integrable-p-mu that*)

lemma *p-mu-has-integral-q-defer-Icc*:
 $((\lambda s. \$p\{s\&x\} * \$\mu\text{-}(x+s)) \text{ has-integral } \$q\{f|t\&x\}) \{f ..f+t\}$ **if** $t \geq 0$ **for**
 $t f :: \text{real}$
proof –
have *negligible* $\{f\}$ **by** *simp*
hence [*simp*]: *negligible* $(\{f ..f+t\} - \{f < ..f+t\})$
by (*smt (verit) Diff-iff atLeastAtMost-iff greaterThanAtMost-iff insertI1 negligible-sing negligible-subset subsetI*)
have [*simp*]: *negligible* $(\{f < ..f+t\} - \{f ..f+t\})$ **by** (*simp add: subset-eq*)
show *?thesis*

apply (*rewrite has-integral-spike-set-eq*[**where** $T = \{f < ..f + t\}$])
apply (*rule negligible-subset*[of $\{f..f + t\} - \{f < ..f + t\}$], *simp*, *blast*)
apply (*rule negligible-subset*[of $\{f < ..f + t\} - \{f..f + t\}$], *simp*, *blast*)
using *p-mu-has-integral-q-defer-Ioc* that **by** *simp*
qed

corollary *p-mu-has-integral-q-Icc*:
 $((\lambda s. \$p\{s \& x\} * \$\mu\text{-}(x+s)) \text{ has-integral } \$q\{t \& x\}) \{0..t\}$ **if** $t \geq 0$ **for** $t::\text{real}$
using *p-mu-has-integral-q-defer-Icc*[**where** $f=0$] that **by** *simp*

corollary *p-mu-integrable-on-Icc*:
 $(\lambda s. \$p\{s \& x\} * \$\mu\text{-}(x+s))$ *integrable-on* $\{0..t\}$ **if** $t \geq 0$ **for** $t::\text{real}$
using *p-mu-has-integral-q-Icc* that **by** *auto*

lemma *e-ennreal-p-mu*: $(\int^+ \xi. \text{ennreal } (T \ x \ \xi) \ \partial(\mathfrak{M} \ \downarrow \ \text{alive } x)) =$
 $(\int^+ s \in \{0..\}. \text{ennreal } (\$p\{s \& x\} * \$\mu\text{-}(x+s) * s) \ \partial \text{lborel})$

proof –
have [*simp*]: *sym-diff* $\{0..\} \{0 < ..\} = \{0::\text{real}\}$ **by** *force*
have $(\int^+ \xi. \text{ennreal } (T \ x \ \xi) \ \partial(\mathfrak{M} \ \downarrow \ \text{alive } x)) = (\int^+ s \in \{0..\}. \text{ennreal } (\text{pdfT } x \ s * s) \ \partial \text{lborel})$
by (*rewrite nn-integral-T-pdfT*[**where** $g = \lambda s. s$]; *simp*)
also have $\dots = (\int^+ s \in \{0 < ..\}. \text{ennreal } (\text{pdfT } x \ s * s) \ \partial \text{lborel})$
by (*rewrite nn-integral-null-delta*; *force*)
also have $\dots = (\int^+ s \in \{0 < ..\}. \text{ennreal } (\$p\{s \& x\} * \$\mu\text{-}(x+s) * s) \ \partial \text{lborel})$
apply (*rule nn-integral-cong-AE*)
using *pdfTx-p-mu-AE* **apply** (*rule AE-mp*, *intro AE-I2*) **by** *force*
also have $\dots = (\int^+ s \in \{0..\}. \text{ennreal } (\$p\{s \& x\} * \$\mu\text{-}(x+s) * s) \ \partial \text{lborel})$
by (*rewrite nn-integral-null-delta*[*THEN sym*]; *force*)
finally show *?thesis* .
qed

lemma *e-LBINT-p-mu*: $\$e' \circ \text{-} x = (\text{LBINT } s: \{0..\}. \$p\{s \& x\} * \$\mu\text{-}(x+s) * s)$
— Note that $0 = 0$ holds when the life expectation diverges.

proof –
let $?f = \lambda s. \$p\{s \& x\} * \$\mu\text{-}(x+s) * s$
have [*simp*]: $(\lambda s. ?f \ s * \text{indicat-real } \{0..\} \ s) \in \text{borel-measurable borel}$
by *measurable (simp-all add: survive-def)*
hence [*simp*]: $(\lambda s. \text{indicat-real } \{0..\} \ s * ?f \ s) \in \text{borel-measurable borel}$
by (*meson measurable-cong mult.commute*)
have [*simp*]: *AE* s *in* *lborel*. $?f \ s * \text{indicat-real } \{0..\} \ s \geq 0$
proof –
have *AE* s *in* *lborel*. $\text{pdfT } x \ s * s * \text{indicat-real } \{0..\} \ s \geq 0$
using *pdfTx-nonneg pdfTx-nonpos-0 x-lt-psi*
by (*intro AE-I2*) (*smt (verit, del-insts) indicator-pos-le mult-eq-0-iff mult-nonneg-nonneg*)
thus *?thesis*
apply (*rule AE-mp*)
using *pdfTx-p-mu-AE* **apply** (*rule AE-mp*)
by (*rule AE-I2*) (*metis atLeast-iff indicator-simps(2) mult-eq-0-iff order-le-less*)
qed

hence $[simp]$: $AE\ s\ \text{in}\ lborel.\ \text{indicat-real}\ \{0..\}\ s * ?f\ s \geq 0$
by $(metis\ (no-types,\ lifting)\ AE-cong\ mult.commute)$
show $?thesis$
proof $(cases\ \langle integrable\ (\mathfrak{M}\ \downarrow\ alive\ x)\ (T\ x) \rangle)$
case $True$
hence $ennreal\ (\$e^{\circ-x}) = \int^{+}\xi.\ ennreal\ (T\ x\ \xi)\ \partial(\mathfrak{M}\ \downarrow\ alive\ x)$
unfolding $life-expect-def$ **apply** $(rewrite\ nn-integral-eq-integral,\ simp-all)$
by $(smt\ (verit)\ AE-I2\ alivex-Tx-pos)$
also have $\dots = \int^{+} s.\ ennreal\ (?f\ s * \text{indicat-real}\ \{0..\}\ s)\ \partial lborel$
by $(simp\ add:\ nn-integral-set-ennreal\ e-ennreal-p-mu)$
also have $\dots = ennreal\ (LBINT\ s:\{0..\}\ .\ ?f\ s)$
proof $-$
have $integrable\ lborel\ (\lambda s.\ ?f\ s * \text{indicat-real}\ \{0..\}\ s)$
proof $-$
have $(\int^{+} s.\ ennreal\ (?f\ s * \text{indicat-real}\ \{0..\}\ s)\ \partial lborel) < \infty$
by $(metis\ calculation\ ennreal-less-top\ infinity-ennreal-def)$
thus $?thesis$ **by** $(intro\ integrableI-nonneg;\ simp)$
qed
thus $?thesis$
unfolding $set-lebesgue-integral-def$
by $(rewrite\ nn-integral-eq-integral,\ simp-all)\ (meson\ mult.commute)$
qed
finally have $ennreal\ (\$e^{\circ-x}) = ennreal\ (LBINT\ s:\{0..\}\ .\ ?f\ s)$.
moreover have $(LBINT\ s:\{0..\}\ .\ ?f\ s) \geq 0$
unfolding $set-lebesgue-integral-def$ **by** $(rule\ integral-nonneg-AE)\ simp$
ultimately show $?thesis$ **using** $e-nonneg$ **by** $simp$
next
case $False$
hence $\$e^{\circ-x} = 0$ **unfolding** $life-expect-def$ **using** $not-integrable-integral-eq$ **by**
 $force$
also have $\dots = (LBINT\ s:\{0..\}\ .\ ?f\ s)$
proof $-$
have $\infty = \int^{+}\xi.\ ennreal\ (T\ x\ \xi)\ \partial(\mathfrak{M}\ \downarrow\ alive\ x)$
using $nn-integral-nonneg-infinite\ False$
by $(smt\ (verit)\ AE-cong\ Tx-alivex-measurable\ alivex-PS.AE-prob-1\ alivex-PS.prob-space\ alivex-Tx-pos\ nn-integral-cong)$
hence $0 = enn2real\ (\int^{+} s \in \{0..\}\ .\ ennreal\ (?f\ s)\ \partial lborel)$ **using** $e-ennreal-p-mu$
by $simp$
also have $\dots = (LBINT\ s:\{0..\}\ .\ ?f\ s)$
unfolding $set-lebesgue-integral-def$ **apply** $(rewrite\ integral-eq-nn-integral,\ simp-all)$
by $(simp\ add:\ indicator-mult-ennreal\ mult.commute)$
finally show $?thesis$ **by** $simp$
qed
finally show $?thesis$.
qed
qed

lemma $e\text{-integral-p-mu}$: $\$e^{\circ-x} = \text{integral}\ \{0..\}\ (\lambda s.\ \$p\ \{s\&x\} * \$\mu\ -(x+s) * s)$

— Note that $0 = 0$ holds when the life expectation diverges.

proof –

have $(LBINT\ s:\{0..\}. \$p-\{s\&x\} * \$\mu-(x+s) * s) = integral\ \{0..\} (\lambda s. \$p-\{s\&x\} * \$\mu-(x+s) * s)$

proof –

have $AE\ s\ in\ lborel.\ s \geq 0 \longrightarrow \$p-\{s\&x\} * \$\mu-(x+s) * s \geq 0$

proof –

have $AE\ s\ in\ lborel.\ \$\mu-(x+s) \geq 0$ **by** $(rule\ AE\text{-translation},\ rule\ mu\text{-nonneg}\text{-}AE)$

with $p\text{-nonneg}$ **show** $?thesis$ **by** $force$

qed

moreover **have** $(\lambda s. \$p-\{s\&x\} * \$\mu-(x+s) * s) \in borel\text{-measurable}\ borel$

unfolding $survive\text{-def}$ **by** $simp$

ultimately **show** $?thesis$ **by** $(intro\ set\text{-borel}\text{-integral}\text{-eq}\text{-integral}\text{-nonneg}\text{-}AE;\ simp)$

qed

thus $?thesis$ **using** $e\text{-LBINT}\text{-}p\text{-}\mu$ **by** $simp$

qed

end

lemma $p\text{-has}\text{-real}\text{-derivative}\text{-}x\text{-cdf}X$:

$((\lambda y. \$p-\{t\&y\})\ has\text{-real}\text{-derivative}\ ((deriv\ svl\ (x+t) * svl\ x - svl\ (x+t) * deriv\ svl\ x) / (svl\ x)^2))\ (at\ x)$

if $svl \equiv cdf\ (distr\ \mathfrak{M}\ borel\ X)\ svl$ **differentiable** **at** x **svl** **differentiable** **at** $(x+t)$

$t \geq 0\ x < \$\psi$ **for** $x\ t :: real$

proof –

have $open\ \{y.\ svl\ \text{differentiable}\ \text{at}\ y\}$ **using** $cdfX\text{-differentiable}\text{-open}\text{-set}$ **that** **by** $simp$

with $that$ **obtain** $e0$ **where** $e0\text{-pos}:\ e0 > 0$ **and** $ball\text{-}e0:\ ball\ x\ e0 \subseteq \{y.\ svl\ \text{differentiable}\ \text{at}\ y\}$

using $open\text{-contains}\text{-ball}$ **by** $blast$

define e **where** $e \equiv if\ \$\psi < \infty$ **then** $min\ e0\ (real\text{-of}\text{-ereal}\ \$\psi - x)$ **else** $e0$

have $e > 0 \wedge ball\ x\ e \subseteq \{y.\ svl\ y \neq 0 \wedge svl\ \text{differentiable}\ \text{at}\ y\}$

proof $(cases\ \langle \$\psi < \infty \rangle)$

case $True$

hence $e > 0$

proof –

have $real\text{-of}\text{-ereal}\ \$\psi - x > 0$ **using** $not\text{-infty}I\ True$ **that** **by** $force$

with $e0\text{-pos}$ **show** $?thesis$ **unfolding** $e\text{-def}$ **using** $True$ **by** $simp$

qed

moreover **have** $ball\ x\ e \subseteq \{y.\ svl\ y \neq 0\}$

proof –

have $e \leq real\text{-of}\text{-ereal}\ \$\psi - x$ **unfolding** $e\text{-def}$ **using** $True$ **by** $simp$

hence $ball\ x\ e \subseteq \{..\ < real\text{-of}\text{-ereal}\ \$\psi\}$ **unfolding** $ball\text{-def}\ dist\text{-real}\text{-def}$ **by** $force$

thus $?thesis$ **using** $that\ cdfX\text{-}0\text{-equiv}$

using $True\ ereal\text{-less}\text{-real}\text{-iff}\ psi\text{-nonneg}$ **by** $auto$

qed

moreover **have** $ball\ x\ e \subseteq \{y.\ svl\ \text{differentiable}\ \text{at}\ y\}$

proof –
have $e \leq e0$ **unfolding** $e\text{-def}$ **using** $True$ **by** $simp$
hence $ball\ x\ e \subseteq ball\ x\ e0$ **by** $force$
with $ball\text{-}e0$ **show** $?thesis$ **by** $simp$
qed
ultimately show $?thesis$ **by** $blast$
next
case $False$
hence $ball\ x\ e \subseteq \{y. svl\ y \neq 0\}$ **using** $ccdfX\text{-}0\text{-equiv}\ that$ **by** $simp$
with $False$ **show** $?thesis$ **unfolding** $e\text{-def}$ **using** $e0\text{-pos}\ ball\text{-}e0$ **by** $force$
qed
hence $e\text{-pos}: e > 0$ **and** $ball\text{-}e: ball\ x\ e \subseteq \{y. svl\ y \neq 0 \wedge svl\ \text{differentiable at } y\}$
by $auto$
hence $ball\text{-}svl: \bigwedge y. y \in ball\ x\ e \implies svl\ y \neq 0 \wedge svl\ \text{field-differentiable at } y$
using $differentiable\text{-}eq\text{-field-differentiable-real}$ **by** $blast$
hence $\bigwedge y. y \in ball\ x\ e \implies \$p\text{-}\{t\&y\} = svl\ (y+t) / svl\ y$
unfolding $survive\text{-}def$ **using** $that\ ccdfX\text{-}0\text{-equiv}$ **by** $(rewrite\ ccdfTx\text{-}ccdfX,$
 $simp\text{-}all)$ $force$
moreover from $ball\text{-}svl$ **have** $((\lambda y. svl\ (y+t) / svl\ y)\ \text{has-real-derivative}$
 $((deriv\ svl\ (x+t) * svl\ x - svl\ (x+t) * deriv\ svl\ x) / (svl\ x)^2))\ (at\ x)$
apply $(rewrite\ power2\text{-}eq\text{-square)}$
apply $(rule\ DERIV\text{-}divide)$
using $DERIV\text{-}deriv\text{-}iff\text{-}real\text{-}differentiable\ DERIV\text{-}shift\ that$ **apply** $blast$
using $that\ DERIV\text{-}deriv\text{-}iff\text{-}real\text{-}differentiable$ **apply** $simp$
by $(simp\ add: e\text{-pos})$
ultimately show $?thesis$
using $e\text{-pos}$
apply $(rewrite\ has\text{-}field\text{-}derivative\text{-}cong\text{-}eventually[\text{where } g = \lambda y. svl\ (y+t) /$
 $svl\ y],\ simp\text{-}all)$
by $(smt\ (verit)\ dist\text{-}commute\ eventually\text{-}at)$
qed

lemma $p\text{-has-real-derivative-x-p-mu}$:
 $((\lambda y. \$p\text{-}\{t\&y\})\ \text{has-real-derivative}\ \$p\text{-}\{t\&x\} * (\$mu\text{-}x - \$mu\text{-}(x+t)))\ (at\ x)$
if $ccdf\ (distr\ \mathfrak{M}\ borel\ X)$ $\text{differentiable at } x$ $ccdf\ (distr\ \mathfrak{M}\ borel\ X)$ differentiable
 $\text{at } (x+t)$
 $t \geq 0\ x < \$\psi$ **for** $x\ t :: real$
proof $(cases\ \langle x+t < \$\psi \rangle)$
case $True$
let $?svl = ccdf\ (distr\ \mathfrak{M}\ borel\ X)$
have $[simp]: ?svl\ x \neq 0$ **using** $that\ ccdfX\text{-}0\text{-equiv}$ **by** $(smt\ (verit)\ le\text{-}ereal\text{-}le$
 $linorder\text{-}not\text{-}le)$
have $[simp]: ?svl\ \text{field-differentiable at } (x+t)$
using $that\ differentiable\text{-}eq\text{-field-differentiable-real}$ **by** $simp$
have $((\lambda y. \$p\text{-}\{t\&y\})\ \text{has-real-derivative}$
 $((deriv\ ?svl\ (x+t) * ?svl\ x - ?svl\ (x+t) * deriv\ ?svl\ x) / (?svl\ x)^2))\ (at\ x)$
using $p\text{-has-real-derivative-x-ccdfX}$ **that** **by** $simp$
moreover have $(deriv\ ?svl\ (x+t) * ?svl\ x - ?svl\ (x+t) * deriv\ ?svl\ x) / (?svl$
 $x)^2 =$

$\$p\text{-}\{t\&x\} * (\$μ\text{-}x - \$μ\text{-}(x+t))$ (is $?LHS = ?RHS$)
proof –
have $deriv\ ?svl\ (x+t) = deriv\ (\lambda y. ?svl\ (y+t))\ x$
using that **by** (metis *DERIV-deriv-iff-real-differentiable DERIV-imp-deriv DERIV-shift*)
hence $?LHS = (deriv\ (\lambda y. ?svl\ (y+t))\ x * ?svl\ x - ?svl\ (x+t) * deriv\ ?svl\ x) / (?svl\ x)^2$
by *simp*
also have $\dots = (deriv\ (\lambda y. ?svl\ (y+t))\ x - ?svl\ (x+t) * deriv\ ?svl\ x / ?svl\ x) / ?svl\ x$
by (*simp add: add-divide-eq-if-simps(4) power2-eq-square*)
also have $\dots = (- ?svl\ (x+t) * \$μ\text{-}(x+t) + ?svl\ (x+t) * \$μ\text{-}x) / ?svl\ x$
proof –
have $deriv\ (\lambda y. ?svl\ (y+t))\ x = - ?svl\ (x+t) * \$μ\text{-}(x+t)$
apply (*rewrite add.commute, rewrite deriv-shift[THEN sym], rewrite add.commute, simp*)
using *add.commute* that **by** (metis *MM-PS.deriv-ccdf-hazard-rate X-RV force-mortal-def*)
moreover have $- ?svl\ (x+t) * deriv\ ?svl\ x / ?svl\ x = ?svl\ (x+t) * \$μ\text{-}x$
using that **by** (*simp add: MM-PS.deriv-ccdf-hazard-rate force-mortal-def*)
ultimately show *?thesis* **by** *simp*
qed
also have $\dots = ?svl\ (x+t) * (\$μ\text{-}x - \$μ\text{-}(x+t)) / ?svl\ x$ **by** (*simp add: mult-diff-mult*)
also have $\dots = ?RHS$ **unfolding** *survive-def* **using** *ccdfTx-ccdfX* that **by** *simp*
ultimately show *?thesis* **by** *simp*
qed
ultimately show *?thesis* **by** *simp*
next
case *False*
hence *ptx-0*: $\$p\text{-}\{t\&x\} = 0$ **using** *p-0-equiv* that **by** *simp*
moreover have $((\lambda y. \$p\text{-}\{t\&y\})\ \text{has-real-derivative}\ 0)$ (at x)
proof –
have $\bigwedge y. x < y \implies y < \$\psi \implies \$p\text{-}\{t\&y\} = 0$
using *False p-0-equiv* that **by** (*smt (verit, ccfv-SIG) ereal-less-le linorder-not-le*)
hence $\forall_F x$ in *at-right* $x. \$p\text{-}\{t\&x\} = 0$
apply (*rewrite eventually-at-right-field*)
using that **by** (*meson ereal-dense2 ereal-le-less less-eq-ereal-def less-ereal.simps*)
hence $((\lambda y. \$p\text{-}\{t\&y\})\ \text{has-real-derivative}\ 0)$ (at-right x)
using *ptx-0* **by** (*rewrite has-field-derivative-cong-eventually[where g= $\lambda\cdot. 0$]; simp*)
thus *?thesis*
using *vector-derivative-unique-within p-has-real-derivative-x-ccdfX* that
by (*metis has-field-derivative-at-within has-real-derivative-iff-has-vector-derivative trivial-limit-at-right-real*)
qed
ultimately show *?thesis* **by** *simp*
qed

corollary *deriv-x-p-mu*: $\text{deriv } (\lambda y. \mathbb{P}\{t \& y\}) x = \mathbb{P}\{t \& x\} * (\mathbb{P}\mu - x - \mathbb{P}\mu - (x+t))$
if *ccdf (distr \mathfrak{M} borel X) differentiable at x ccdf (distr \mathfrak{M} borel X) differentiable at $(x+t)$*
 $t \geq 0 \ x < \mathbb{P}\psi$ **for** $x \ t :: \text{real}$
using *that p-has-real-derivative-x-p-mu DERIV-imp-deriv by blast*

lemma *e-has-derivative-mu-e*: $((\lambda x. \mathbb{P}e^{\circ-x}) \text{ has-real-derivative } (\mathbb{P}\mu - x * \mathbb{P}e^{\circ-x} - 1))$ (at x)
if $\bigwedge x. x \in \{a <.. < b\} \implies \text{set-integrable lborel } \{x..\}$ (ccdf (distr \mathfrak{M} borel X))
ccdf (distr \mathfrak{M} borel X) differentiable at $x \ x \in \{a <.. < b\} \ b \leq \mathbb{P}\psi$
for $a \ b \ x :: \text{real}$

proof –

let $?svl = \text{ccdf (distr } \mathfrak{M} \text{ borel } X)$
have $x\text{-lt-psi}[simp]$: $x < \mathbb{P}\psi$ **using** *that ereal-le-less by simp*
hence $svlx\text{-neq0}[simp]$: $?svl \ x \neq 0$ **by** (*simp add: ccdfX-0-equiv linorder-not-le*)
define $d :: \text{real}$ **where** $d \equiv \min (b-x) (x-a)$
have $d\text{-pos}$: $d > 0$ **unfolding** $d\text{-def}$ **using** *that ereal-less-real-iff by force*
have $e\text{-svl}$: $\bigwedge y. y < \mathbb{P}\psi \implies \mathbb{P}e^{\circ-y} = (\text{LBINT } t:\{0..\}. ?svl (y+t)) / ?svl \ y$
apply (*rewrite e-LBINT-p, simp*)
apply (*rewrite set-integral-divide-zero[THEN sym]*)
apply (*rule set-lebesgue-integral-cong, simp-all*)
unfolding *survive-def* **using** ccdfTx-ccdfX **by** *force*
have $((\lambda y. \text{LBINT } t:\{0..\}. ?svl (y+t)) \text{ has-real-derivative } (- ?svl \ x))$ (at x)

proof –

{ **fix** y **assume** $\text{dist } y \ x < d$
hence $y\text{-ab}$: $y \in \{a <.. < b\}$ **unfolding** $d\text{-def}$ *dist-real-def* **by** *force*
hence *set-integrable lborel* $\{y..\}$ $?svl$ **using** *that by simp*
hence *set-integrable lborel* $(\text{einterval } y \ \infty)$ $?svl$
by (*rewrite set-integrable-discrete-difference[where $X = \{y\}$]; simp*) *force*
moreover **have** $\bigwedge u. ((\lambda u. u-y) \text{ has-real-derivative } (1-0))$ (at u)
by (*rule derivative-intros*) *auto*
moreover **have** $\bigwedge u. y < u \implies \text{isCont } (\lambda t. ?svl (y+t)) (u-y)$
apply (*rewrite add.commute, rewrite isCont-shift, simp*)
using *ccdfX-continuous continuous-on-eq-continuous-within* **by** *blast*
moreover **have** $((\text{ereal} \circ (\lambda u. u-y) \circ \text{real-of-ereal}) \longrightarrow 0)$ (at-right (ereal

y))

by (*smt (verit, ccfv-SIG) LIM-zero Lim-cong-within ereal-tendsto-simps1(2) ereal-tendsto-simps2(1) tendsto-ident-at zero-ereal-def*)
moreover **have** $((\text{ereal} \circ (\lambda u. u-y) \circ \text{real-of-ereal}) \longrightarrow \infty)$ (at-left ∞)

proof –

have $\bigwedge r \ u. r+y+1 \leq u \implies r < u-y$ **by** *auto*
hence $\bigwedge r. \forall_F \ u \text{ in at-top. } r < u - y$ **by** (*rule eventually-at-top-linorderI*)
thus *?thesis* **by** (*rewrite ereal-tendsto-simps, rewrite tendsto-PInfty*) *simp*
qed

ultimately **have** $(\text{LBINT } t=0..\infty. ?svl (y+t)) = (\text{LBINT } u=y..\infty. ?svl (y+(u-y)) * 1)$

using *distrX-RD.ccdf-nonneg* **by** (*intro interval-integral-substitution-nonneg(2); simp*)

moreover have $(LBINT\ t:\{0..\}. \ ?svl\ (y+t)) = (LBINT\ t=0..\infty. \ ?svl\ (y+t))$
unfolding *interval-lebesgue-integral-def einterval-def* **apply** *simp*
by *(rule set-integral-discrete-difference[where X={0}]; force)*
moreover have $(LBINT\ u=y..\infty. \ ?svl\ (y+(u-y)*1)) = (LBINT\ u:\{y..\}. \ ?svl\ u)$
unfolding *interval-lebesgue-integral-def einterval-def* **apply** *simp*
by *(rule set-integral-discrete-difference[where X={y}]; force)*
ultimately have $(LBINT\ t:\{0..\}. \ ?svl\ (y+t)) = (LBINT\ u:\{y..\}. \ ?svl\ u)$ **by**
simp }
hence $\forall_F\ y\ \text{in}\ nhds\ x. (LBINT\ t:\{0..\}. \ ?svl\ (y+t)) = (LBINT\ u:\{y..\}. \ ?svl\ u)$
using *d-pos* **by** *(rewrite eventually-nhds-metric) auto*
moreover have $((\lambda y. LBINT\ u:\{y..\}. \ ?svl\ u)$ *has-real-derivative* $(-\ ?svl\ x))$
(at x)
proof –
have $((\lambda y. \text{integral}\ \{y..b\}\ \ ?svl)$ *has-real-derivative* $(-\ ?svl\ x))$ *(at x within*
 $\{a..b\})$
using *that* **apply** *(intro integral-has-real-derivative'; simp)*
using *ccdfX-continuous continuous-on-subset* **by** *blast*
hence $((\lambda y. \text{integral}\ \{y..b\}\ \ ?svl)$ *has-real-derivative* $(-\ ?svl\ x))$ *(at x)*
using *that* **apply** *(rewrite at-within-open[where S={a<..**b**}, THEN sym],*
simp-all)
by *(rule DERIV-subset[where s={a..b}]) auto*
moreover have $\forall_F\ y\ \text{in}\ nhds\ x. (LBINT\ u:\{y..b\}. \ ?svl\ u) = \text{integral}\ \{y..b\}$
 $\ ?svl$
apply *(rewrite eventually-nhds-metric)*
using *d-pos* **by** *(metis ccdfX-integrable-Icc set-borel-integral-eq-integral(2))*
ultimately have $((\lambda y. LBINT\ u:\{y..b\}. \ ?svl\ u)$ *has-real-derivative* $(-\ ?svl$
 $x))$ *(at x)*
by *(rewrite DERIV-cong-ev; simp)*
hence $((\lambda y. (LBINT\ u:\{y..b\}. \ ?svl\ u) + (LBINT\ u:\{b<..\}. \ ?svl\ u))$ *has-real-derivative*
 $(-\ ?svl\ x))$ *(at x)*
by *(rewrite to - ?svl x + 0 add-0-right[THEN sym], rule DERIV-add; simp)*
moreover have $\forall_F\ y\ \text{in}\ nhds\ x.$
 $(LBINT\ u:\{y..\}. \ ?svl\ u) = (LBINT\ u:\{y..b\}. \ ?svl\ u) + (LBINT\ u:\{b<..\}. \ ?svl\ u)$
 $\ ?svl\ u)$
proof –
{ fix y assume $dist\ y\ x < d$
hence $y-ab: y \in \{a<..**b**\}$ **unfolding** *d-def dist-real-def* **by** *force*
hence *set-integrable* *lborel* $\{y..\}$ $\ ?svl$ **using** *that* **by** *simp*
hence *set-integrable* *lborel* $\{y..b\}$ $\ ?svl$ *set-integrable* *lborel* $\{b<..\}$ $\ ?svl$
apply *(rule set-integrable-subset, simp-all)+*
using $y-ab$ **by** *force*
moreover have $\{y..b\} \cap \{b<..\} = \{\}$ $\{y..\} = \{y..b\} \cup \{b<..\}$ **using** $y-ab$
by *force+*
ultimately have
 $(LBINT\ u:\{y..\}. \ ?svl\ u) = (LBINT\ u:\{y..b\}. \ ?svl\ u) + (LBINT\ u:\{b<..\}. \ ?svl\ u)$
 $\ ?svl\ u)$
using *set-integral-Un* **by** *simp* }
thus *?thesis* **using** *d-pos* **by** *(rewrite eventually-nhds-metric) blast*

qed
ultimately show *?thesis* **by** (rewrite has-field-derivative-cong-ev; simp)
qed
ultimately show *?thesis* **by** (rewrite DERIV-cong-ev; simp)
qed
moreover have (*?svl has-real-derivative (deriv ?svl x)*) (at x)
using that DERIV-deriv-iff-real-differentiable **by** blast
ultimately have (($\lambda y. (LBINT t:\{0..\}. ?svl (y+t)) / ?svl y$) has-real-derivative
(($- ?svl x$) * *?svl x* - ($LBINT t:\{0..\}. ?svl (x+t)$) * *deriv ?svl x*) / (*?svl x* *
?svl x)) (at x)
by (rule DERIV-divide) simp
moreover have eventually ($\lambda y. (LBINT t:\{0..\}. ?svl (y+t)) / ?svl y = \$e^{\circ}y$)
(at x)
proof -
{ **fix** y **assume** $dist\ y\ x < d\ y \neq x$
hence $y < \$\psi$
unfolding dist-real-def d-def **using** that ereal-le-less **by** fastforce
hence $\$e^{\circ}y = (LBINT t:\{0..\}. ?svl (y+t)) / ?svl y$ **by** (rule e-svl) }
thus *?thesis*
apply (rewrite eventually-at-filter, rewrite eventually-nhds-metric)
using d-pos that **by** metis
qed
ultimately have (($\lambda y. \$e^{\circ}y$) has-real-derivative
(($- ?svl x$) * *?svl x* - ($LBINT t:\{0..\}. ?svl (x+t)$) * *deriv ?svl x*) / (*?svl x* *
?svl x)) (at x)
using e-svl **by** (rewrite has-field-derivative-cong-eventually[THEN sym]; simp)
moreover have
(($- ?svl x$) * *?svl x* - ($LBINT t:\{0..\}. ?svl (x+t)$) * *deriv ?svl x*) / (*?svl x* *
?svl x) =
 $\$ \mu \cdot x * \$e^{\circ}x - 1$ (**is** ?LHS = ?RHS)
proof -
have ?LHS = $-1 + (LBINT t:\{0..\}. ?svl (x+t)) / ?svl x * (-\ deriv\ ?svl\ x /$
?svl x)
by simp (smt (verit) svlx-neq0 add-divide-distrib divide-eq-minus-1-iff
mult-minus-left real-divide-square-eq)
also have ... = $-1 + \$e^{\circ}x * \$ \mu \cdot x$ **using** e-svl mu-deriv-ccdf that **by** force
also have ... = ?RHS **by** simp
finally show *?thesis* .
qed
ultimately show *?thesis* **by** simp
qed

corollary e-has-derivative-mu-e': (($\lambda x. \$e^{\circ}x$) has-real-derivative ($\$ \mu \cdot x * \$e^{\circ}x$
- 1)) (at x)
if $\bigwedge x. x \in \{a < .. < b\} \implies$ ccdf (distr \mathfrak{M} borel X) integrable-on {x..}
ccdf (distr \mathfrak{M} borel X) differentiable at x $x \in \{a < .. < b\}$ $b \leq \$\psi$
for a b x :: real
using that **apply** (intro e-has-derivative-mu-e[where a=a], simp-all)
using distrX-RD.ccdf-nonneg **by** (rewrite integrable-on-iff-set-integrable-nonneg;

simp)

5.2.11 Properties of Curtate Life Expectation

context

fixes $x::\text{real}$

assumes $x\text{-lt-}\psi[\textit{simp}]$: $x < \psi$

begin

lemma *isCont-p-nat*: $\textit{isCont} (\lambda t. \mathbb{P}\{t \leq x\}) (k + (1::\text{real}))$ **for** $k::\text{nat}$

proof –

fix $k::\text{nat}$

have *continuous-on* $\{0 < ..\}$ $(\lambda t. \mathbb{P}\{t \leq x\})$

unfolding *survive-def*

using *cdfTx-continuous-on-nonneg continuous-on-subset Ioi-le-Ico x-lt-psi* **by**

blast

hence $\forall t \in \{0 < ..\}. \textit{isCont} (\lambda t. \mathbb{P}\{t \leq x\}) t$

using *continuous-on-eq-continuous-at open-greaterThan* **by** *blast*

thus $\textit{isCont} (\lambda t. \mathbb{P}\{t \leq x\}) (\text{real } k+1)$ **by** *force*

qed

lemma *curt-e-sum-p-smooth*: $\mathbb{E}x = (\sum k. \mathbb{P}\{k+1 \leq x\})$ **if** *summable* $(\lambda k. \mathbb{P}\{k+1 \leq x\})$

using *curt-e-sum-p isCont-p-nat that* **by** *simp*

lemma *curt-e-rec-smooth*: $\mathbb{E}x = \mathbb{P}x * (1 + \mathbb{E}x(x+1))$ **if** *summable* $(\lambda k. \mathbb{P}\{k+1 \leq x\})$

$x+1 < \psi$

using *curt-e-rec isCont-p-nat that* **by** *simp*

lemma *curt-e-sum-p-finite-smooth*: $\mathbb{E}x = (\sum k < n. \mathbb{P}\{k+1 \leq x\})$ **if** $x+n+1 >$

ψ **for** $n::\text{nat}$

using *curt-e-sum-p-finite isCont-p-nat that* **by** *simp*

lemma *temp-curt-e-sum-p-smooth*: $\mathbb{E}\{x:n\} = (\sum k < n. \mathbb{P}\{k+1 \leq x\})$ **for** $n::\text{nat}$

using *temp-curt-e-sum-p isCont-p-nat* **by** *simp*

lemma *temp-curt-e-rec-smooth*: $\mathbb{E}\{x:n\} = \mathbb{P}x * (1 + \mathbb{E}\{x+1:n-1\})$

if $x+1 < \psi$ $n \neq 0$ **for** $n::\text{nat}$

using *temp-curt-e-rec isCont-p-nat that* **by** *simp*

end

end

5.3 Limited Survival Function

locale *limited-survival-function* = *survival-model* +

assumes *psi-limited[simp]*: $\psi < \infty$

begin

definition *ult-age* :: $\text{nat} (\mathbb{P}\omega)$

where $\$w \equiv \text{LEAST } x::\text{nat. cdf } (\text{distr } \mathfrak{M} \text{ borel } X) x = 0$
 — the conventional notation for ultimate age

lemma *ccdfX-ceil-psi-0*: $\text{ccdf } (\text{distr } \mathfrak{M} \text{ borel } X) \lceil \text{real-of-ereal } \$\psi \rceil = 0$
proof –
have $\text{real-of-ereal } \$\psi \leq \lceil \text{real-of-ereal } \$\psi \rceil$ **by** *simp*
thus *?thesis* **using** *ccdfX-0-equiv psi-limited cdfX-psi-0 le-ereal-le* **by** *presburger*
qed

lemma *ccdfX-omega-0*: $\text{ccdf } (\text{distr } \mathfrak{M} \text{ borel } X) \$w = 0$
unfolding *ult-age-def*
proof (*rule LeastI-ex*)
have $\lceil \text{real-of-ereal } \$\psi \rceil \geq 0$ **using** *psi-nonneg real-of-ereal-pos* **by** *fastforce*
thus $\exists x::\text{nat. cdf } (\text{distr } \mathfrak{M} \text{ borel } X) (\text{real } x) = 0$
using *ccdfX-ceil-psi-0* **by** (*metis of-int-of-nat-eq zero-le-imp-eq-int*)
qed

corollary *psi-le-omega*: $\$w \leq \ψ
using *ccdfX-0-equiv cdfX-omega-0* **by** *blast*

corollary *omega-pos*: $\$w > 0$
using *psi-le-omega order.strict-iff-not* **by** *fastforce*

lemma *omega-ceil-psi*: $\$w = \lceil \text{real-of-ereal } \$\psi \rceil$
proof (*rule antisym*)
let *?cpsi* = $\lceil \text{real-of-ereal } \$\psi \rceil$
have $\star: ?cpsi \geq 0$ **using** *psi-nonneg real-of-ereal-pos* **by** *fastforce*
moreover **have** $\text{ccdf } (\text{distr } \mathfrak{M} \text{ borel } X) ?cpsi = 0$ **by** (*rule cdfX-ceil-psi-0*)
ultimately **have** $\$w \leq \text{nat } ?cpsi$
unfolding *ult-age-def* **using** *wellorder-Least-lemma of-nat-nat* **by** *smt*
thus $\text{int } \$w \leq ?cpsi$ **using** *le-nat-iff* \star **by** *blast*
next
show $\lceil \text{real-of-ereal } \$\psi \rceil \leq \text{int } \w
using *psi-le-omega* **by** (*simp add: ceiling-le-iff real-of-ereal-ord-simps(2)*)
qed

lemma *ccdfX-0-equiv-nat*: $\text{ccdf } (\text{distr } \mathfrak{M} \text{ borel } X) x = 0 \iff x \geq \w **for** $x::\text{nat}$
proof
assume $\text{ccdf } (\text{distr } \mathfrak{M} \text{ borel } X) (\text{real } x) = 0$
thus $x \geq \$w$ **unfolding** *ult-age-def* **using** *wellorder-Least-lemma* **by** *fastforce*
next
assume $x \geq \$w$
hence $\text{ereal } (\text{real } x) \geq \ψ **using** *psi-le-omega le-ereal-le of-nat-mono* **by** *blast*
thus $\text{ccdf } (\text{distr } \mathfrak{M} \text{ borel } X) (\text{real } x) = 0$ **using** *ccdfX-0-equiv* **by** *simp*
qed

lemma *psi-le-iff-omega-le*: $\$w \leq x \iff \$w \leq x$ **for** $x::\text{nat}$
using *ccdfX-0-equiv cdfX-0-equiv-nat* **by** *presburger*

```

context
  fixes  $x::nat$ 
  assumes  $x\text{-lt-}\omega[simp]: x < \omega$ 
begin

lemma  $x\text{-lt-}\psi[simp]: x < \psi$ 
  using  $x\text{-lt-}\omega$   $\psi\text{-le-iff-}\omega\text{-le}$  by (meson linorder-not-less)

lemma  $p\text{-}0\text{-}1\text{-}nat: \mathbb{P}\{0 \& x\} = 1$ 
  using  $p\text{-}0\text{-}1$  by  $simp$ 

lemma  $p\text{-}0\text{-}equiv\text{-}nat: \mathbb{P}\{t \& x\} = 0 \iff x+t \geq \omega$  for  $t::nat$ 
  using  $\psi\text{-le-iff-}\omega\text{-le}$   $p\text{-}0\text{-}equiv$  by (metis of-nat-add x-lt-psi)

lemma  $q\text{-}0\text{-}0\text{-}nat: \mathbb{Q}\{0 \& x\} = 0$ 
  using  $p\text{-}q\text{-}1$   $p\text{-}0\text{-}1\text{-}nat$  by (smt (verit) x-lt-psi)

lemma  $q\text{-}1\text{-}equiv\text{-}nat: \mathbb{Q}\{t \& x\} = 1 \iff x+t \geq \omega$  for  $t::nat$ 
  using  $p\text{-}q\text{-}1$   $p\text{-}0\text{-}equiv\text{-}nat$  by (smt (verit) x-lt-psi)

lemma  $q\text{-}defer\text{-}old\text{-}0\text{-}nat: \mathbb{Q}\{f | t \& x\} = 0$  if  $\omega \leq x+f$  for  $f t :: nat$ 
  using  $q\text{-}defer\text{-}old\text{-}0$   $\psi\text{-le-iff-}\omega\text{-le}$  that by (metis of-nat-0-le-iff of-nat-add x-lt-psi)

lemma  $curt\text{-}e\text{-}sum\text{-}P\text{-}finite\text{-}nat: \mathbb{E}x = (\sum_{k < n} \mathcal{P}(\xi \text{ in } \mathfrak{M}. T x \xi \geq k + 1 \mid T x \xi > 0))$ 
  if  $x+n \geq \omega$  for  $n::nat$ 
  apply (rule curt-e-sum-P-finite, simp)
  using that  $\psi\text{-le-iff-}\omega\text{-le}$  by (smt (verit) le-ereal-less of-nat-add)

lemma  $curt\text{-}e\text{-}sum\text{-}p\text{-}finite\text{-}nat: \mathbb{E}x = (\sum_{k < n} \mathbb{P}\{k+1 \& x\})$ 
  if  $\bigwedge k::nat. k < n \implies isCont (\lambda t. \mathbb{P}\{t \& x\})$   $(real\ k + 1) x+n \geq \omega$  for  $n::nat$ 
  apply (rule curt-e-sum-p-finite, simp-all add: that)
  using that  $\psi\text{-le-iff-}\omega\text{-le}$  by (smt (verit) le-ereal-less of-nat-add)

end

lemma  $q\text{-}\omega\text{-}1: \mathbb{Q}(\omega-1) = 1$ 
  using  $q\text{-}1\text{-}equiv\text{-}nat$ 
  by (metis diff-less dual-order.refl le-diff-conv of-nat-1 omega-pos zero-less-one)

end

end
theory Life-Table
  imports Survival-Model
begin

```

6 Life Table

Define a life table axiomatically.

```

locale life-table =
  fixes l :: real  $\Rightarrow$  real ($l'-- [101] 200)
  assumes l-0-pos: 0 < l 0
    and l-neg-nil:  $\bigwedge x. x \leq 0 \implies l x = l 0$ 
    and l-PInfty-0: (l  $\longrightarrow$  0) at-top
    and l-antimono: antimono l
    and l-right-continuous:  $\bigwedge x. \text{continuous (at-right } x) l$ 
begin

```

6.1 Basic Properties of Life Table

```

lemma l-0-neq-0[simp]: $l-0  $\neq$  0
  using l-0-pos by simp

```

```

lemma l-nonneg[simp]: $l-x  $\geq$  0 for x::real
  using l-antimono l-PInfty-0 by (rule antimono-at-top-le)

```

```

lemma l-bounded[simp]: $l-x  $\leq$  $l-0 for x::real
  using l-neg-nil l-antimono by (smt (verit) antimonoD)

```

```

lemma l-measurable[measurable, simp]: l  $\in$  borel-measurable borel
  by (rule borel-measurable-antimono, rule l-antimono)

```

```

lemma l-left-continuous-nonpos: continuous (at-left x) l if x  $\leq$  0 for x::real
proof –

```

```

  have $l-x = $l-0 using l-neg-nil that by blast
  moreover have continuous (at-left x) ( $\lambda$ -. $l-0) by simp
  moreover have eventually ( $\lambda y. $l-y = $l-0$ ) (at-left x)
    apply (rule eventually-at-leftI[of x-1], simp-all)
    using that l-neg-nil by (smt (verit))
  ultimately show ?thesis by (rewrite continuous-at-within-cong[where g= $\lambda$ -.
  $l-0]; simp)
qed

```

```

lemma l-integrable-Icc: set-integrable lborel {a..b} l for a b :: real
  unfolding set-integrable-def
  apply (rule integrableI-bounded-set[where A={a..b} and B=$l-0], simp-all)
  using emeasure-compact-finite by auto

```

```

corollary l-integrable-on-Icc: l integrable-on {a..b} for a b :: real
  using l-integrable-Icc by (rewrite integrable-on-iff-set-integrable-nonneg[THEN
  sym]; simp)

```

```

lemma l-integrable-Icc-shift: set-integrable lborel {a..b} ( $\lambda t. $l-(x+t)$ ) for a b x ::
  real
  using set-integrable-Icc-shift l-integrable-Icc by (metis (full-types) add-diff-cancel-right')

```

corollary *l-integrable-on-Icc-shift*: $(\lambda t. \mathbb{I}l(x+t))$ *integrable-on* $\{a..b\}$ **for** $x a b :: \text{real}$
using *l-integrable-Icc-shift* **by** (*rewrite integrable-on-iff-set-integrable-nonneg*[*THEN sym*]; *simp*)

lemma *l-normal-antimono*: *antimono* $(\lambda x. \mathbb{I}l-x / \mathbb{I}l-0)$
using *divide-le-cancel l-0-pos l-antimono* **unfolding** *antimono-def* **by** *fastforce*

lemma *compl-l-normal-right-continuous*: *continuous* (*at-right* x) $(\lambda x. 1 - \mathbb{I}l-x / \mathbb{I}l-0)$ **for** $x :: \text{real}$
using *l-0-pos l-right-continuous* **by** (*intro continuous-intros*; *simp*)

lemma *compl-l-normal-NInfty-0*: $((\lambda x. 1 - \mathbb{I}l-x / \mathbb{I}l-0) \longrightarrow 0)$ *at-bot*
apply (*rewrite tendsto-cong*[**where** $g = \lambda \cdot. 0$], *simp-all*)
by (*smt* (*verit*) *div-self eventually-at-bot-linorder l-0-pos l-neg-nil*)

lemma *compl-l-normal-PInfty-1*: $((\lambda x. 1 - \mathbb{I}l-x / \mathbb{I}l-0) \longrightarrow 1)$ *at-top*
using *l-0-pos l-PInfty-0* **by** (*intro tendsto-eq-intros*) *simp-all+*

lemma *compl-l-real-distribution*: *real-distribution* (*interval-measure* $(\lambda x. 1 - \mathbb{I}l-x / \mathbb{I}l-0)$)
using *l-normal-antimono compl-l-normal-right-continuous compl-l-normal-NInfty-0 compl-l-normal-PInfty-1*
by (*intro real-distribution-interval-measure*; *simp add: antimono-def*)

definition *total* $:: \text{real} \Rightarrow \text{real}$ ($\mathbb{I}T'-$ [101] 200) **where** $\mathbb{I}T-x \equiv \text{LBINT } y:\{x..\}.$
 $\mathbb{I}l-y$
— the number of lives older than the ones aged x
— The parameter x must be nonnegative.

lemma *T-nonneg*[*simp*]: $\mathbb{I}T-x \geq 0$ **for** $x :: \text{real}$
unfolding *total-def* **by** *simp*

definition *total-finite* $\equiv \text{set-integrable lborel } \{0..\} l$

lemma *total-finite-iff-set-integrable-Ici*:
total-finite $\longleftrightarrow \text{set-integrable lborel } \{x..\} l$ **for** $x :: \text{real}$
unfolding *total-finite-def* **using** *set-integrable-Ici-equiv l-integrable-Icc* **by** *blast*

lemma *total-finite-iff-integrable-on-Ici*: *total-finite* $\longleftrightarrow l$ *integrable-on* $\{x..\}$ **for** $x :: \text{real}$
using *total-finite-iff-set-integrable-Ici integrable-on-iff-set-integrable-nonneg l-nonneg*
by (*metis atLeast-borel l-measurable measurable-lborel2 sets-lborel*)

lemma *total-finite-iff-summable*: *total-finite* $\longleftrightarrow \text{summable}$ $(\lambda k. \mathbb{I}l-(x+k))$ **for** $x :: \text{real}$
apply (*rewrite total-finite-iff-set-integrable-Ici*)
apply (*rule set-integrable-iff-summable*[*of x, simplified*], *simp-all*)
using *l-antimono* **unfolding** *antimono-def monotone-on-def* **by** *simp*

lemma $T\text{-tendsto-}0$: $((\lambda x. \$T\text{-}x) \longrightarrow 0)$ *at-top* **if** *total-finite*
proof –
have $\bigwedge x. x \geq 0 \implies \$T\text{-}x = \$T\text{-}0 - (LBINT\ y:\{0..x\}. \$l\text{-}y)$
proof –
fix $x::real$ **assume** $asm: x \geq 0$
let $?A = \{x..\}$ **and** $?B = \{0..x\}$
have $\{0..\} = ?A \cup ?B$ **using** asm **by** $auto$
thus $\$T\text{-}x = \$T\text{-}0 - (LBINT\ y:\{0..x\}. \$l\text{-}y)$
unfolding $total\text{-}def$ **apply** $(rewrite\ eq\text{-}diff\text{-}eq)$
using $that\ total\text{-}finite\text{-}iff\text{-}set\text{-}integrable\text{-}Ici\ l\text{-}integrable\text{-}Icc$
apply $(rewrite\ set\text{-}integral\text{-}Un\text{-}AE[THEN\ sym],\ simp\text{-}all)$
using $AE\text{-}lborel\text{-}singleton\ add\text{-}0\ asm\ le\text{-}add\text{-}same\text{-}cancel2\ le\text{-}numeral\text{-}extra(3)$
by $force$
qed
hence $\forall_F\ x\ in\ at\text{-}top. \$T\text{-}x = \$T\text{-}0 - (LBINT\ y:\{0..x\}. \$l\text{-}y)$
by $(rule\ eventually\text{-}at\text{-}top\text{-}linorderI[of\ 0])$
moreover **have** $((\lambda x. LBINT\ y:\{0..x\}. \$l\text{-}y) \longrightarrow \$T\text{-}0)$ *at-top*
using $that\ unfolding\ total\text{-}def\ total\text{-}finite\text{-}def$
by $(intro\ tendsto\text{-}set\text{-}lebesgue\text{-}integral\text{-}at\text{-}top;\ simp)$
ultimately **show** $?thesis$
apply $(rewrite\ tendsto\text{-}cong,\ simp\text{-}all)$
using $LIM\text{-}zero\text{-}iff'$ **by** $force$
qed

definition $lives :: real \Rightarrow real \Rightarrow real$ $(\$L'\text{-}\{\&x\}\ [0,0]\ 200)$
where $\$L'\text{-}\{n\&x\} \equiv LBINT\ y:\{x..x+n\}. \$l\text{-}y$
– the number of lives between ages x and $x+n$
– The parameter x must be nonnegative.
– The parameter n is usually nonnegative, but theoretically it can be negative.

abbreviation $lives\text{-}1 :: real \Rightarrow real$ $(\$L'\text{-}\ [101]\ 200)$
where $\$L\text{-}x \equiv \$L'\text{-}\{1\&x\}$

lemma $l\text{-}has\text{-}integral\text{-}L$: $(l\ has\text{-}integral\ \$L'\text{-}\{n\&x\}\ \{x..x+n\})$ **for** $x\ n :: real$
unfolding $lives\text{-}def$ **by** $(rule\ has\text{-}integral\text{-}set\text{-}integral\text{-}real)\ (rule\ l\text{-}integrable\text{-}Icc)$

lemma $L\text{-}neg\text{-}0[simp]$: $\$L'\text{-}\{n\&x\} = 0$ **if** $n < 0$ **for** $x\ n :: real$
unfolding $lives\text{-}def$ **using** $that$ **by** $(rewrite\ to\ \{\}\ atLeastatMost\text{-}empty;\ simp)$

lemma $L\text{-}nonneg[simp]$: $\$L'\text{-}\{n\&x\} \geq 0$ **for** $x\ n :: real$
unfolding $lives\text{-}def$ **by** $simp$

lemma $L\text{-}T$: $\$L'\text{-}\{n\&x\} = \$T\text{-}x - \$T\text{-}(x+n)$ **if** *total-finite* $n \geq 0$ **for** $x\ n :: real$
proof –
have $\{x..x+n\} \cup \{x+n..\} = \{x..\}$ **using** $that$ **by** $force$
moreover **have**
 $(LBINT\ y:\{x..x+n\} \cup \{x+n..\}. \$l\text{-}y) = (LBINT\ y:\{x..x+n\}. \$l\text{-}y) + (LBINT\ y:\{x+n..\}. \$l\text{-}y)$

proof –
have $AE\ y\ in\ lborel.\ \neg\ (y \in \{x..x+n\} \wedge y \in \{x+n.. \})$ **by** (rule $AE-I'$ [**where** $N=\{x+n\}$]; *force*)
moreover have *set-integrable lborel* $\{x..x+n\}$ l **by** (rule *l-integrable-Icc*)
moreover have *set-integrable lborel* $\{x+n.. \}$ l
using *that total-finite-iff-set-integrable-Ici* **by** *simp*
ultimately show *?thesis* **by** (*intro set-integral-Un-AE; simp*)
qed
ultimately show *?thesis* **unfolding** *total-def lives-def* **by** *simp*
qed

lemma $L\text{-sums-}T$: $(\lambda k.\ \$L\text{-}(x+k))\ sums\ \$T\text{-}x$ **if** *total-finite* **for** $x::real$

proof –
have $(\lambda k::nat.\ \$T\text{-}(x+k)) \longrightarrow 0$
using *T-tendsto-0*
apply (rule *filterlim-compose*[**where** $f=\lambda k::nat.\ x+k$ **and** $g=total$], *simp add:*
that)
using *filterlim-real-sequentially filterlim-tendsto-add-at-top* **by** *blast*
hence $(\lambda k.\ \$T\text{-}(x+k) - \$T\text{-}(x + Suc\ k))\ sums\ \$T\text{-}x$
by (*simp*) (rule *telescope-sums'*[*of* $\lambda k.\ \$T\text{-}(x+k)$ 0 , *simplified*])
thus *?thesis* **using** *that L-T* **by** (*rewrite sums-cong, simp-all*) *smt*
qed

definition $death :: real \Rightarrow real \Rightarrow real$ ($\$d'\text{-}\{\&x\}$ $[0,0]$ 200)

where $\$d\text{-}\{t\&x\} \equiv max\ 0\ (\$l\text{-}x - \$l\text{-}(x+t))$
– the number of deaths between ages x and $x+t$
– The parameter t is usually nonnegative, but theoretically it can be negative.

abbreviation $death1 :: real \Rightarrow real$ ($\$d'\text{-}$ $[101]$ 200)

where $\$d\text{-}x \equiv \$d\text{-}\{1\&x\}$

lemma $death\text{-def-nonneg}$: $\$d\text{-}\{t\&x\} = \$l\text{-}x - \$l\text{-}(x+t)$ **if** $t \geq 0$ **for** $t\ x :: real$

using *that l-antimono* **unfolding** $death\text{-def}$ *antimono-def* **by** *simp*

lemma $d\text{-nonpos-0}$: $\$d\text{-}\{t\&x\} = 0$ **if** $t \leq 0$ **for** $t\ x :: real$

using *that l-antimono* **unfolding** $death\text{-def}$ *antimono-def* **by** *simp*

corollary $d\text{-0-0}$: $\$d\text{-}\{0\&x\} = 0$ **for** $x::real$

using $d\text{-nonpos-0}$ **by** *simp*

lemma $d\text{-nonneg}[simp]$: $\$d\text{-}\{t\&x\} \geq 0$ **for** $t\ x :: real$

unfolding $death\text{-def}$ **by** *simp*

lemma $dx\text{-l}$: $\$d\text{-}x = \$l\text{-}x - \$l\text{-}(x+1)$ **for** $x::real$

using $death\text{-def-nonneg}$ **by** *simp*

lemma $sum\text{-}dx\text{-l}$: $(\sum k < n.\ \$d\text{-}(x+k)) = \$l\text{-}x - \$l\text{-}(x+n)$ **for** $x::real$ **and** $n::nat$

proof (*induction n*)

case 0

```

thus ?case by simp
next
  case (Suc n)
  thus ?case
    using dx-l
    by (metis Set-Interval.comm-monoid-add-class.sum.lessThan-Suc
      add-diff-cancel-left' diff-diff-eq2 of-nat-Suc)
qed

```

```

corollary d-sums-l: ( $\lambda k. \$d(x+k)$ ) sums  $\$l-x$  for  $x::real$ 
  unfolding sums-def
  apply (rewrite sum-dx-l)
  apply (rule tendsto-diff[where  $b=0$ , simplified], simp)
  using l-PIfty-0 filterlim-compose filterlim-real-sequentially filterlim-tendsto-add-at-top
    tendsto-const by blast

```

```

lemma add-d:  $\$d\{t\&x\} + \$d\{t' \& x+t\} = \$d\{t+t' \& x\}$  if  $t \geq 0$   $t' \geq 0$  for  $t$ 
 $t' :: real$ 
  using death-def-nonneg that by (smt (verit))

```

```

definition die-central ::  $real \Rightarrow real \Rightarrow real$  ( $\$m'\{-\&- \} [0,0] 200$ )
  where  $\$m\{n\&x\} \equiv \$d\{n\&x\} / \$L\{n\&x\}$ 
  — central death rate

```

```

abbreviation die-central-1 ::  $real \Rightarrow real$  ( $\$m'\{- [101] 200$ )
  where  $\$m-x \equiv \$m\{1\&x\}$ 

```

6.2 Construction of Survival Model from Life Table

```

definition life-table-measure ::  $real$  measure  $\mathfrak{M}$ 
  where  $\mathfrak{M} \equiv interval-measure (\lambda x. 1 - \$l-x / \$l-0)$ 

```

```

lemma prob-space-actuary-MM:  $prob-space-actuary$   $\mathfrak{M}$ 
  unfolding life-table-measure-def using compl-l-real-distribution real-distribution-def
  by (intro prob-space-actuary.intro) force

```

```

definition survival-model-X ::  $real \Rightarrow real$  ( $X$ ) where  $X \equiv \lambda x. x$ 

```

```

lemma survival-model-MM-X:  $survival-model$   $\mathfrak{M}$   $X$ 

```

```

proof —

```

```

  let ?F =  $\lambda x. 1 - \$l-x / \$l-0$ 

```

```

  show survival-model  $\mathfrak{M}$   $X$ 

```

```

    unfolding life-table-measure-def survival-model-X-def

```

```

  proof (rule survival-model.intro)

```

```

    show prob-space-actuary (interval-measure ?F)

```

```

      using prob-space-actuary-MM unfolding life-table-measure-def by simp

```

```

  show survival-model-axioms (interval-measure ?F) ( $\lambda x. x$ )

```

```

  proof —

```

```

    have [simp]:  $\{\xi::real. \xi \leq 0\} = \{..0\}$  by blast

```

```

have measure (interval-measure ( $\lambda x. 1 - \$l-x / \$l-0$ ))  $\{..0\} = 0$ 
using l-normal-antimono compl-l-normal-right-continuous compl-l-normal-NInfty-0
  by (rewrite measure-interval-measure-Iic, simp-all add: antimono-def)
hence emeasure (interval-measure ( $\lambda x. 1 - \$l-x / \$l-0$ ))  $\{..0\} = \text{ennreal } 0$ 
  apply (rewrite finite-measure.emeasure-eq-measure, simp-all)
  using compl-l-real-distribution prob-space-def real-distribution-def by blast
thus ?thesis
  apply (intro survival-model-axioms.intro, simp)
  apply (rewrite AE-iff-null, simp)
  by (rewrite not-less) auto
qed
qed
qed

end

```

```

sublocale life-table  $\subseteq$  survival-model  $\mathfrak{M} X$ 
  by (rule survival-model-MM-X)

```

```

context life-table
begin

```

```

interpretation distrX-RD: real-distribution distr  $\mathfrak{M}$  borel X
  using MM-PS.real-distribution-distr by simp

```

6.2.1 Relations between Life Table and Survival Function for X

```

lemma ccdfX-l-normal: ccdf (distr  $\mathfrak{M}$  borel X) = ( $\lambda x. \$l-x / \$l-0$ )

```

```

proof (rule ext)

```

```

  let ?F =  $\lambda x. 1 - \$l-x / \$l-0$ 

```

```

  interpret F-FBM: finite-borel-measure interval-measure ?F

```

```

    using compl-l-real-distribution real-distribution.finite-borel-measure-M by blast

```

```

  show  $\bigwedge x. \text{ccdf} (\text{distr } \mathfrak{M} \text{ borel } X) x = \$l-x / \$l-0$ 

```

```

    unfolding ccdf-def life-table-measure-def survival-model-X-def

```

```

    apply (rewrite measure-distr, simp-all)

```

```

    using l-normal-antimono compl-l-normal-right-continuous

```

```

      compl-l-normal-NInfty-0 compl-l-normal-PInfty-1

```

```

    by (rewrite F-FBM.measure-interval-measure-Ioi; simp add: antimono-def)

```

```

qed

```

```

corollary deriv-ccdfX-l: deriv (ccdf (distr  $\mathfrak{M}$  borel X))  $x = \text{deriv } l x / \$l-0$ 

```

```

  if l differentiable at  $x$  for  $x::\text{real}$ 

```

```

  using differentiable-eq-field-differentiable-real that

```

```

  by (rewrite ccdfX-l-normal, rewrite deriv-cdivide-right; simp)

```

```

notation death-pt ( $\psi$ )

```

```

lemma l-0-equiv:  $\$l-x = 0 \iff x \geq \psi$  for  $x::\text{real}$ 

```

```

  using ccdfX-l-normal ccdfX-0-equiv by simp

```


lemma *d-old-0*: $\mathbb{P}\{t \leq x\} = 0$ if $x \geq \psi$ $t \geq 0$ for $x t :: \text{real}$
unfolding *death-def* **using** *l-0-equiv* that **by** (*smt* (*verit*) *le-ereal-le*)

lemma *d-l-equiv*: $\mathbb{P}\{t \leq x\} = \mathbb{P}\{x+t \leq \psi\}$ if $t \geq 0$ for $x t :: \text{real}$
using *death-def-nonneg* *l-0-equiv* that **by** *simp*

lemma *continuous-ccdfX-l*: continuous F ($\text{ccdf} (\text{distr } \mathfrak{M} \text{ borel } X)$) \longleftrightarrow continuous F l
for $F :: \text{real filter}$
proof –
have continuous F ($\text{ccdf} (\text{distr } \mathfrak{M} \text{ borel } X)$) \longleftrightarrow continuous F ($\lambda x. \mathbb{P}\{x \leq l\}$)
using *ccdfX-l-normal* **by** *simp*
also have ... \longleftrightarrow continuous F l **using** *continuous-cdivide-iff* *l-0-neq-0* **by** *blast*
finally show *?thesis* .
qed

lemma *has-real-derivative-ccdfX-l*:
($\text{ccdf} (\text{distr } \mathfrak{M} \text{ borel } X)$ has-real-derivative D) (at x) \longleftrightarrow
(l has-real-derivative $\mathbb{P}\{x \leq l\}$) (at x)
for $D x :: \text{real}$
proof –
have ($\text{ccdf} (\text{distr } \mathfrak{M} \text{ borel } X)$ has-real-derivative D) (at x) \longleftrightarrow
($\lambda x. \mathbb{P}\{x \leq l\}$ has-real-derivative D) (at x)
by (*rule* *has-field-derivative-cong-eventually*; *simp* *add*: *ccdfX-l-normal*)
also have ... \longleftrightarrow ($\lambda x. \mathbb{P}\{x \leq l\}$ has-real-derivative $\mathbb{P}\{x \leq l\}$) (at x)
by (*rule* *DERIV-cmult-iff*, *simp*)
also have ... \longleftrightarrow (l has-real-derivative $\mathbb{P}\{x \leq l\}$) (at x) **by** *simp*
finally show *?thesis* .
qed

corollary *differentiable-ccdfX-l*:
 $\text{ccdf} (\text{distr } \mathfrak{M} \text{ borel } X)$ differentiable (at x) \longleftrightarrow l differentiable (at x)
for $D x :: \text{real}$
using *has-real-derivative-ccdfX-l*
by (*metis* *l-0-neq-0* *mult.commute* *nonzero-divide-eq-eq* *real-differentiable-def*)

lemma *PX-l-normal*: $\mathbb{P}(\xi \text{ in } \mathfrak{M}. X \xi > x) = \mathbb{P}\{x \leq l\}$ for $x :: \text{real}$
using *MM-PS.ccdf-distr-P* *ccdfX-l-normal* *X-RV* **by** (*metis* (*mono-tags*, *lifting*)
Collect-cong)

lemma *set-integrable-ccdfX-l*:
set-integrable l borel A ($\text{ccdf} (\text{distr } \mathfrak{M} \text{ borel } X)$) \longleftrightarrow set-integrable l borel A l
if $A \in \text{sets } l$ borel **for** $A :: \text{real set}$
proof –
have set-integrable l borel A ($\text{ccdf} (\text{distr } \mathfrak{M} \text{ borel } X)$) \longleftrightarrow
set-integrable l borel A ($\lambda x. \mathbb{P}\{x \leq l\}$)
by (*rule* *set-integrable-cong*; *simp* *add*: *ccdfX-l-normal*)
also have ... \longleftrightarrow set-integrable l borel A l **by** *simp*

finally show ?thesis .
qed

corollary *integrable-ccdfX-l*: $\text{integrable lborel } (\text{ccdf } (\text{distr } \mathfrak{M} \text{ borel } X)) \longleftrightarrow \text{integrable lborel } l$
using *set-integrable-ccdfX-l*[where $A=UNIV$] by (*simp add: set-integrable-def*)

lemma *integrable-on-ccdfX-l*:

$\text{ccdf } (\text{distr } \mathfrak{M} \text{ borel } X) \text{ integrable-on } A \longleftrightarrow l \text{ integrable-on } A$ for $A :: \text{real set}$

proof –

have $\text{ccdf } (\text{distr } \mathfrak{M} \text{ borel } X) \text{ integrable-on } A \longleftrightarrow (\lambda x. \$l-x / \$l-0) \text{ integrable-on } A$

by (*rule integrable-cong*) (*simp add: ccdfX-l-normal*)

also have $\dots \longleftrightarrow l \text{ integrable-on } A$

using *integrable-on-cdivide-iff*[of $\$l-0$ l] by *simp*

finally show ?thesis .

qed

6.2.2 Relations between Life Table and Cumulative Distributive Function for X

lemma *cdfX-l-normal*: $\text{cdf } (\text{distr } \mathfrak{M} \text{ borel } X) = (\lambda x. 1 - \$l-x / \$l-0)$ for $x::\text{real}$
using *ccdfX-l-normal distrX-RD.cdf-ccdf distrX-RD.prob-space* by *presburger*

lemma *deriv-cdfX-l*: $\text{deriv } (\text{cdf } (\text{distr } \mathfrak{M} \text{ borel } X)) x = - \text{deriv } l x / \$l-0$

if l differentiable at x for $x::\text{real}$

using *distrX-RD.cdf-ccdf differentiable-eq-field-differentiable-real that differentiable-ccdfX-l*

deriv-diff deriv-ccdfX-l that by *simp*

lemma *continuous-cdfX-l*: $\text{continuous } F (\text{cdf } (\text{distr } \mathfrak{M} \text{ borel } X)) \longleftrightarrow \text{continuous } F l$

for $F :: \text{real filter}$

using *distrX-RD.continuous-cdf-ccdf continuous-ccdfX-l* by *simp*

lemma *has-real-derivative-cdfX-l*:

$(\text{cdf } (\text{distr } \mathfrak{M} \text{ borel } X) \text{ has-real-derivative } D) (\text{at } x) \longleftrightarrow$

$(l \text{ has-real-derivative } - (\$l-0 * D)) (\text{at } x)$

for $D x :: \text{real}$

using *distrX-RD.has-real-derivative-cdf-ccdf has-real-derivative-ccdfX-l* by *simp*

lemma *differentiable-cdfX-l*:

$\text{cdf } (\text{distr } \mathfrak{M} \text{ borel } X) \text{ differentiable } (\text{at } x) \longleftrightarrow l \text{ differentiable } (\text{at } x)$ for $D x :: \text{real}$

using *differentiable-eq-field-differentiable-real distrX-RD.differentiable-cdf-ccdf differentiable-ccdfX-l* by *simp*

lemma *PX-compl-l-normal*: $\mathcal{P}(\xi \text{ in } \mathfrak{M}. X \xi \leq x) = 1 - \$l-x / \$l-0$ for $x::\text{real}$

using *PX-l-normal* by (*metis MM-PS.prob-compl X-compl-gt-le X-gt-event*)

6.2.3 Relations between Life Table and Survival Function for $T(x)$

context

fixes $x::real$
assumes $x\text{-lt-psi}[simp]: x < \psi$

begin

notation $futr\text{-life} (T)$

interpretation $alive\text{-}PS$: *prob-space* $\mathfrak{M} \downarrow$ *alive* x

by (rule $MM\text{-}PS.cond\text{-}prob\text{-}space\text{-}correct$, $simp\text{-}all$ add: $alive\text{-}def$)

interpretation $distrTx\text{-}RD$: *real-distribution* $distr (\mathfrak{M} \downarrow$ *alive* $x)$ *borel* $(T x)$ **by** $simp$

lemma $lx\text{-}neq\text{-}0[simp]: \$l\text{-}x \neq 0$

using $l\text{-}0\text{-}equiv\ x\text{-}lt\text{-}psi\ linorder\text{-}not\text{-}less$ **by** $blast$

corollary $lx\text{-}pos[simp]: \$l\text{-}x > 0$

using $lx\text{-}neq\text{-}0\ l\text{-}nonneg$ **by** ($smt (verit)$)

lemma $ccdfTx\text{-}l\text{-}normal$: $ccdf (distr (\mathfrak{M} \downarrow$ *alive* $x)$ *borel* $(T x))\ t = \$l\text{-}(x+t) / \$l\text{-}x$
if $t \geq 0$ **for** $t::real$

using $ccdfTx\text{-}PX\ PX\text{-}l\text{-}normal\ l\text{-}0\text{-}neq\text{-}0$ **that** **by** $simp$

lemma $deriv\text{-}ccdfTx\text{-}l$:

$deriv (ccdf (distr (\mathfrak{M} \downarrow$ *alive* $x)$ *borel* $(T x)))\ t = deriv (\lambda t. \$l\text{-}(x+t) / \$l\text{-}x)\ t$
if $t > 0$ **l** *differentiable at* $(x+t)$ **for** $t::real$

proof –

have $\forall_F\ s$ *in nhds* t . $ccdf (distr (\mathfrak{M} \downarrow$ *alive* $x)$ *borel* $(T x))\ s = \$l\text{-}(x+s) / \$l\text{-}x$

apply ($rewrite\ eventually\text{-}nhds\text{-}metric$)

using $that\ ccdfTx\text{-}l\text{-}normal\ dist\text{-}real\text{-}def$ **by** ($intro\ exI[of\ -\ t]$) $auto$

thus $?thesis$ **by** ($rule\ deriv\text{-}cong\text{-}ev$) $simp$

qed

lemma $continuous\text{-}at\text{-}within\text{-}ccdfTx\text{-}l$:

$continuous (at\ t\ within\ \{0..\}) (ccdf (distr (\mathfrak{M} \downarrow$ *alive* $x)$ *borel* $(T x))) \longleftrightarrow$

$continuous (at\ (x+t)\ within\ \{x..\})\ l$

if $t \geq 0$ **for** $t::real$

using $continuous\text{-}ccdfX\text{-}ccdfTx$ **that** $continuous\text{-}ccdfX\text{-}l$ **by** $force$

lemma $isCont\text{-}ccdfTx\text{-}l$:

$isCont (ccdf (distr (\mathfrak{M} \downarrow$ *alive* $x)$ *borel* $(T x)))\ t \longleftrightarrow isCont\ l\ (x+t)$ **if** $t > 0$ **for** $t::real$

using $that\ continuous\text{-}ccdfX\text{-}l\ isCont\text{-}ccdfX\text{-}ccdfTx$ **by** $force$

lemma $has\text{-}real\text{-}derivative\text{-}ccdfTx\text{-}l$:

$(ccdf (distr (\mathfrak{M} \downarrow$ *alive* $x)$ *borel* $(T x)))\ has\text{-}real\text{-}derivative\ D (at\ t) \longleftrightarrow$

$(l\ has\text{-}real\text{-}derivative\ \$l\text{-}x * D) (at\ (x+t))$

if $t > 0$ **for** $t D :: \text{real}$
proof –
have $(\text{ccdf } (\text{distr } \mathfrak{M} \mid \text{alive } x) \text{ borel } (T x)) \text{ has-real-derivative } D) (at t) \longleftrightarrow$
 $(\text{ccdf } (\text{distr } \mathfrak{M} \mid \text{alive } x) \text{ borel } (T x)) \text{ has-real-derivative}$
 $(\$l-x / \$l-0 * D / \mathcal{P}(\xi \text{ in } \mathfrak{M}. X \xi > x)) (at t)$
using $PX-l\text{-normal}$ **by force**
also have $\dots = (\text{ccdf } (\text{distr } \mathfrak{M} \text{ borel } X) \text{ has-real-derivative } (\$l-x / \$l-0 * D))$
 $(at (x+t))$
using $\text{has-real-derivative-ccdfX-ccdfTx}$ **that by simp**
also have $\dots = (l \text{ has-real-derivative } (\$l-x * D)) (at (x+t))$
using $\text{has-real-derivative-ccdfX-l}$ **by simp**
finally show $?thesis$.
qed

lemma $\text{differentiable-ccdfTx-l}$:
 $\text{ccdf } (\text{distr } \mathfrak{M} \mid \text{alive } x) \text{ borel } (T x) \text{ differentiable at } t \longleftrightarrow l \text{ differentiable at}$
 $(x+t)$
if $t > 0$ **for** $t::\text{real}$
using $\text{differentiable-ccdfX-ccdfTx}$ $\text{differentiable-ccdfX-l}$ **that by force**

lemma $PTx-l\text{-normal}$: $\mathcal{P}(\xi \text{ in } \mathfrak{M}. T x \xi > t \mid T x \xi > 0) = \$l-(x+t) / \$l-x$ **if** $t \geq$
 0 **for** $t::\text{real}$
using ccdfTx-l-normal **that by** $(\text{simp add: ccdfTx-cond-prob})$

6.2.4 Relations between Life Table and Cumulative Distributive Function for $T(x)$

lemma $\text{cdfTx-compl-l-normal}$: $\text{cdf } (\text{distr } \mathfrak{M} \mid \text{alive } x) \text{ borel } (T x) t = 1 -$
 $\$l-(x+t) / \$l-x$
if $t \geq 0$ **for** $t::\text{real}$
using $\text{distrTx-RD.cdf-ccdf}$ cdfTx-l-normal **that** $\text{distrTx-RD.prob-space}$ **by auto**

lemma deriv-cdfTx-l :
 $\text{deriv } (\text{cdf } (\text{distr } \mathfrak{M} \mid \text{alive } x) \text{ borel } (T x)) t = - \text{deriv } (\lambda t. \$l-(x+t) / \$l-x) t$
if $t > 0$ l **differentiable at** $(x+t)$ **for** $t::\text{real}$
using deriv-ccdfTx-l $\text{differentiable-cdfX-cdfTx}$ $\text{differentiable-cdfX-l}$ $\text{distrTx-RD.deriv-cdf-ccdf}$
that by fastforce

lemma $\text{continuous-at-within-cdfTx-l}$:
 $\text{continuous } (at t \text{ within } \{0..\}) (\text{cdf } (\text{distr } \mathfrak{M} \mid \text{alive } x) \text{ borel } (T x)) \longleftrightarrow$
 $\text{continuous } (at (x+t) \text{ within } \{x..\}) l$
if $t \geq 0$ **for** $t::\text{real}$
using $\text{that continuous-cdfX-l}$ $\text{continuous-cdfX-cdfTx}$ **by force**

lemma isCont-cdfTx-l :
 $\text{isCont } (\text{cdf } (\text{distr } \mathfrak{M} \mid \text{alive } x) \text{ borel } (T x)) t \longleftrightarrow \text{isCont } l (x+t)$ **if** $t > 0$ **for**
 $t::\text{real}$
using $\text{that continuous-cdfX-l}$ isCont-cdfX-cdfTx **by force**

lemma *has-real-derivative-cdfTx-l*:
 $(cdf (distr (\mathfrak{M} \downarrow \text{alive } x) \text{ borel } (T x)) \text{ has-real-derivative } D) (at t) \longleftrightarrow$
 $(l \text{ has-real-derivative } - \$l-x * D) (at (x+t))$
if $t > 0$ **for** $t D :: \text{real}$
using *has-real-derivative-ccdfTx-l* that *distrTx-RD.has-real-derivative-cdf-ccdf* **by**
auto

lemma *differentiable-cdfTx-l*:
 $cdf (distr (\mathfrak{M} \downarrow \text{alive } x) \text{ borel } (T x)) \text{ differentiable at } t \longleftrightarrow l \text{ differentiable (at}$
 $(x+t))$
if $t > 0$ **for** $t :: \text{real}$
using *differentiable-cdfX-l* that *differentiable-cdfX-cdfTx* **by** *auto*

lemma *PTx-compl-l-normal*: $\mathcal{P}(\xi \text{ in } \mathfrak{M}. T x \xi \leq t \mid T x \xi > 0) = 1 - \$l-(x+t) /$
 $\$l-x$
if $t \geq 0$ **for** $t :: \text{real}$
using *cdfTx-compl-l-normal* that **by** (*simp add: cdfTx-cond-prob*)

6.2.5 Life Table and Actuarial Notations

notation *survive* ($\$p'\{-\&- \} [0,0] 200$)
notation *survive-1* ($\$p'--[101] 200$)
notation *die* ($\$q'\{-\&- \} [0,0] 200$)
notation *die-1* ($\$q'--[101] 200$)
notation *die-defer* ($\$q'\{-|\&- \} [0,0,0] 200$)
notation *die-defer-1* ($\$q'\{-|\&- \} [0,0] 200$)
notation *life-expect* ($\$e'\circ'--[101] 200$)
notation *temp-life-expect* ($\$e'\circ'\{-:- \} [0,0] 200$)
notation *curt-life-expect* ($\$e'--[101] 200$)
notation *temp-curt-life-expect* ($\$e'\{-:- \} [0,0] 200$)

lemma *p-l*: $\$p\{-t\&x\} = \$l-(x+t) / \$l-x$ **if** $t \geq 0$ **for** $t :: \text{real}$
unfolding *survive-def* **using** *ccdfTx-l-normal* that **by** *simp*

corollary *p-1-l*: $\$p-x = \$l-(x+1) / \$l-x$
using *p-l* **by** *simp*

lemma *isCont-p-l*: $isCont (\lambda s. \$p\{-s\&x\}) t \longleftrightarrow isCont l (x+t)$ **if** $t > 0$ **for** $t :: \text{real}$
proof –

have $\forall_F s \text{ in nhds } t. \$p\{-s\&x\} = \$l-(x+s) / \$l-x$
apply (*rewrite eventually-nhds-metric*)
apply (*rule exI[of - t], auto simp add: that*)
by (*rewrite p-l; simp add: dist-real-def*)
hence $isCont (\lambda s. \$p\{-s\&x\}) t \longleftrightarrow isCont (\lambda s. \$l-(x+s) / \$l-x) t$ **by** (*rule isCont-cong*)
also have $\dots \longleftrightarrow isCont (\lambda s. \$l-(x+s)) t$ **using** *continuous-cdivide-iff lx-neq-0*
by *metis*
also have $\dots \longleftrightarrow isCont l (x+t)$ **using** *isCont-shift* **by** (*force simp add: add commute*)

finally show *?thesis* .
qed

lemma *total-finite-iff-p-set-integrable-Ici*:
 $total-finite \iff set-integrable\ lborel\ \{0..\} (\lambda t. \$p-\{t\&x\})$
apply (*rewrite set-integrable-cong-AE*[**where** $g=\lambda t. \$l-(x+t) / \$l-x$]; *simp*)
using *survive-def* **apply** *simp*
using *p-l* **apply** (*intro AE-I2*, *simp*)
by (*metis l-integrable-Icc-shift set-integrable-Ici-equiv set-integrable-Ici-shift total-finite-iff-set-integrable-Ici*)

lemma *p-PTx-ge-l-isCont*: $\$p-\{t\&x\} = \mathcal{P}(\xi\ in\ \mathfrak{M}. T\ x\ \xi \geq t \mid T\ x\ \xi > 0)$
if *isCont* $l\ (x+t)\ t > 0$ **for** $t::real$
using *p-PTx-ge-ccdf-isCont* *that continuous-ccdfX-l* **by** *force*

lemma *q-defer-l*: $\$q-\{f|t\&x\} = (\$l-(x+f) - \$l-(x+f+t)) / \$l-x$ **if** $f \geq 0\ t \geq 0$ **for** $f\ t :: real$
apply (*rewrite q-defer-p*, *simp-all add: that*)
using *that* **by** (*rewrite p-l*, *simp*) + (*smt (verit) diff-divide-distrib*)

corollary *q-defer-d-l*: $\$q-\{f|t\&x\} = \$d-\{t\ \&\ x+f\} / \$l-x$ **if** $f \geq 0\ t \geq 0$ **for** $f\ t :: real$
using *q-defer-l* *that death-def-nonneg* **by** *simp*

corollary *q-defer-1-d-l*: $\$q-\{f|\&x\} = \$d-(x+f) / \$l-x$ **if** $f \geq 0$ **for** $f::real$
using *q-defer-d-l* *that* **by** *simp*

lemma *q-d-l*: $\$q-\{t\&x\} = \$d-\{t\&x\} / \$l-x$ **for** $t::real$
proof (*cases* $\langle t \geq 0 \rangle$)
case *True*
thus *?thesis* **using** *q-defer-d-l[of 0]* **by** *simp*
next
case *False*
thus *?thesis* **using** *q-nonpos-0 d-nonpos-0* **by** *simp*
qed

corollary *q-1-d-l*: $\$q-x = \$d-x / \$l-x$
using *q-d-l* **by** *simp*

lemma *LBINT-p-l*: $(LBINT\ t:A. \$p-\{t\&x\}) = (LBINT\ t:A. \$l-(x+t)) / \$l-x$
if $A \subseteq \{0..\}$ $A \in sets\ lborel$ **for** $A :: real\ set$
— Note that $0 = 0$ holds when the integral diverges.
proof —
have [*simp*]: $\bigwedge t. t \in A \implies \$p-\{t\&x\} = \$l-(x+t) / \$l-x$ **using** *p-l* *that* **by** *blast*
hence $(LBINT\ t:A. \$p-\{t\&x\}) = (LBINT\ t:A. \$l-(x+t)) / \$l-x$
using *that* **by** (*rewrite set-lebesgue-integral-cong*[**where** $g=\lambda t. \$l-(x+t) / \$l-x$];
simp)
also have $\dots = (LBINT\ t:A. \$l-(x+t)) / \$l-x$ **by** (*rewrite set-integral-divide-zero*)
simp

finally show *?thesis* .
qed

corollary *e-LBINT-l*: $\$e^{\circ}x = (LBINT\ t:\{0..n\}.\ \$l(x+t)) / \$l-x$
— Note that $0 = 0$ holds when the integral diverges.
by (*simp add: e-LBINT-p LBINT-p-l*)

corollary *e-LBINT-l-Icc*: $\$e^{\circ}x = (LBINT\ t:\{0..n\}.\ \$l(x+t)) / \$l-x$ **if** $x+n \geq \$\psi$ **for** $n::real$
using *e-LBINT-p-Icc* **by** (*rewrite LBINT-p-l[THEN sym]; simp add: that*)

lemma *temp-e-LBINT-l*: $\$e^{\circ}\{x:n\} = (LBINT\ t:\{0..n\}.\ \$l(x+t)) / \$l-x$ **if** $n \geq 0$ **for** $n::real$
using *temp-e-LBINT-p* **by** (*rewrite LBINT-p-l[THEN sym]; simp add: that*)

lemma *integral-p-l*: $integral\ A\ (\lambda t.\ \$p\{t\&x\}) = (integral\ A\ (\lambda t.\ \$l(x+t))) / \$l-x$
if $A \subseteq \{0..n\}$ $A \in sets\ lborel$ **for** $A :: real\ set$
— Note that $0 = 0$ holds when the integral diverges.
using *that apply* (*rewrite set-borel-integral-eq-integral-nonneg[THEN sym], simp-all*)
apply (*simp add: survive-def*)
apply (*rewrite set-borel-integral-eq-integral-nonneg[THEN sym], simp-all*)
by (*rule LBINT-p-l; simp*)

corollary *e-integral-l*: $\$e^{\circ}x = integral\ \{0..n\}\ (\lambda t.\ \$l(x+t)) / \$l-x$
— Note that $0 = 0$ holds when the integral diverges.
by (*simp add: e-integral-p integral-p-l*)

corollary *e-integral-l-Icc*:
 $\$e^{\circ}x = integral\ \{0..n\}\ (\lambda t.\ \$l(x+t)) / \$l-x$ **if** $x+n \geq \$\psi$ **for** $n::real$
using *e-integral-p-Icc* **by** (*rewrite integral-p-l[THEN sym]; simp add: that*)

lemma *e-pos-total-finite*: $\$e^{\circ}x > 0$ **if** *total-finite*
using *e-pos total-finite-iff-p-set-integrable-Ici* **that** **by** *simp*

lemma *temp-e-integral-l*:
 $\$e^{\circ}\{x:n\} = integral\ \{0..n\}\ (\lambda t.\ \$l(x+t)) / \$l-x$ **if** $n \geq 0$ **for** $n::real$
using *temp-e-integral-p* **by** (*rewrite integral-p-l[THEN sym]; simp add: that*)

lemma *curt-e-sum-l*: $\$e-x = (\sum k.\ \$l(x+k+1)) / \$l-x$ **if** *total-finite* $\wedge k::nat.$ *isCont* $l(x+k+1)$

proof —

have *summable* $(\lambda k.\ \$l(x+(k+1::nat)))$
using *that total-finite-iff-summable* **by** (*rewrite summable-iff-shift[of $\lambda k.\ \$l(x+k)$ 1] simp*)
moreover **hence** *summable* $(\lambda k.\ \$p\{k+1\&x\})$ **by** (*rewrite p-l, simp-all add: add.commute*)
moreover **have** $\wedge k::nat.$ *isCont* $(\lambda t.\ \$p\{t\&x\}) (k+1)$
using *isCont-p-l* **that** **by** (*simp add: add.assoc*)
ultimately show *?thesis*

apply (*rewrite curt-e-sum-p, simp-all*)
apply (*rewrite p-l, simp*)
by (*rewrite suminf-divide*) (*simp add: add.commute, simp add: add.assoc*)
qed

lemma curt-e-sum-l-finite: $\$e\text{-}x = (\sum k < n. \$l\text{-}(x+k+1)) / \$l\text{-}x$
if $\bigwedge k :: \text{nat}. k < n \implies \text{isCont } l \text{ } (x+k+1) \text{ } x+n+1 > \ψ **for** $n :: \text{nat}$
apply (*rewrite curt-e-sum-p-finite[of x n], simp-all add: that*)
using *isCont-p-l that* **apply** (*simp add: add.assoc*)
apply (*rewrite sum-divide-distrib, rule sum.cong, simp*)
using *p-l* **by** (*smt (verit) of-nat-0-le-iff*)

lemma temp-curt-e-sum-p: $\$e\text{-}\{x:n\} = (\sum k < n. \$l\text{-}(x+k+1)) / \$l\text{-}x$
if $\bigwedge k :: \text{nat}. k < n \implies \text{isCont } l \text{ } (x+k+1)$ **for** $n :: \text{nat}$
apply (*rewrite temp-curt-e-sum-p[of x n], simp-all add: that*)
using *isCont-p-l that* **apply** (*simp add: add.assoc*)
apply (*rewrite sum-divide-distrib, rule sum.cong, simp*)
using *p-l* **by** (*smt (verit) of-nat-0-le-iff*)

lemma e-T-l: $\$e\text{'o-}x = \$T\text{-}x / \$l\text{-}x$
unfolding *total-def*
apply (*rewrite e-LBINT-l, simp-all*)
by (*metis add-cancel-left-left diff-add-cancel lborel-set-integral-Ici-shift*)

lemma temp-e-L-l: $\$e\text{'o-}\{x:n\} = \$L\text{-}\{n\&x\} / \$l\text{-}x$ **if** $n \geq 0$ **for** $n :: \text{real}$
unfolding *lives-def* **using** *that*
apply (*rewrite temp-e-LBINT-l, simp-all*)
using *diff-self add-diff-cancel-left' lborel-set-integral-Icc-shift* **by** *metis*

lemma m-q-e: $\$m\text{-}\{n\&x\} = \$q\text{-}\{n\&x\} / \$e\text{'o-}\{x:n\}$ **if** $n \geq 0$ **for** $n :: \text{real}$
proof –
have $\$m\text{-}\{n\&x\} = (\$d\text{-}\{n\&x\} / \$l\text{-}x) / (\$L\text{-}\{n\&x\} / \$l\text{-}x)$ **unfolding** *die-central-def*
by *simp*
thus *?thesis* **using** *q-d-l temp-e-L-l* **that** **by** *simp*
qed

end

lemma l-p: $\$l\text{-}x / \$l\text{-}0 = \$p\text{-}\{x\&0\}$ **for** $x :: \text{real}$
using *ccdfX-l-normal cdfX-p* **by** *force*

lemma e-p-e-total-finite: $\$e\text{'o-}x = \$e\text{'o-}\{x:n\} + \$p\text{-}\{n\&x\} * \$e\text{'o-}(x+n)$
if *total-finite* $n \geq 0 \text{ } x+n < \ψ **for** $x \text{ } n :: \text{real}$
using *e-p-e* **that** *total-finite-iff-p-set-integrable-Ici* **by** (*smt (verit) ereal-less-le*)

proposition x-ex-const-equiv-total-finite: $x + \$e\text{'o-}x = y + \$e\text{'o-}y \iff \$q\text{-}\{y-x\&x\} = 0$
if *total-finite* $x \leq y \text{ } y < \$\psi$ **for** $x \text{ } y :: \text{real}$
using *x-ex-const-equiv* **that** *total-finite-iff-p-set-integrable-Ici* *p-set-integrable-shift*

by blast

corollary *x-ex-const-iff-l-const*: $x + \mathbb{e}'_o x = y + \mathbb{e}'_o y \iff \mathbb{L}x = \mathbb{L}y$
 if total-finite $x \leq y$ $y < \mathbb{L}\psi$ for $x y :: \text{real}$
 using *x-ex-const-equiv-total-finite* that
 by (smt (verit, ccfv-threshold) divide-cancel-right ereal-less-le
 l-0-equiv life-table.death-def-nonneg life-table.q-d-l life-table-axioms q-1-equiv)

end

6.3 Piecewise Differentiable Life Table

locale *smooth-life-table* = *life-table* +
 assumes *l-piecewise-differentiable[simp]*: l piecewise-differentiable-on UNIV
begin

lemma *smooth-survival-function-MM-X*: *smooth-survival-function* $\mathfrak{M} X$

proof (rule *smooth-survival-function.intro*)

show *survival-model* $\mathfrak{M} X$ by (rule *survival-model-axioms*)

show *smooth-survival-function-axioms* $\mathfrak{M} X$

proof

show *ccdf* (*distr* \mathfrak{M} *borel* X) *piecewise-differentiable-on* UNIV

apply (rewrite *ccdfX-l-normal*)

apply (rewrite *divide-inverse*, rewrite *mult commute*)

using *l-piecewise-differentiable piecewise-differentiable-scaleR*[of l] by *simp*

qed

qed

end

sublocale *smooth-life-table* \subseteq *smooth-survival-function* $\mathfrak{M} X$

by (rule *smooth-survival-function-MM-X*)

context *smooth-life-table*

begin

notation *force-mortal* ($\mathbb{L}\mu'$ -- [101] 200)

lemma *l-continuous[simp]*: *continuous-on* UNIV l

using *l-piecewise-differentiable piecewise-differentiable-on-imp-continuous-on* by *fastforce*

lemma *l-nondifferentiable-finite-set[simp]*: *finite* $\{x. \neg l \text{ differentiable at } x\}$

using *differentiable-ccdfX-l cdfX-nondifferentiable-finite-set* by *simp*

lemma *l-differentiable-borel-set[measurable, simp]*: $\{x. l \text{ differentiable at } x\} \in \text{sets borel}$

using *differentiable-ccdfX-l cdfX-differentiable-borel-set* by *simp*

lemma *l-differentiable-AE*: AE x in $lborel$. l differentiable at x
using *differentiable-ccdfX-l cdfX-differentiable-AE* **by** *simp*

lemma *deriv-l-measurable[measurable]*: $deriv\ l \in borel\text{-measurable}\ borel$

proof –

let $?S = \{x. \neg l\ \text{differentiable at } x\}$

have $\bigwedge x. x \notin ?S \implies \mathbb{R} * deriv\ (ccdf\ (distr\ \mathfrak{M}\ borel\ X))\ x = deriv\ l\ x$

using *deriv-ccdfX-l* **by** *simp*

thus *?thesis*

apply –

by (*rule measurable-discrete-difference*

[**where** $X = ?S$ **and** $f = \lambda x. \mathbb{R} * deriv\ (ccdf\ (distr\ \mathfrak{M}\ borel\ X))\ x$])

(*simp-all add: countable-finite*)

qed

lemma *pdfX-l-normal*:

$pdfX\ x = (if\ l\ \text{differentiable at } x\ \text{then } -\ deriv\ l\ x / \mathbb{R} * l\ \text{else } 0)$ **for** $x :: real$

unfolding *pdfX-def*

using *differentiable-eq-field-differentiable-real differentiable-cdfX-l deriv-cdfX-l* **by** *simp*

lemma *mu-deriv-l*: $\mu\text{-}x = -\ deriv\ l\ x / \mathbb{R} * l$ **if** l differentiable at x **for** $x :: real$

using *mu-pdfX that cdfX-l-normal that pdfX-l-normal* **by** (*simp add: differentiable-cdfX-l*)

lemma *mu-nonneg-differentiable-l*: $\mu\text{-}x \geq 0$ **if** l differentiable at x **for** $x :: real$

using *differentiable-cdfX-l mu-nonneg-differentiable that* **by** *simp*

lemma *mu-deriv-ln-l*:

$\mu\text{-}x = -\ deriv\ (\lambda x. \ln\ (\mathbb{R} * l))\ x$ **if** l differentiable at x $x < \mathbb{R} * \psi$ **for** $x :: real$

proof –

have $\forall_F\ x\ \text{in nhds } x. \ln\ (\mathbb{R} * l / \mathbb{R} * l) = \ln\ (\mathbb{R} * l) - \ln\ (\mathbb{R} * l)$

proof (*cases* $\langle \mathbb{R} * \psi < \infty \rangle$)

case *True*

thus *?thesis*

apply (*rewrite eventually-nhds-metric*)

apply (*intro exI[of - real-of-ereal $\mathbb{R} * \psi - x$], auto*)

using *that True not-inftyI* **apply** *fastforce*

apply (*rewrite ln-div, simp-all*)

using *lx-pos dist-real-def not-inftyI that(2)* **by** *fastforce*

next

case *False*

hence $\bigwedge x. \mathbb{R} * l > 0$ **using** *l-0-equiv* **by** *force*

thus *?thesis* **by** (*intro always-eventually, rewrite ln-div; simp*)

qed

hence $deriv\ (\lambda x. \ln\ (\mathbb{R} * l / \mathbb{R} * l))\ x = deriv\ (\lambda x. \ln\ (\mathbb{R} * l))\ x$

apply (*rewrite deriv-cong-ev[of - $\lambda x. \ln\ (\mathbb{R} * l) - \ln\ (\mathbb{R} * l)$], simp-all*)

apply (*rewrite deriv-diff, simp-all*)

unfolding *field-differentiable-def* **using** *that*

by (metis DERIV-lm-divide-chain lx-pos real-differentiable-def)
 thus ?thesis using cdfX-l-normal mu-deriv-lm that differentiable-cdfX-l by force
 qed

lemma deriv-l-shift: deriv l (x+t) = deriv (λt. \$l(x+t)) t
 if l differentiable at (x+t) for x t :: real
 using deriv-shift differentiable-eq-field-differentiable-real that by simp

context
 fixes x::real
 assumes x-lt-psi[simp]: x < \$ψ
 begin

lemma p-mu-l: \$p-{t&x} * \$μ-(x+t) = - deriv l (x+t) / \$l-x
 if l differentiable at (x+t) t > 0 x+t < \$ψ for t::real
 using p-l mu-deriv-l that by simp

lemma p-mu-l-AE: AE s in lborel. 0 < s ∧ x+s < \$ψ → \$p-{s&x} * \$μ-(x+s)
 = - deriv l (x+s) / \$l-x

proof -
 have AE s in lborel. l differentiable at (x+s)
 apply (rule AE-borel-affine[of 1 λu. l differentiable at u x, simplified])
 unfolding pred-def using l-differentiable-AE by simp-all
 moreover have AE s in lborel.
 l differentiable at (x+s) → 0 < s ∧ x+s < \$ψ → \$p-{s&x} * \$μ-(x+s) =
 - deriv l (x+s) / \$l-x
 using p-mu-l by (intro AE-I2) simp
 ultimately show ?thesis by (rule AE-mp)
 qed

lemma LBINT-l-mu-q: (LBINT s:{f<..f+t}. \$l-(x+s) * \$μ-(x+s)) / \$l-x = \$q-{f|t&x}
 if t ≥ 0 f ≥ 0 for t f :: real

proof -
 have ∧s. s∈{f<..f+t} ⇒ \$p-{s&x} = \$l-(x+s) / \$l-x using p-l that by simp
 hence \$q-{f|t&x} = (LBINT s:{f<..f+t}. \$l-(x+s) / \$l-x * \$μ-(x+s))
 using LBINT-p-mu-q-defer
 by (smt (verit) greaterThanAtMost-borel set-lebesgue-integral-cong sets-lborel
 that x-lt-psi)
 also have ... = (LBINT s:{f<..f+t}. \$l-(x+s) * \$μ-(x+s)) / \$l-x
 using set-integral-divide-zero by simp
 finally show ?thesis by simp
 qed

lemma set-integrable-l-mu: set-integrable lborel {f<..f+t} (λs. \$l-(x+s) * \$μ-(x+s))
 if t ≥ 0 f ≥ 0 for t f :: real

proof -
 have set-integrable lborel {f<..f+t} (λs. \$l-(x+s) * \$μ-(x+s) / \$l-x)
 using p-l set-integrable-p-mu that
 by (rewrite set-integrable-cong[where f'=λs. \$p-{s&x} * \$μ-(x+s)]) simp-all+

thus ?thesis by simp
qed

lemma *l-mu-has-integral-q-defer*:

$((\lambda s. \$l(x+s) * \$\mu-(x+s) / \$l-x) \text{ has-integral } \$q-\{f|t\&x\}) \{f..f+t\}$

if $t \geq 0$ $f \geq 0$ for $t f :: \text{real}$

using *p-l* that *p-mu-has-integral-q-defer-Icc*

by (rewrite has-integral-cong[of - - $\lambda s. \$p-\{s\&x\} * \$\mu-(x+s)$]; simp)

corollary *l-mu-has-integral-q*:

$((\lambda s. \$l(x+s) * \$\mu-(x+s) / \$l-x) \text{ has-integral } \$q-\{t\&x\}) \{0..t\}$ if $t \geq 0$ for $t::\text{real}$

using *l-mu-has-integral-q-defer*[where $f=0$] that by simp

lemma *l-mu-has-integral-d*:

$((\lambda s. \$l(x+s) * \$\mu-(x+s)) \text{ has-integral } \$d-\{t \& x+f\}) \{f..f+t\}$

if $t \geq 0$ $f \geq 0$ for $t f :: \text{real}$

proof –

have $((\lambda s. \$l-x * (\$p-\{s\&x\} * \$\mu-(x+s))) \text{ has-integral } \$l-x * \$q-\{f|t\&x\}) \{f..f+t\}$

apply (rule has-integral-mult-right)

by (rule *p-mu-has-integral-q-defer-Icc*; simp add: that)

thus ?thesis

using that apply (rewrite in asm *q-defer-d-l*, simp-all)

apply (rewrite has-integral-cong[where $g=\lambda s. \$l-x * (\$p-\{s\&x\} * \$\mu-(x+s))$])

by (rewrite *p-l*; simp)

qed

corollary *l-mu-has-integral-d-1*:

$((\lambda s. \$l(x+s) * \$\mu-(x+s)) \text{ has-integral } \$d-(x+f)) \{f..f+1\}$ if $t \geq 0$ $f \geq 0$ for $t f :: \text{real}$

using *l-mu-has-integral-d*[where $t=1$] that by simp

lemma *e-LBINT-l*: $\$e\circ-x = (\text{LBINT } s:\{0..\}. \$l(x+s) * \$\mu-(x+s) * s) / \$l-x$

— Note that $0 = 0$ holds when the life expectation diverges.

proof –

have $\bigwedge s. s \in \{0..\} \implies \$p-\{s\&x\} = \$l(x+s) / \$l-x$ using *p-l* by simp

hence $\$e\circ-x = (\text{LBINT } s:\{0..\}. \$l(x+s) / \$l-x * \$\mu-(x+s) * s)$

using *e-LBINT-p-mu*

by (smt (verit) *atLeast-borel set-lebesgue-integral-cong sets-lborel x-lt-psi*)

also have $\dots = (\text{LBINT } s:\{0..\}. \$l(x+s) * \$\mu-(x+s) * s) / \$l-x$

using *set-integral-divide-zero* by simp

finally show ?thesis .

qed

lemma *e-integral-l*: $\$e\circ-x = \text{integral } \{0..\} (\lambda s. \$l(x+s) * \$\mu-(x+s) * s) / \$l-x$

— Note that $0 = 0$ holds when the life expectation diverges.

proof –

have *AE* s in *lborel*. $\$\mu-(x+s) \geq 0$ by (rule *AE-translation*, rule *mu-nonneg-AE*)

hence $(\text{LBINT } s:\{0..\}. \$l(x+s) * \$\mu-(x+s) * s) = \text{integral } \{0..\} (\lambda s. \$l(x+s) * \$\mu-(x+s) * s)$

by (intro set-borel-integral-eq-integral-nonneg-AE; force)
thus ?thesis using e-LBINT-l by simp
qed

lemma m-LBINT-p-mu: $\$m\text{-}\{n\&x\} = (LBINT\ t:\{0<..n\}.\ \$p\text{-}\{t\&x\} * \$\mu\text{-}(x+t))$
 $/ (LBINT\ t:\{0..n\}.\ \$p\text{-}\{t\&x\})$
if $n \geq 0$ for $n::real$
using that
apply (rewrite m-q-e, simp-all)
apply (rewrite LBINT-p-mu-q[simplified], simp-all)
by (rewrite temp-e-LBINT-p; simp)

lemma m-integral-p-mu:
 $\$m\text{-}\{n\&x\} = integral\ \{0..n\}\ (\lambda t.\ \$p\text{-}\{t\&x\} * \$\mu\text{-}(x+t)) / integral\ \{0..n\}\ (\lambda t.\$
 $\$p\text{-}\{t\&x\})$
if $n \geq 0$ for $n::real$
using that
apply (rewrite m-q-e, simp-all)
apply (rewrite integral-unique[OF p-mu-has-integral-q-Icc])
apply simp-all[2]
by (rewrite temp-e-integral-p; simp)

end

lemma deriv-x-p-mu-l: $deriv\ (\lambda y.\ \$p\text{-}\{t\&y\})\ x = \$p\text{-}\{t\&x\} * (\$\mu\text{-}x - \$\mu\text{-}(x+t))$
if l differentiable at x l differentiable at $(x+t)$ $t \geq 0$ $x < \$\psi$ for $x\ t :: real$
using deriv-x-p-mu that differentiable-ccdfX-l by blast

lemma e-has-derivative-mu-e-l: $((\lambda x.\ \$e\text{'o-}x)$ has-real-derivative $(\$\mu\text{-}x * \$e\text{'o-}x -$
 $1))$ (at x)
if total-finite l differentiable at x $x \in \{a <..< b\}$ $b \leq \$\psi$ for $a\ b\ x :: real$
using total-finite-iff-set-integrable-Ici that
e-has-derivative-mu-e differentiable-ccdfX-l set-integrable-ccdfX-l
by force

corollary e-has-derivative-mu-e-l': $((\lambda x.\ \$e\text{'o-}x)$ has-real-derivative $(\$\mu\text{-}x * \$e\text{'o-}x$
 $- 1))$ (at x)
if total-finite l differentiable at x $x \in \{a <..< b\}$ $b \leq \$\psi$ for $a\ b\ x :: real$
using that by (intro e-has-derivative-mu-e-l[where a=a]; simp)

context

fixes $x::real$
assumes $x\text{-lt-psi}$ [simp]: $x < \$\psi$
begin

lemma curt-e-sum-l-smooth: $\$e\text{-}x = (\sum k.\ \$l\text{-}(x+k+1)) / \$l\text{-}x$ if total-finite
proof –
have [simp]: summable $(\lambda k.\ \$l\text{-}(x+k+1))$
using total-finite-iff-summable[of $x+1$] that

by (*metis* (*no-types*, *lifting*) *add.commute add.left-commute summable-def sums-cong*)
 hence *summable* ($\lambda k. \$p-\{k+1\&x\}$) by (*rewrite p-l; simp add: add.assoc*)
 hence $\$e-x = (\sum k. \$p-\{k+1\&x\})$ using *curt-e-sum-p-smooth* by *simp*
 also have $\dots = (\sum k. \$l-(x+k+1) / \$l-x)$ by (*rewrite p-l; simp add: add.assoc*)
 also have $\dots = (\sum k. \$l-(x+k+1)) / \$l-x$ by (*rewrite suminf-divide; simp*)
 finally show *?thesis* .
 qed

lemma *curt-e-sum-l-finite-smooth*: $\$e-x = (\sum k < n. \$l-(x+k+1)) / \$l-x$ if $x+n+1 > \$\psi$ for $n::nat$
 apply (*rewrite curt-e-sum-p-finite-smooth[of x n], simp-all add: that*)
 apply (*rewrite p-l, simp-all*)
 by (*smt (verit) sum.cong sum-divide-distrib*)

lemma *temp-curt-e-sum-l-smooth*: $\$e-\{x:n\} = (\sum k < n. \$l-(x+k+1)) / \$l-x$ for $n::nat$
 apply (*rewrite temp-curt-e-sum-p-smooth[of x n], simp*)
 apply (*rewrite p-l, simp-all*)
 by (*smt (verit) sum.cong sum-divide-distrib*)

end

end

6.4 Interpolations

context *life-table*
begin

definition *linear-interpolation* \equiv
 $\forall (x::nat)(t::real). 0 \leq t \wedge t \leq 1 \longrightarrow \$l-(x+t) = (1-t)*\$l-x + t*\$l-(x+1)$

lemma *linear-l*: $\$l-(x+t) = (1-t)*\$l-x + t*\$l-(x+1)$
 if *linear-interpolation* $0 \leq t \leq 1$ for $x::nat$ and $t::real$
 using *that unfolding linear-interpolation-def* by (*metis of-nat-1 of-nat-add*)

lemma *linear-l-d*: $\$l-(x+t) = \$l-x - t*\$d-x$
 if *linear-interpolation* $0 \leq t \leq 1$ for $x::nat$ and $t::real$
 using *death-def-nonneg that unfolding linear-interpolation-def*
 by (*smt (verit) distrib-left left-diff-distrib'*)

lemma *linear-p-q*: $\$p-\{t\&x\} = 1 - t*\$q-x$
 if *linear-interpolation* $0 \leq t \leq 1$ $x < \$\psi$ for $x::nat$ and $t::real$
 using *that*
 apply (*rewrite p-l, simp-all*)
 apply (*rewrite q-d-l, simp-all*)
 using *divide-self[of \$l-(real x)] linear-l-d*
 by (*smt (verit, ccfv-SIG) add-divide-distrib lx-neq-0*)

lemma *linear-q*: $\$q\{t\&x\} = t*\$q\text{-}x$
if *linear-interpolation* $0 \leq t \leq 1$ $x < \$\psi$ **for** $x::\text{nat}$ **and** $t::\text{real}$
using *that linear-p-q p-q-1* **by** (*smt (verit)*)

lemma *linear-L-l-d*: $\$L\text{-}x = \$l\text{-}x - \$d\text{-}x / 2$ **if** *linear-interpolation* **for** $x::\text{nat}$
proof –
have $\$L\text{-}(real\ x) = (LBINT\ t:\{0..1\}. \$l\text{-}(real\ x + t))$
unfolding *lives-def* **using** *lborel-set-integral-Icc-shift[of real x real x + 1 l real x]*
by (*simp add: add commute*)
also have $\dots = (LBINT\ t:\{0..1\}. \$l\text{-}(real\ x) - t*\$d\text{-}(real\ x))$
using *linear-l-d* **that** **by** (*intro set-lebesgue-integral-cong; simp*)
also have $\dots = \$l\text{-}(real\ x) - \$d\text{-}(real\ x) / 2$
proof –
have $(LBINT\ t:\{0::real..1\}. t) = 1/2$
unfolding *set-lebesgue-integral-def* **using** *integral-power[of 0 1 1]* **by** (*simp add: mult commute*)
hence $(LBINT\ t:\{0..1\}. t*\$d\text{-}(real\ x)) = \$d\text{-}(real\ x) / 2$ **by** *auto*
moreover have *set-integrable lborel {0..1} ($\lambda t. t*\$d\text{-}(real\ x)$)*
apply (*rule set-integrable-mult-left[where f=id and a=\$d-(real x), simplified]*)
unfolding *set-integrable-def* **using** *integrable-power[of 0 1 1]* **by** (*simp add: mult commute*)
moreover have $(LBINT\ t:\{0::real..1\}. \$l\text{-}(real\ x)) = \$l\text{-}(real\ x)$
unfolding *set-lebesgue-integral-def* **by** *simp*
ultimately show *?thesis* **using** *set-integrable-def* **by** (*rewrite set-integral-diff; force*)
qed
finally show *?thesis* .
qed

lemma *linear-L-l-d'*: $\$L\text{-}x = \$l\text{-}(x+1) + \$d\text{-}x / 2$ **if** *linear-interpolation* **for** $x::\text{nat}$
proof –
have $\$L\text{-}(real\ x) = \$l\text{-}(real\ x) - \$d\text{-}(real\ x) + \$d\text{-}(real\ x) / 2$ **using** *that linear-L-l-d* **by** *simp*
also have $\dots = \$l\text{-}(real\ (x+1)) + \$d\text{-}(real\ x) / 2$ **using** *dx-l* **by** (*smt (verit) of-nat-1 of-nat-add*)
finally show *?thesis* .
qed

lemma *linear-l-continuous*: *continuous-on UNIV l* **if** *linear-interpolation*
unfolding *continuous-on-def*
proof
fix $u::\text{real}$
show $l - u \rightarrow \$l\text{-}u$
proof (*cases $\langle u \leq 0 \rangle$*)
case *True*
hence $(l \longrightarrow \$l\text{-}u)$ (*at-left u*) **using** *l-left-continuous-nonpos continuous-within*
by *auto*
thus *?thesis*

```

    apply (rule filterlim-split-at-real)
    using l-right-continuous continuous-within by auto
next
case False
hence u-pos: u > 0 by simp
thus ?thesis
proof (cases ⟨u = real-of-int [u]⟩)
  case True
  from this u-pos obtain x::nat where ux: u = Suc x
  by (metis gr0-implies-Suc of-int-0-less-iff of-int-of-nat-eq pos-int-cases)
  have ((λt. (1-t)*$l-(real x) + t*$l-(real x + 1)) → $l-(real x + 1)) (at-left
1)
    apply (rewrite in (- → □) - add-0[THEN sym], rule tendsto-add)
    apply (simp add: LIM-zero-iff' tendsto-mult-left-zero)
    by (rewrite in (- → □) - mult-1[THEN sym], rule tendsto-mult-right)
simp
hence ((λt. $l-(real x + t)) → $l-(real x + 1)) (at-left 1)
  apply (rewrite tendsto-cong; simp)
  apply (rule eventually-at-leftI[of 0]; simp)
  using that by (rewrite linear-l; simp add: add.commute)
moreover have ((λt. $l-(real x + t)) → $l-(real x + 1)) (at-right 1)
  using l-right-continuous apply (rule continuous-within-tendsto-compose,
simp-all)
  apply (rule eventually-at-right-less)
  by (rule tendsto-intros, simp-all)
ultimately show ?thesis
  apply (rewrite ux)+
  apply (rewrite filterlim-shift-iff[where d=x, THEN sym])
  by (rule filterlim-split-at-real; simp add: comp-def add.commute)
next
case False
let ?x = nat [u]
let ?t = u - real ?x
let ?e = min ?t (1 - ?t)
from False have ?t > 0 using u-pos by linarith
moreover have ?t < 1 by linarith
ultimately have e-pos: ?e > 0 by simp
hence
  ∀F v in nhds u. $l-v = (1 - (v - real ?x))*$l-(real ?x) + (v - real ?x)*$l-(real
?x + 1)
proof -
  { fix v::real assume vu-e: dist v u < ?e
    hence v - real ?x ≥ 0 using dist-real-def by force
    moreover have v - real ?x ≤ 1 using dist-real-def vu-e by force
    ultimately have $l-v = (1 - (v - real ?x))*$l-(real ?x) + (v - real
?x)*$l-(real ?x + 1)
      using linear-l that by (smt (verit, ccfv-threshold) linear-interpolation-def)
  }
  thus ?thesis using eventually-nhds-metric e-pos by blast

```


qed
moreover have
 $isCont (\lambda v. (1 - (v - real ?x)) * \$l-(real ?x) + (v - real ?x) * \$l-(real ?x + 1)) u$
by (*rule continuous-intros*) +
ultimately have $isCont l u$ **using** $isCont-cong$ **by force**
thus $?thesis$ **by** (*simp add: isContD*)
qed
qed
qed

lemma $linear-l-sums-T-l: (\lambda k. \$l-(x + Suc k)) \text{ sums } (\$T-x - \$l-x / 2)$
if $linear-interpolation \text{ total-finite}$ **for** $x::nat$
proof –
have $\bigwedge k::nat. \$l-(real (x + Suc k)) = \$L-(real (x+k)) - \$d-(real (x+k)) / 2$
using $linear-L-l-d'$ **that** **by** (*smt (verit) Suc-eq-plus1 add-Suc-right*)
moreover have $(\lambda k::nat. \$L-(real (x+k))) \text{ sums } \$T-x$ **using** $L-sums-T$ **that** **by** *simp*
moreover have $(\lambda k::nat. \$d-(real (x+k)) / 2) \text{ sums } (\$l-(real x) / 2)$
using $sums-divide d-sums-l$ **by auto**
ultimately show $?thesis$
apply (*rewrite sums-cong, simp*)
by (*rule sums-diff; simp*)
qed

corollary $linear-T-suminf-l: \$T-x = (\sum k. \$l-(x+k+1)) + \$l-x / 2$
if $linear-interpolation \text{ total-finite}$ **for** $x::nat$
using $linear-l-sums-T-l$ **that** $sums-unique$ **by** (*smt (z3) Suc-eq-plus1 add-Suc-right suminf-cong*)

lemma $linear-mx-q: \$m-x = \$q-x / (1 - \$q-x / 2)$ **if** $linear-interpolation x < \ψ
for $x::nat$
proof –
have [*simp*]: $\$l-(real x) \neq 0$ **using** $that$ **by** *simp*
have $\$m-(real x) = \$d-(real x) / (\$l-(real x) - \$d-(real x) / 2)$
unfolding $die-central-def$ **using** $linear-L-l-d$ **that** **by** *simp*
also have $\dots = (\$d-(real x) / \$l-(real x)) / ((\$l-(real x) - \$d-(real x) / 2) / \$l-(real x))$
by *simp*
also have $\dots = (\$d-(real x) / \$l-(real x)) / (1 - (\$d-(real x) / \$l-(real x)) / 2)$
by (*rewrite diff-divide-distrib*) *simp*
also have $\dots = \$q-(real x) / (1 - \$q-(real x) / 2)$ **using** $that q-d-l$ **by** *simp*
finally show $?thesis$.
qed

lemma $linear-e-curt-e: \$e' \circ -x = \$e-x + 1/2$
if $linear-interpolation \text{ total-finite}$ $x < \$\psi$ **for** $x::nat$
proof –
have $\$e' \circ -(real x) = ((\sum k::nat. \$l-(real (x+k+1))) + \$l-(real x) / 2) / \$l-(real$

x)
using *e-T-l linear-T-suminf-l* **that by** *simp*
also have $\dots = (\sum k::nat. \$l(\text{real } (x+k+1))) / \$l(\text{real } x) + (\$l(\text{real } x) / 2) / \$l(\text{real } x)$
using *add-divide-distrib* **by** *blast*
also have $\dots = \$e(\text{real } x) + 1/2$
using *that apply* (*rewrite curt-e-sum-l, simp-all*)
using *linear-l-continuous* **by** (*rule continuous-on-interior, simp-all add: that add.commute*)
finally show *?thesis* .
qed
end

context *smooth-life-table*
begin

lemma *linear-l-has-derivative-at-frac*:
 $((\lambda s. \$l(x+s)) \text{ has-real-derivative } - \$d-x) \text{ (at } t)$
if *linear-interpolation* $0 < t < 1$ **for** $x::nat$ **and** $t::real$
proof –
let $?x = \text{real } x$
have $((\lambda s. \$l-?x - s*\$d-?x) \text{ has-real-derivative } (0 - \$d-?x)) \text{ (at } t)$
apply (*rule derivative-intros, simp*)
apply (*rule DERIV-cmult-right*[*of id 1, simplified*])
by (*metis DERIV-ident eq-id-iff*)
moreover have $\forall_F s \text{ in nhds } t. \$l-?x + s = \$l-?x - s*\$d-?x$
proof –
let $?r = \min t (1-t)$
have $?r > 0$ **using** *that* **by** *simp*
moreover
{ **fix** s **assume** $\text{dist } s t < ?r$
hence $\$l-?x + s = \$l-?x - s*\$d-?x$ **using** *linear-l-d* **that** *dist-real-def* **by**
force **}**
ultimately show *?thesis* **using** *eventually-nhds-metric* **that by** *blast*
qed
ultimately show *?thesis* **by** (*rewrite DERIV-cong-ev; simp*)
qed

lemma *linear-l-has-derivative-at-frac'*:
 $(l \text{ has-real-derivative } - \$d-x) \text{ (at } y)$
if *linear-interpolation* $x < y < x+1$ **for** $x::nat$ **and** $y::real$
apply (*rewrite DERIV-at-within-shift*[**where** $x=y$ – *real x* **and** $z=\text{real } x$ **and** $S=UNIV$, *simplified*])
using *linear-l-has-derivative-at-frac* **that by** *simp*

lemma *linear-l-differentiable-on-frac*:
 $l \text{ differentiable-on } \{x <.. < x+1\}$ **if** *linear-interpolation* **for** $x::nat$
proof –

```

{ fix y::real assume y ∈ {real x <..  

  hence l differentiable at y  

  using linear-l-has-derivative-at-frac' that real-differentiable-def by auto }  

thus ?thesis unfolding differentiable-on-def by (metis differentiable-at-withinI)  

qed

```

lemma *linear-l-has-right-derivative-at-nat*:

(*l* has-real-derivative $- \$d\cdot x$) (at-right *x*) if linear-interpolation for *x*::nat

proof –

```

let ?x = real x  

have [simp]: plus ?x ‘ {0<..  

  unfolding image-def greaterThan-def apply simp  

  by (metis Groups.ab-semigroup-add-class.add commute  

    add-minus-cancel neg-less-iff-less real-0-less-add-iff)  

have ((λs. $l-(?x + s)) has-real-derivative  $- \$d\cdot ?x$ ) (at-right 0)  

apply (rewrite has-field-derivative-cong-eventually[where g=λs. $l-?x - s*$d-?x])  

using linear-l-d that apply (intro eventually-at-rightI[of 0 1], simp-all)  

apply (rule has-field-derivative-at-within)  

apply (rewrite diff-0[of $d-?x, THEN sym])  

apply (rule DERIV-diff, simp)  

apply (rule DERIV-cmult-right[of id 1, simplified])  

by (metis DERIV-ident eq-id-iff)  

thus ?thesis  

  by (rewrite DERIV-at-within-shift[where z=?x and x=0 and S={0<..  

simplified]) simp  

qed

```

lemma *linear-l-has-left-derivative-at-nat*:

(*l* has-real-derivative $- \$d\cdot(\text{real } x - 1)$) (at-left *x*) if linear-interpolation for *x*::nat

proof (cases *x*)

case 0

hence (*l* has-real-derivative 0) (at-left (real *x*))

apply (rewrite has-field-derivative-cong-eventually[where g=λ-. \$l-0]; simp)

apply (rule eventually-at-leftI[of -1]; simp)

using l-neg-nil less-eq-real-def by blast

moreover have $\$d\cdot(\text{real } x - 1) = 0$ using 0 dx-l l-neg-nil less-eq-real-def by fastforce

ultimately show ?thesis by auto

next

case (Suc *y*)

let ?x = real *x* and ?y = real *y*

have [simp]: ?y + 1 = ?x using Suc by simp

have [simp]: plus ?y ‘ {..
 using Suc unfolding image-def lessThan-def apply simp

by (metis (no-types, opaque-lifting) Groups.ab-semigroup-add-class.add commute
 add-less-cancel-right diff-add-cancel)

have $\$l\cdot ?x = \$l\cdot \text{real } y - \$d\cdot \text{real } y$ using Suc by (simp add: dx-l)

moreover have $\forall_F s$ in at-left 1. $\$l\cdot(?y + s) = \$l\cdot ?y - s*\$d\cdot ?y$

apply (rule eventually-at-leftI[of 0], simp-all)
using Suc linear-l-d that **by** simp
moreover have (($\lambda s. \$l-?y - s*\$d-?y$) has-real-derivative - $\$d-?y$) (at-left 1)
apply (rule has-field-derivative-at-within)
apply (rewrite diff-0[of $\$d-?y$, THEN sym])
apply (rule DERIV-diff, simp)
apply (rule DERIV-cmult-right[of id 1, simplified])
by (metis DERIV-ident eq-id-iff)
ultimately have (($\lambda s. \$l-(?y + s)$) has-real-derivative - $\$d-?y$) (at-left 1)
by (rewrite has-field-derivative-cong-eventually[**where** $g=\lambda s. \$l-?y - s*\$d-?y$];
simp)
thus ?thesis
apply (rewrite DERIV-at-within-shift[**where** $S=\{..<1\}$ and $z=?y$ and $x=1$,
simplified])
using Suc **by** simp
qed

lemma linear-l-has-derivative-at-nat-iff-d:

(l has-real-derivative - $\$d-x$) (at x) \longleftrightarrow $\$d-x = \$d-(\text{real } x - 1)$
if linear-interpolation **for** $x::\text{nat}$

proof -

let ? $x = \text{real } x$

have (l has-real-derivative - $\$d-?x$) (at ? x) \longleftrightarrow

(l has-real-derivative - $\$d-?x$) (at-right ? x) \wedge

(l has-real-derivative - $\$d-?x$) (at-left ? x)

using has-real-derivative-at-split **by** auto

also have ... \longleftrightarrow $\$d-?x = \$d-(?x - 1)$ (**is** ?LHS = ?RHS)

proof

assume ?LHS

hence (l has-real-derivative - $\$d-?x$) (at-left ? x) **by** simp

moreover have (l has-real-derivative - $\$d-(?x - 1)$) (at-left ? x)

using that linear-l-has-left-derivative-at-nat **by** simp

ultimately show ?RHS

using has-real-derivative-iff-has-vector-derivative vector-derivative-unique-within
trivial-limit-at-left-real **by** (smt (verit, ccfv-SIG))

next

assume ?RHS

thus ?LHS

using that linear-l-has-right-derivative-at-nat linear-l-has-left-derivative-at-nat

by metis

qed

finally show ?thesis .

qed

lemma linear-l-differentiable-at-nat-iff-d:

l differentiable at x \longleftrightarrow $\$d-x = \$d-(\text{real } x - 1)$

if linear-interpolation **for** $x::\text{nat}$

proof

let ? $x = \text{real } x$

assume l differentiable at x
from this obtain D **where** $DERIV$ - l : (l has-real-derivative D) (at $?x$)
using $real$ -differentiable-def **by** $blast$
hence (l has-real-derivative D) (at-right $?x$)
using has -field-derivative-at-within **by** $blast$
moreover have at $?x$ within $\{real\ x < ..\} \neq \perp$ **by** $simp$
moreover have (l has-real-derivative $- \$d\ ?x$) (at-right $?x$)
using $linear$ - l -has-right-derivative-at-nat that **by** $simp$
ultimately have $D = - \$d\ ?x$
using that has-real-derivative-iff-has-vector-derivative $vector$ -derivative-unique-within
by $blast$
thus $\$d\ ?x = \$d\ (?x - 1)$
using $linear$ - l -has-derivative-at-nat-iff-d that $DERIV$ - l **by** $blast$
next
assume $\$d\ (real\ x) = \$d\ (real\ x - 1)$
thus l differentiable at ($real\ x$)
using $linear$ - l -has-derivative-at-nat-iff-d that $real$ -differentiable-def **by** $blast$
qed

lemma $linear$ - l -limited: $\$ \psi < \infty$ **if** $linear$ -interpolation
proof –
let $?ND = \{y. \neg l\ \text{differentiable at } y\}$
obtain $xn::nat$ **where** xn -def: $Max\ ?ND < real\ xn$ **using** $reals$ -Archimedean2
by $blast$
hence xn -dif: $\bigwedge x::nat. x \geq xn \implies l\ \text{differentiable at } (real\ x)$
proof –
fix $x::nat$ **assume** $x \geq xn$
with xn -def **have** $real\ x > Max\ ?ND$ **by** $simp$
hence $real\ x \notin ?ND$ **using** $notI\ Max.coboundedI\ l$ -nondifferentiable-finite-set
 leD **by** $blast$
thus l differentiable at ($real\ x$) **by** $simp$
qed
hence d -const: $\bigwedge x::nat. x \geq xn \implies \$d\ (real\ x) = \$d\ (real\ xn)$
proof –
fix $x::nat$ **assume** $x \geq xn$
moreover have
 $\bigwedge y::nat. xn \leq y \implies \$d\ (real\ y) = \$d\ (real\ xn) \implies \$d\ (real\ (Suc\ y)) = \$d\ (real\ xn)$
using $linear$ - l -differentiable-at-nat-iff-d that
by ($smt\ (verit, best)\ of$ -nat-Suc of -nat-le-iff xn -dif)
ultimately show $\$d\ (real\ x) = \$d\ (real\ xn)$
using nat -induct-at-least[**where** $P = \lambda x::nat. \$d\ (real\ x) = \$d\ (real\ xn)$] **by**
 $simp$
qed
have $\neg \$d\ (real\ xn) > 0$
proof ($rule\ notI$)
assume $\$d\ (real\ xn) > 0$
from this obtain $N::nat$ **where** N -def: $N * \$d\ (real\ xn) > \$l\ (real\ xn)$
using $reals$ -Archimedean3 **by** $blast$

hence $\$l(\text{real } (xn+N)) < 0$
 proof –
 have $\$l(\text{real } (xn+N)) = \$l(\text{real } xn) - (\sum k < N. \$d(\text{real } xn + \text{real } k))$ **using**
sum-dx-l **by** *simp*
 also have $\dots = \$l(\text{real } xn) - (\sum k < N. \$d(\text{real } xn))$
 using *d-const* **by** (*metis le-add1 of-nat-add*)
 also have $\dots = \$l(\text{real } xn) - N * \$d(\text{real } xn)$ **by** *simp*
 also have $\dots < 0$ **using** *N-def* **by** *simp*
 finally show *?thesis* .
 qed
 thus *False* **by** (*smt (verit, ccfv-SIG) l-nonneg*)
 qed
 hence *dxn0*: $\$d(\text{real } xn) = 0$ **by** (*smt (verit) d-nonneg*)
 hence $\bigwedge x::\text{nat}. x \geq xn \implies \$l(\text{real } x) = \$l(\text{real } xn)$
 proof –
 fix *x::nat*
 assume $xn \leq x$
 moreover have
 $\bigwedge y::\text{nat}. xn \leq y \implies \$l(\text{real } y) = \$l(\text{real } xn) \implies \$l(\text{real } (\text{Suc } y)) = \$l(\text{real } xn)$
by (*smt (verit, ccfv-threshold) dxn0 d-const dx-l of-nat-Suc*)
 ultimately show $\$l(\text{real } x) = \$l(\text{real } xn)$
 using *nat-induct-at-least*[**where** $P = \lambda x::\text{nat}. \$l(\text{real } x) = \$l(\text{real } xn)$] **by**
simp
 qed
 hence $(\lambda x::\text{nat}. \$l(\text{real } x)) \longrightarrow \$l(\text{real } xn)$
 using *eventually-sequentially* **by** (*intro tendsto-eventually*) *blast*
 moreover have $(\lambda x::\text{nat}. \$l(\text{real } x)) \longrightarrow 0$
 using *l-PInfty-0* **by** (*simp add: filterlim-compose filterlim-real-sequentially*)
 ultimately have $\$l(\text{real } xn) = 0$ **by** (*simp add: LIMSEQ-unique*)
 thus *?thesis* **by** *force*
 qed
lemma *linear-mu-q*: $\$ \mu(x+t) = \$q-x / (1 - t*\$q-x)$
 if *linear-interpolation l differentiable at (x+t)* $0 < t < 1$ $x+t < \$\psi$
 for *x::nat* and *t::real*
 proof –
 have [*simp*]: $\text{ereal } x < \$\psi$ **using** *that* **by** (*simp add: ereal-less-le*)
 have [*simp*]: $\$l(\text{real } x) \neq 0$ **by** *simp*
 have [*simp*]: *l field-differentiable at (real x + t)*
 using *differentiable-eq-field-differentiable-real* **that** **by** *simp*
 define *d* where $d \equiv \min t (1-t)$
 have *d-pos*: $d > 0$ **unfolding** *d-def* **using** *that* **by** *simp*
 have $\$p\{t \ \& \ \text{real } x\} \neq 0$ **using** *that* **by** (*simp add: ereal-less-le p-0-equiv*)
 moreover have $(\lambda s. \$p\{s \ \& \ \text{real } x\})$ *differentiable at t*
 proof –
 have $(\lambda s. \$l(\text{real } x + s) / \$l(\text{real } x))$ *field-differentiable at t*
 using *that* **apply** (*intro derivative-intros*)
 apply (*rewrite add commute, rewrite field-differentiable-shift*[*THEN sym*])

by (rewrite add commute) simp-all
 thus ?thesis
 apply (rewrite differentiable-eq-field-differentiable-real)
 apply (rule field-differentiable-transform-within[where d=d], simp-all add:
 d-pos)
 apply (rewrite p-l, simp-all) unfolding d-def using dist-real-def by auto
 qed
 ultimately have $\$ \mu\text{-}(\text{real } x + t) = - \text{deriv } (\lambda s. \$ p\text{-}\{s \ \& \ \text{real } x\}) t / \$ p\text{-}\{t \ \& \ \text{real } x\}$
 using that deriv-t-p-mu by simp
 also have $\dots = \$ q\text{-}(\text{real } x) / (1 - t * \$ q\text{-}(\text{real } x))$
 proof -
 have $\bigwedge s. \text{dist } s \ t < d \implies \$ p\text{-}\{s \ \& \ \text{real } x\} = 1 - s * \$ q\text{-}(\text{real } x)$
 proof -
 fix s assume dist s t < d
 hence $0 \leq s \leq 1$ unfolding d-def using that dist-real-def by auto
 thus $\$ p\text{-}\{s \ \& \ \text{real } x\} = 1 - s * \$ q\text{-}(\text{real } x)$ by (intro linear-p-q; simp add: that)
 qed
 hence $\text{deriv } (\lambda s. \$ p\text{-}\{s \ \& \ \text{real } x\}) t = \text{deriv } (\lambda s. 1 - s * \$ q\text{-}(\text{real } x)) t$
 using d-pos apply (intro deriv-cong-ev, simp-all)
 by (rewrite eventually-nhds-metric) auto
 also have $\dots = - \$ q\text{-}(\text{real } x)$
 apply (rewrite deriv-diff, simp-all)
 by (rule derivative-intros) auto
 finally have $\text{deriv } (\lambda s. \$ p\text{-}\{s \ \& \ \text{real } x\}) t = - \$ q\text{-}(\text{real } x)$.
 thus ?thesis using linear-p-q that by simp
 qed
 finally show ?thesis .
 qed

definition exponential-interpolation \equiv
 $\forall (x::\text{nat})(t::\text{real}). x+1 < \$ \psi \longrightarrow 0 \leq t \wedge t < 1 \longrightarrow \$ \mu\text{-}(x+t) = \$ \mu\text{-}x$
 — Without $x+1 < \$ \psi$, the smooth life table could not be limited.

lemma exponential-mu: $\$ \mu\text{-}(x+t) = \$ \mu\text{-}x$
 if exponential-interpolation $x+1 < \$ \psi$ $0 \leq t < 1$ for $x::\text{nat}$ and $t::\text{real}$
 using that unfolding exponential-interpolation-def by simp

corollary exponential-mu': $\$ \mu\text{-}y = \$ \mu\text{-}x$
 if exponential-interpolation $x \leq y$ $y < x+1$ $x+1 < \$ \psi$ for $x::\text{nat}$ and $y::\text{real}$
 proof -
 let ?t = y - real x
 have $0 \leq ?t$ and $?t < 1$ using that by simp-all
 moreover have $\$ \mu\text{-}y = \$ \mu\text{-}(\text{real } x + ?t)$ by simp
 ultimately show ?thesis using exponential-mu that by presburger
 qed

lemma exponential-integral-mu: $\text{integral } \{x..<x+t\} (\lambda y. \$ \mu\text{-}y) = \$ \mu\text{-}x * t$
 if exponential-interpolation $x+1 < \$ \psi$ $0 \leq t < 1$ for $x::\text{nat}$ and $t::\text{real}$

proof –

have $\text{integral } \{ \text{real } x \text{ ..} < \text{real } x + t \} (\lambda y. \$\mu\text{-}y) =$
 $\text{integral } \{ \text{real } x \text{ ..} < \text{real } x + t \} (\lambda y. \$\mu\text{-}(\text{real } x))$
using *exponential-mu'* **that by** (*intro integral-cong; simp*)
also have $\dots = \text{integral } \{ \text{real } x \text{ ..} \text{real } x + t \} (\lambda y. \$\mu\text{-}(\text{real } x))$
apply (*rule integral-subset-negligible, force*)
using that by (*rewrite Icc-minus-Ico; simp*)
also have $\dots = \$\mu\text{-}x * t$ **using that by** (*rewrite integral-const-real*) *simp*
finally show $\text{integral } \{ \text{real } x \text{ ..} < \text{real } x + t \} (\lambda y. \$\mu\text{-}y) = \$\mu\text{-}(\text{real } x) * t$.
qed

lemma *exponential-p-mu*: $\$p\text{-}x = \exp(-\$ \mu\text{-}x)$ **if** *exponential-interpolation* $x+1 < \$\psi$ **for** $x::\text{nat}$

proof –

have $\$p\text{-}x = \exp(-\text{integral } \{ \text{real } x \text{ ..} < \text{real } x + 1 \} (\lambda y. \$\mu\text{-}y))$
using that apply (*rewrite p-exp-integral-mu; simp add: add.commute*)
apply (*rule integral-subset-negligible[THEN sym], force*)
by (*rewrite Icc-minus-Ico; simp*)
also have $\dots = \exp(-\$ \mu\text{-}(\text{real } x))$ **using that by** (*rewrite exponential-integral-mu; simp*)
finally show *?thesis* .
qed

corollary *exponential-mu-p*: $\$ \mu\text{-}x = -\ln(\$p\text{-}x)$ **if** *exponential-interpolation* $x+1 < \$\psi$ **for** $x::\text{nat}$
using *exponential-p-mu that by simp*

corollary *exponential-mu-xt-p*: $\$ \mu\text{-}(x+t) = -\ln(\$p\text{-}x)$
if *exponential-interpolation* $x+1 < \$\psi$ $0 \leq t < 1$ **for** $x::\text{nat}$ **and** $t::\text{real}$
using that *exponential-mu exponential-mu-p by presburger*

corollary *exponential-q-mu*: $\$q\text{-}x = 1 - \exp(-\$ \mu\text{-}x)$
if *exponential-interpolation* $x+1 < \$\psi$ **for** $x::\text{nat}$
using *exponential-p-mu that p-q-1*
by (*smt (verit, ccfv-SIG) ereal-less-le not-add-less1 of-nat-less-imp-less*)

lemma *exponential-p*: $\$p\text{-}\{t \& x\} = (\$p\text{-}x).\hat{t}$
if *exponential-interpolation* $x+1 < \$\psi$ $0 \leq t \leq 1$ **for** $x::\text{nat}$ **and** $t::\text{real}$

proof –

have [*simp*]: $\text{real } x + t < \$\psi$ **using that** *ereal-less-le* **by** *auto*
have $\$p\text{-}\{t \& \text{real } x\} = \exp(-\text{integral } \{ \text{real } x \text{ ..} < \text{real } x + t \} (\lambda y. \$\mu\text{-}y))$
using that apply (*rewrite p-exp-integral-mu, simp-all*)
apply (*rule integral-subset-negligible[THEN sym], force*)
by (*rewrite Icc-minus-Ico; simp*)
also have $\dots = \exp(-\$ \mu\text{-}(\text{real } x) * t)$
using that by (*rewrite exponential-integral-mu; simp*)
also have $\dots = (\$p\text{-}(\text{real } x)).\hat{t}$
using *exponential-p-mu that*
by (*smt (verit) exp-not-eq-zero exponential-mu-p mult.commute powr-def*)

finally show ?thesis .
qed

lemma exponential-q: $q\{t \& x\} = 1 - (1 - q\{x\})^t$
 if exponential-interpolation $x+1 < \psi$ $0 \leq t \leq 1$ for $x::nat$ and $t::real$
proof –
 have $q\{t \& real\ x\} = 1 - p\{t \& real\ x\}$
 using p-q-1 that by (smt (verit) ereal-less-le le-add1 of-nat-mono)
 also have $\dots = 1 - (p\{real\ x\})^t$ using that by (rewrite exponential-p; simp)
 also have $\dots = 1 - (1 - q\{real\ x\})^t$
 using p-q-1 that by (smt (verit) ereal-less-le not-add-less1 of-nat-less-imp-less)
 finally show ?thesis .
 qed

lemma exponential-l-p: $l(x+t) = l\ x * (p\ x)^t$
 if exponential-interpolation $x+1 < \psi$ $0 \leq t \leq 1$ for $x::nat$ and $t::real$
proof –
 have ereal (real x) < ψ using that ereal-less-le by auto
 hence $l(real\ x + t) = l\ x * p\{t \& real\ x\}$ using that by (rewrite p-l; simp)
 also have $\dots = l(real\ x) * (p\{real\ x\})^t$ using that by (rewrite exponential-p; simp)
 finally show ?thesis .
 qed

lemma exponential-l-has-derivative-at-frac:
 (($\lambda s. l(x+s)$) has-real-derivative ($- l\ x * \mu\ x * (p\ x)^t$)) (at t)
 if exponential-interpolation $x+1 < \psi$ $0 < t < 1$ for $x::nat$ and $t::real$
proof –
 let $?x = real\ x$
 have (($\lambda s. (p\ ?x)^s$) has-real-derivative ($-\ \mu\ ?x * (p\ ?x)^t$)) (at t)
 using that exponential-p-mu has-real-derivative-pow2
 by (metis Groups.ab-semigroup-mult-class.mult.commute exp-gt-zero ln-exp)
 hence (($\lambda s. l\ ?x * (p\ ?x)^s$) has-real-derivative ($- l\ ?x * \mu\ ?x * (p\ ?x)^t$)) (at t)
 by (rewrite minus-mult-commute, rewrite mult.assoc) (rule DERIV-cmult)
 moreover have $\forall_F s$ in nhds t. $l(?x + s) = l\ ?x * (p\ ?x)^s$
proof –
 let $?r = \min\ t\ (1-t)$
 have $?r > 0$ using that by simp
 moreover have $\bigwedge s. dist\ s\ t < ?r \implies l(?x + s) = l\ ?x * (p\ ?x)^s$
 using dist-real-def that exponential-l-p by force
 ultimately show ?thesis using eventually-nhds-metric by blast
 qed
 ultimately show ?thesis by (rewrite DERIV-cong-ev[where $g = \lambda s. l\ ?x * (p\ ?x)^s$]; simp)
 qed

lemma exponential-l-has-derivative-at-frac':
 (l has-real-derivative ($- l\ x * \mu\ x * (p\ x)^{y-x}$)) (at y)

if exponential-interpolation $x+1 < \psi$ $x < y$ $y < x+1$ for $x::\text{nat}$ and $y::\text{real}$
apply (rewrite DERIV-at-within-shift[where $x=y$ - real x and $z=\text{real } x$ and $S=\text{UNIV}$, simplified])
using exponential-l-has-derivative-at-frac that **by** simp

lemma exponential-l-differentiable-on-frac:
l differentiable-on $\{x <.. < x+1\}$ if exponential-interpolation $x+1 < \psi$ for $x::\text{nat}$
proof -
{ **fix** $y::\text{real}$ **assume** $y \in \{\text{real } x <.. < \text{real } (x+1)\}$
hence *l* differentiable at y
using exponential-l-has-derivative-at-frac' that real-differentiable-def **by** auto
}
thus ?thesis **unfolding** differentiable-on-def **by** (metis differentiable-at-withinI)
qed

lemma exponential-l-has-right-derivative-at-nat:
(*l* has-real-derivative $(-\ \$l-x * \$\mu-x)$) (at-right x)
if exponential-interpolation $x+1 < \psi$ for $x::\text{nat}$
proof -
let $?x = \text{real } x$
have [simp]: plus $?x$ ' $\{0 <.. \} = \{?x <.. \}$
unfolding image-def greaterThan-def **apply** simp
by (metis Groups.ab-semigroup-add-class.add-commute
add-minus-cancel neg-less-iff-less real-0-less-add-iff)
have [simp]: $\$p-x > 0$ **using** exponential-p-mu that **by** auto
hence [simp]: $\$p-x \neq 0$ **by** force
have $(\lambda s. \$l-(?x + s))$ has-real-derivative $(-\ \$l-?x * \$\mu-?x)$ (at-right 0)
apply (rewrite has-field-derivative-cong-eventually[where $g=\lambda s. \$l-?x * (\$p-?x). \hat{s}$])
using exponential-l-p that **apply** (intro eventually-at-rightI[of 0 1]; simp)
using powr-zero-eq-one **apply** simp
apply (rewrite minus-mult-commute, rule DERIV-cmult)
apply (rule has-field-derivative-at-within)
using that **apply** (rewrite exponential-mu-p; simp)
using has-real-derivative-powr2[of $\$p-x$ 0] powr-zero-eq-one **by** force
thus ?thesis
by (rewrite DERIV-at-within-shift[where $z=?x$ and $x=0$ and $S=\{0 <.. \}$,
simplified]) simp
qed

lemma exponential-l-has-left-derivative-at-nat:
(*l* has-real-derivative $(-\ \$l-x * \$\mu-(\text{real } x - 1))$) (at-left x)
if exponential-interpolation $x < \psi$ for $x::\text{nat}$
proof (cases x)
case 0
hence (*l* has-real-derivative 0) (at-left (real x))
apply (rewrite has-field-derivative-cong-eventually[where $g=\lambda-. \$l-0$]; simp)
apply (rule eventually-at-leftI[of -1]; simp)
using l-neg-nil less-eq-real-def **by** blast
moreover **have** $-\ \$l-x * \$\mu-(\text{real } x - 1) = 0$ **using** mu-unborn-0 0 **by** simp

ultimately show *?thesis* **by** *auto*
next
let $?x = \text{real } x$
case (*Suc y*)
let $?y = \text{real } y$
have [*simp*]: $?y + 1 = ?x$ **using** *Suc* **by** *simp*
hence [*simp*]: $\text{ereal } ?y < \$\psi$ **using that by** (*smt (verit) ereal-less-le*)
have [*simp*]: $\$p-?y > 0$ **using** *Suc exponential-p-mu* **that by** *auto*
have [*simp*]: $\text{plus } ?y \text{ ' } \{..<1\} = \{..<?x\}$
using *Suc unfolding image-def lessThan-def* **apply** *simp*
by (*metis (no-types, opaque-lifting) Groups.ab-semigroup-add-class.add commute add-less-cancel-right diff-add-cancel*)
have $((\lambda t. (\$p-?y).\hat{t}) \text{ has-real-derivative } (\$p-?y * -\$mu-?y))$ (*at-left 1*)
apply (*rule DERIV-subset[where s=UNIV]; simp*)
using that apply (*rewrite exponential-mu-p; simp add: add.commute[of 1 ?y]*)
by (*rule has-real-derivative-pow2[of \$p-?y 1, simplified]*)
hence $((\lambda t. \$l-?y * (\$p-?y).\hat{t}) \text{ has-real-derivative } (\$l-?y * \$p-?y * -\$mu-?y))$
(*at-left 1*)
by (*metis DERIV-cmult mult-ac(1)*)
moreover have $\$l-?y * \$p-?y = \$l-?x$ **using** *Suc p-1-l* **by** *simp*
ultimately have $((\lambda t. \$l-?y * (\$p-?y).\hat{t}) \text{ has-real-derivative } (- \$l-?x * \$mu-?y))$
(*at-left 1*)
by *simp*
moreover have $\forall_F t \text{ in } \text{at-left } 1. \$l-(?y + t) = \$l-?y * (\$p-?y).\hat{t}$
apply (*rule eventually-at-leftI[of 0]; simp*)
by (*rewrite exponential-l-p; simp add: that add.commute[of 1 ?y]*)
ultimately have $((\lambda t. \$l-(?y + t)) \text{ has-real-derivative } (- \$l-?x * \$mu-?y))$ (*at-left 1*)
by (*rewrite has-field-derivative-cong-eventually[where g=\lambda t. \\$l-?y * (\\$p-?y).\hat{t}; simp add: p-1-l]*)
thus *?thesis*
apply (*rewrite DERIV-at-within-shift[where S=\{..<1\} and z=?y and x=1, simplified]*)
using *Suc* **by** *simp*
qed

lemma *exponential-l-has-derivative-at-nat-iff-mu*:

$(l \text{ has-real-derivative } (- \$l-x * \$mu-x))$ (*at x*) \longleftrightarrow $\$mu-x = \$mu-(\text{real } x - 1)$

if *exponential-interpolation* $x+1 < \$\psi$ **for** $x::\text{nat}$

proof –

let $?x = \text{real } x$

have [*simp*]: $?x < \$\psi$ **using that by** (*simp add: ereal-less-le*)

hence [*simp*]: $\$l-?x \neq 0$ **by** *simp*

have $(l \text{ has-real-derivative } (- \$l-?x * \$mu-?x))$ (*at ?x*) \longleftrightarrow

$(l \text{ has-real-derivative } (- \$l-?x * \$mu-?x))$ (*at-right ?x*) \wedge

$(l \text{ has-real-derivative } (- \$l-?x * \$mu-?x))$ (*at-left ?x*)

using *has-real-derivative-at-split* **by** *auto*

also have $\dots \longleftrightarrow - \$l-?x * \$mu-?x = - \$l-?x * \$mu-(?x - 1)$ (**is** *?LHS* \longleftrightarrow *?RHS*)

proof
 assume $?LHS$
 hence $(l \text{ has-real-derivative } (- \$l-?x * \$\mu-?x))$ (at-left $?x$) **by simp**
 moreover have $(l \text{ has-real-derivative } (- \$l-?x * \$\mu-(?x - 1)))$ (at-left $?x$)
 using that exponential-l-has-left-derivative-at-nat **by force**
 ultimately show $?RHS$
 using has-real-derivative-iff-has-vector-derivative vector-derivative-unique-within
 trivial-limit-at-left-real **by blast**
next
 assume $?RHS$
 thus $?LHS$
 using that exponential-l-has-right-derivative-at-nat exponential-l-has-left-derivative-at-nat
 by force
qed
 also have $\dots \longleftrightarrow \$\mu-?x = \$\mu-(?x - 1)$ **by simp**
 finally show $?thesis$.
qed

lemma exponential-l-differentiable-at-nat-iff-mu:
 l differentiable at $x \longleftrightarrow \$\mu-x = \$\mu-(\text{real } x - 1)$
 if exponential-interpolation $x+1 < \$\psi$ for $x::\text{nat}$

proof
 let $?x = \text{real } x$
 assume l differentiable at $?x$
 from this obtain D where $DERIV-l: (l \text{ has-real-derivative } D)$ (at $?x$)
 using real-differentiable-def **by blast**
 hence $(l \text{ has-real-derivative } D)$ (at-right $?x$)
 using has-field-derivative-at-within **by blast**
 moreover have at $?x$ within $\{\text{real } x < ..\} \neq \perp$ **by simp**
 moreover have $(l \text{ has-real-derivative } (- \$l-\text{real } x * \$\mu-\text{real } x))$ (at-right $?x$)
 using exponential-l-has-right-derivative-at-nat that **by simp**
 ultimately have $D = - \$l-?x * \$\mu-?x$
 using that has-real-derivative-iff-has-vector-derivative vector-derivative-unique-within
 by blast
 thus $\$ \mu-?x = \$ \mu-(?x - 1)$
 using exponential-l-has-derivative-at-nat-iff-mu that $DERIV-l$ **by blast**
next
 assume $\$ \mu-(\text{real } x) = \$ \mu-(\text{real } x - 1)$
 thus l differentiable at $(\text{real } x)$
 using exponential-l-has-derivative-at-nat-iff-mu that real-differentiable-def **by blast**
qed

lemma exponential-L-d-mu: $\$L-x = \$d-x / \$\mu-x$
 if exponential-interpolation $\$ \mu-x \neq 0$ $x+1 < \$\psi$ for $x::\text{nat}$
proof –
 have $[simp]: \text{ereal } (\text{real } x) < \ψ **using that ereal-less-le by simp**
 have $[simp]: \$l-(\text{real } x) \neq 0$ **by simp**
 have $p\text{-pos}[simp]: \$p-(\text{real } x) > 0$ **using that by (simp add: exponential-p-mu)**

have [simp]: $\$p\text{-}(\text{real } x) \neq 1$ **using** that exponential-p-mu **by** simp
have $\$L\text{-}(\text{real } x) = (\text{LBINT } t:\{0..1\}. \$l\text{-}(\text{real } x + t))$
unfolding lives-def **by** (rewrite lborel-set-integral-Icc-shift[where t=x]) simp
also have $\dots = \text{integral } \{0..1\} (\lambda t. \$l\text{-}(\text{real } x + t))$
by (rule set-borel-integral-eq-integral-nonneg; simp)
also have $\dots = \text{integral } \{0..1\} (\lambda t. \$l\text{-}(\text{real } x) * (\$p\text{-}(\text{real } x)). \hat{t})$
apply (rule integral-cong)
using that **by** (rewrite exponential-l-p; simp)
also have $\dots = \$l\text{-}(\text{real } x) * \text{integral } \{0..1\} (\lambda t. (\$p\text{-}(\text{real } x)). \hat{t})$
using integral-mult-right **by** blast
also have $\dots = \$l\text{-}(\text{real } x) * (\$q\text{-}(\text{real } x) / - \ln (\$p\text{-}(\text{real } x)))$
proof -
have $\text{integral } \{0..1\} (\lambda t. (\$p\text{-}(\text{real } x)). \hat{t}) = ((\$p\text{-}(\text{real } x)). \hat{1} - 1) / \ln (\$p\text{-}(\text{real } x))$
apply (rule integral-unique)
by (intro has-integral-pow2-from-0; simp)
also have $\dots = \$q\text{-}(\text{real } x) / - \ln (\$p\text{-}(\text{real } x))$
using p-pos **apply** (rewrite powr-one, linarith)
using that **by** (rewrite in - - \sqcap p-q-1[of x 1, THEN sym]; simp)
finally show ?thesis **by** simp
qed
also have $\dots = \$d\text{-}(\text{real } x) / \$\mu\text{-}(\text{real } x)$ **using** that exponential-mu-p **by** (rewrite q-d-l; simp)
finally show ?thesis .
qed

lemma exponential-mx-mu: $\$m\text{-}x = \$\mu\text{-}x$ **if** exponential-interpolation $x+1 < \$\psi$
for $x::\text{nat}$
proof (cases $\langle \$\mu\text{-}(\text{real } x) = 0 \rangle$)
have lx-neq0: $\$l\text{-}(\text{real } x) \neq 0$ **using** ereal-less-le that **by** simp
case True
hence $\$q\text{-}(\text{real } x) = 0$ **using** exponential-q-mu that **by** simp
hence $\$d\text{-}(\text{real } x) = 0$
using q-1-d-l that **by** (metis lx-neq0 d-old-0 divide-eq-0-iff linorder-not-less zero-le-one)
hence $\$m\text{-}(\text{real } x) = 0$ **unfolding** die-central-def **by** simp
also have $\dots = \$\mu\text{-}(\text{real } x)$ **using** True **by** simp
finally show ?thesis .
next
case False
thus ?thesis **unfolding** die-central-def **using** exponential-L-d-mu that
by (smt (verit) divide-eq-0-iff divide-mult-cancel exp-eq-one-iff exponential-q-mu linorder-not-less mu-beyond-0 nonzero-mult-div-cancel-left q-1-d-l)
qed

lemma exponential-d-mu-sums-T: $(\lambda k. \$d\text{-}(x+k) / \$\mu\text{-}(x+k))$ sums $\$T\text{-}x$
if exponential-interpolation total-finite $\bigwedge k::\text{nat}. \$\mu\text{-}(x+k) \neq 0$ **for** $x::\text{nat}$
proof -
have $\neg \$\psi < \infty$

proof
 assume $\psi < \infty$
 from this obtain $y::\text{nat}$ where $\psi < \text{ereal}(\text{real } y)$ using *less-PIInf-Ex-of-nat*
 by *fastforce*
 hence $xy: \psi < \text{ereal}(\text{real}(x+y))$ by (*simp add: less-ereal-le*)
 hence $\mu(\text{real}(x+y)) \neq 0$ using *that* by *simp*
 moreover have $\mu(\text{real}(x+y)) = 0$ using *xy mu-beyond-0* by *simp*
 ultimately show *False ..*
qed
 hence $\psi = \infty$ by *simp*
 moreover hence $\bigwedge k::\text{nat}. d(\text{real}(x+k)) / \mu(\text{real}(x+k)) = L(\text{real}(x+k))$
 using *that* by (*rewrite exponential-L-d-mu; simp*)
 ultimately show *?thesis*
 apply (*rewrite sums-cong; simp*)
 by (*rule L-sums-T, simp add: that*)
qed

lemma *exponential-e-d-l-mu*: $(\lambda k. d(x+k) / (l \cdot x * \mu(x+k))) \text{ sums } e^{\circ-x}$
 if *exponential-interpolation total-finite* $\bigwedge k::\text{nat}. \mu(x+k) \neq 0$ for $x::\text{nat}$

proof –
 let $?x = \text{real } x$
 have $\neg \text{ereal } ?x \geq \psi$ using *that mu-beyond-0* by (*metis add-cancel-right-right*)
 hence [*simp*]: $\text{ereal } ?x < \psi$ by *simp*
 have $(\lambda k. d(\text{real}(x+k)) / \mu(\text{real}(x+k)) / l \cdot ?x \text{ sums } (T \cdot ?x / l \cdot ?x))$
 using *sums-divide exponential-d-mu-sums-T* that by *force*
 thus *?thesis* by (*rewrite e-T-l; simp add: mult.commute*)
qed

end

6.5 Limited Life Table

locale *limited-life-table* = *life-table* +
 assumes *l-limited*: $\exists x::\text{real}. l \cdot x = 0$
begin

lemma *limited-survival-function-MM-X*: *limited-survival-function* $\mathfrak{M} X$

proof (*rule limited-survival-function.intro*)
 show *survival-model* $\mathfrak{M} X$ by (*rule survival-model-MM-X*)
 show *limited-survival-function-axioms* $\mathfrak{M} X$
 unfolding *limited-survival-function-axioms-def* using *l-limited death-pt-def* by
fastforce
qed

end

sublocale *limited-life-table* \subseteq *limited-survival-function* $\mathfrak{M} X$
 by (*rule limited-survival-function-MM-X*)

```

context limited-life-table
begin

notation ult-age ( $\$w$ )

lemma l-omega-0:  $\$l-\$w = 0$ 
  using ccdfX-l-normal ccdfX-omega-0 by simp

lemma l-0-equiv-nat:  $\$l-x = 0 \iff x \geq \$w$  for  $x::nat$ 
  using ccdfX-l-normal ccdfX-0-equiv-nat by simp

lemma d-l-equiv-nat:  $\$d-\{t\&x\} = \$l-x \iff x+t \geq \$w$  if  $t \geq 0$  for  $x\ t :: nat$ 
  by (metis d-l-equiv of-nat-0-le-iff of-nat-add psi-le-iff-omega-le)

corollary d-1-omega-l:  $\$d-(\$w-1) = \$l-(\$w-1)$ 
  using d-l-equiv-nat[of 1 \$w-1] omega-pos by simp

lemma limited-life-table-imp-total-finite: total-finite
proof –
  have  $\{0..\} = \{0 .. real\ \$w\} \cup \{real\ \$w <..\}$  by force
  moreover have set-integrable lborel  $\{0 .. real\ \$w\}$  l by (rule l-integrable-Icc)
  moreover have set-integrable lborel  $\{real\ \$w <..\}$  l
    apply (rewrite set-integrable-cong[where f'= $\lambda$ -. 0], simp-all)
    using l-0-equiv-nat apply (meson l-0-equiv le-ereal-le order-le-less)
    unfolding set-integrable-def by simp
  ultimately have set-integrable lborel  $\{0..\}$  l
    using set-integrable-Un
    by (smt (verit, del-insts) Ici-subset-Ioi-iff add-mono-thms-linordered-field(1)
      atLeast-borel l-0-pos set-integrable-subset sets-lborel total-finite-iff-set-integrable-Ici)
  thus ?thesis unfolding total-finite-def by simp
qed

context
  fixes  $x::nat$ 
  assumes x-lt-omega[simp]:  $x < \$w$ 
begin

lemma curt-e-sum-l-finite-nat:  $\$e-x = (\sum k < n. \$l-(x+k+1)) / \$l-x$ 
  if  $\bigwedge k::nat. k < n \implies isCont\ l\ (x+k+1)$   $x+n \geq \$w$  for  $n::nat$ 
  apply (rewrite curt-e-sum-l-finite[of x n], simp)
  using that le-ereal-less psi-le-omega apply (metis of-nat-1 of-nat-add, force)
  by (simp add: add.commute)

end

end

end
theory Examples

```

```

imports Life-Table
begin

```

7 Examples

The following lemma is a verification of the solution to the multiple choice question No. 3 of Exam LTAM Spring 2022 by Society of Actuaries.

```

context smooth-survival-function
begin

```

```

lemma SoA-LTAM-2022-Spring-MCQ-No3:

```

```

  assumes  $\bigwedge x::real. 0 \leq x \implies x \leq 100 \implies cdf (distr \mathfrak{M} \text{ borel } X) x = (1 - 0.01*x).\hat{0.5}$ 

```

```

  shows  $|1000*\mu_{25} - 6.7| < 0.05$ 

```

```

proof -

```

```

  let ?svl = cdf (distr  $\mathfrak{M}$  borel X)

```

```

  have [simp]: ereal 25 <  $\psi$ 

```

```

    apply (rewrite not-le[THEN sym])

```

```

    using assms by (rewrite cdfX-0-equiv[THEN sym]) simp

```

```

  have  $\star$ : (( $\lambda x. \ln (1 - x/100)$ ) has-real-derivative (-1/75)) (at 25)

```

```

proof -

```

```

  have (( $\lambda x. (1 - x/100)$ ) has-real-derivative (0 - 1/100)) (at 25)

```

```

    apply (intro derivative-intros)

```

```

    by (rule DERIV-cdivide) simp

```

```

  hence (( $\lambda x. \ln (1 - x/100)$ ) has-real-derivative (1 / (1 - 25/100) * (-1/100))) (at 25)

```

```

    by (intro derivative-intros) auto

```

```

  thus ?thesis by simp

```

```

qed

```

```

have  $\exp \circ (\lambda x::real. 0.5 * \ln (1 - 0.01*x))$  field-differentiable at 25

```

```

  apply (rule field-differentiable-compose, simp-all)

```

```

  apply (rule derivative-intros, simp-all)

```

```

  using  $\star$  field-differentiable-def apply blast

```

```

  using field-differentiable-within-exp by blast

```

```

hence ?svl differentiable at 25

```

```

  apply (rewrite differentiable-eq-field-differentiable-real)

```

```

  by (rule field-differentiable-transform-within[where d=1])

```

```

    (simp-all add: powr-def dist-real-def assms)

```

```

hence  $\mu_{25} = - \text{deriv } (\lambda x. \ln (?svl x)) 25$  by (rule mu-deriv-ln; simp)

```

```

also have ... = 0.005 / (1 - 0.01*25)

```

```

proof -

```

```

  have  $\forall_F x$  in nhds 25.  $\ln (?svl x) * 2 = \ln (1 - x/100)$ 

```

```

proof -

```

```

  { fix  $x::real$  assume dist x 25 < 1

```

```

    hence asm: 0 < x < 100 using dist-real-def by auto

```

```

    hence  $\ln (?svl x) * 2 = \ln ((1 - 0.01*x).\hat{0.5}) * 2$  using assms by (smt

```

```

(verit))

```

```

    also have ... = (0.5 *  $\ln (1 - 0.01*x)$ ) * 2 using asm by (rewrite

```


\ln -power) auto
finally have $\ln (?svl x) * 2 = \ln (1 - x/100)$ **by** *simp* }
thus *?thesis* **by** (rewrite eventually-nhds-metric) (smt (verit, del-insts))
qed
hence $((\lambda x. \ln (?svl x)) \text{ has-real-derivative } (-0.005 / (1 - 0.01 * 25)))$ (at 25)
using \star **apply** (rewrite DERIV-cong-ev[**where** $g = \lambda x. 0.5 * \ln (1 - 0.01 * x)$],
simp-all)
by (rule derivative-eq-intros) auto
thus *?thesis* **using** DERIV-imp-deriv **by** force
qed
finally show *?thesis* **by** *simp*
qed
end

The following lemma is a verification of the solution to the problem No. 2. (1)-1 of Life Insurance Mathematics 2016 by the Institute of Actuaries of Japan, slightly modified; see the remark below.

context *smooth-life-table*
begin

lemma *IAJ-Life-Insurance-Math-2016-2-1-1*:

fixes $a b :: \text{real}$

assumes $-1 < a$ $a < 0$ $0 < b$ $-b/a \leq \psi$ **and**

total-finite **and**

$\bigwedge x. 0 < x \implies x < -b/a \implies l$ *differentiable at x* **and**

$\bigwedge x. 0 \leq x \implies x < -b/a \implies \mathbb{R}e^{\circ-x} = a*x + b$

shows $\bigwedge x. 0 \leq x \implies x < -b/a \implies \mathbb{R}l-x = \mathbb{R}l-0 * (b / (a*x + b)). \sim (a+1)/a$

proof –

fix x **assume** *asm-x*: $0 \leq x$ $x < -b/a$

hence *x-lt-psi[simp]*: $\mathbb{R}eal x < \psi$ **using** *assms ereal-le-less* **by** *presburger*

from *asm-x* **have** *axb-pos[simp]*: $a*x + b > 0$

using *assms* **by** (smt (verit, ccfv-threshold) *mult.commute neg-less-divide-eq*)

have *mu*: $\bigwedge t. t \in \{0 < .. < -b/a\} \implies \mathbb{R}\mu-t = (a + 1) / (a*t + b)$

proof –

fix t **assume** *asm-t*: $t \in \{0 < .. < -b/a\}$

with *assms* **have** $((\lambda u. \mathbb{R}e^{\circ-u}) \text{ has-real-derivative } (\mathbb{R}\mu-t * \mathbb{R}e^{\circ-t} - 1))$ (at t)

by (intro *e-has-derivative-mu-e-l'*[**where** $a=0$]; *simp*)

moreover have $((\lambda u. \mathbb{R}e^{\circ-u}) \text{ has-real-derivative } a)$ (at t)

proof –

let $?d = \min t (-b/a - t)$

have $?d > 0$ **using** *assms asm-t* **by** *simp*

moreover have $\bigwedge u. \text{dist } u t < ?d \implies \mathbb{R}e^{\circ-u} = a*u + b$ **using** *assms*

dist-real-def **by** auto

ultimately have $\forall_F u$ *in nhds t*. $\mathbb{R}e^{\circ-u} = a*u + b$ **by** (rewrite eventually-nhds-metric) *blast*

thus *?thesis*

using *assms* **apply** (rewrite DERIV-cong-ev[**where** $g = \lambda t. a*t + b$], *simp-all*)

by (intro derivative-eq-intros) auto

qed
ultimately have $\mu - t * e^{\circ-t} - 1 = a$ **using** *DERIV-unique* **by** *blast*
moreover have $e^{\circ-t} = a*t + b$ **using** *assms asm-t* **by** *simp*
ultimately show $\mu - t = (a + 1) / (a*t + b)$
using *assms* **by** (*smt (verit) mult-eq-0-iff nonzero-mult-div-cancel-right*)
qed
hence $p\{x \neq 0\} = (b / (a*x + b)).\wedge((a+1)/a)$
proof –
have *integ-mu: integral {0..x} (λt. μ-t) = (a + 1) / a * ln ((a*x + b) / b)*
proof –
have *integral {0..x} (λt. μ-t) = integral {0<..x} (λt. μ-t)*
apply (*rule integral-spike-set*)
apply (*rule negligible-subset[where s={0}]; force*)
by (*rule negligible-subset[where s={}]; simp*)
also have $\dots = \text{integral } \{0 < ..x\} (\lambda t. ((a + 1) / a) * (a / (a*t + b)))$
using *asm-x assms* **by** (*intro integral-cong*) (*rewrite mu; simp*)
also have $\dots = (a + 1) / a * \text{integral } \{0 < ..x\} (\lambda t. a / (a*t + b))$
by (*metis integral-mult-right*)
also have $\dots = (a + 1) / a * \ln ((a*x + b) / b)$
proof –
have *integral {0<..x} (λt. a / (a*t + b)) = integral {0..x} (λt. a / (a*t +*
b))
apply (*rule integral-spike-set*)
apply (*rule negligible-subset[where s={}]; simp*)
by (*rule negligible-subset[where s={0}]; force*)
also have $\dots = \ln (a*x + b) - \ln (a*0 + b)$
apply (*rule integral-unique*)
using *assms asm-x* **apply** (*intro inverse-fun-has-integral-ln, simp-all*)
using *axb-pos assms* **apply** (*smt (verit) mult-less-cancel-left*)
apply (*intro continuous-intros*)
by (*intro derivative-eq-intros*) *auto*
also have $\dots = \ln ((a*x + b) / b)$ **using** *assms* **by** (*rewrite ln-div; simp*)
finally have *integral {0<..x} (λt. a / (a*t + b)) = ln ((a*x + b) / b) .*
thus *?thesis* **by** *simp*
qed
finally show *?thesis* .
qed
thus *?thesis*
apply (*rewrite p-exp-integral-mu, simp-all add: asm-x*)
unfolding *powr-def* **using** *assms*
by *simp (smt (verit) axb-pos ln-div*
nonzero-minus-divide-divide nonzero-minus-divide-right right-diff-distrib')
qed
thus $l - x = l - 0 * (b / (a*x + b)).\wedge((a+1)/a)$
using *assms asm-x* **apply** (*rewrite in asm p-l, simp-all*)
by (*metis divide-mult-cancel l-0-neq-0 mult commute*)
qed

REMARK. The original problem lacks the following hypotheses: (i) $0 < b$,

(ii) $-b/a \leq \psi$, (iii) $\forall x. 0 < x < -b/a \rightarrow l$ differentiable at x , (iv) $\forall x. 0 \leq x < -b/a \rightarrow l$ integrable-on $\{x.. \}$. Moreover, the hypothesis $\forall x. 0 \leq x < -b/a$ is originally $\forall x. 0 \leq x \leq -b/a$.

end

end