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RFC 9666

Area Proxy for IS-IS

Abstract

Link-state routing protocols have hierarchical abstraction already built into them. However, when lower levels are used for transit, they must expose their internal topologies to each other, thereby leading to scaling issues.

To avoid such issues, this document discusses extensions to the IS-IS routing protocol that allow Level 1 (L1) areas to provide transit but only inject an abstraction of the Level 1 topology into Level 2 (L2). Each Level 1 area is represented as a single Level 2 node, thereby enabling a greater scale.

Status of This Memo

This document is not an Internet Standards Track specification; it is published for examination, experimental implementation, and evaluation.

This document defines an Experimental Protocol for the Internet community. This document is a product of the Internet Engineering Task Force (IETF). It represents the consensus of the IETF community. It has received public review and has been approved for publication by the Internet Engineering Steering Group (IESG). Not all documents approved by the IESG are candidates for any level of Internet Standard; see Section 2 of RFC 7841.

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1. Introduction

The IS-IS routing protocol [ISO10589] supports a two-level hierarchy of abstraction. The fundamental unit of abstraction is the "area", which is a (hopefully) connected set of systems running IS-IS at the same level. Level 1, the lowest level, is abstracted by routers that participate in both Level 1 and Level 2, and they inject area information into Level 2. Level 2 systems seeking to access Level 1 use this abstraction to compute the shortest path to the Level 1 area. The full topology database of Level 1 is not injected into Level 2, rather, only a summary of the address space contained within the area is injected. Therefore, the scalability of the Level 2 Link State Database (LSDB) is protected.

This works well if the Level 1 area is tangential to the Level 2 area. This also works well if there are several routers in both Levels 1 and 2 and they are adjacent to one another, so Level 2 traffic will never need to transit Level 1 only routers. Level 1 will not contain any Level 2 topology and Level 2 will only contain area abstractions for Level 1.

Unfortunately, this scheme does not work so well if the Level 1 only area needs to provide transit for Level 2 traffic. For Level 2 Shortest Path First (SPF) computations to work correctly, the transit topology must also appear in the Level 2 LSDB. This implies that all routers that could provide transit plus any links that might also provide Level 2 transit must also become part of the Level 2 topology. If this is a relatively tiny portion of the Level 1 area, this is not overly painful.

However, with today's data center topologies, this is problematic. A common application is to use a Layer 3 Leaf-Spine (L3LS) topology, which is a folded 3-stage Clos fabric [Clos]. It can also be thought of as a complete bipartite graph. In such a topology, the desire is to use Level 1 to contain the routing dynamics of the entire L3LS topology and then use Level 2 for the remainder of the network. Leaves in the L3LS topology are appropriate for connection outside of the data center itself, so they would provide connectivity for Level 2. If there are multiple connections to Level 2 for redundancy or other areas, these would also be made to the leaves in the topology. This creates a difficulty because there are now multiple Level 2 leaves in the topology, with connectivity between the leaves provided by the spines.

Following the current rules of IS-IS, all spine routers would necessarily be part of the Level 2 topology plus all links between a Level 2 leaf and the spines. In the limit, where all leaves need to support Level 2, it implies that the entire L3LS topology becomes part of Level 2. This is seriously problematic, as it more than doubles the LSDB held in the L3LS topology and eliminates any benefits of the hierarchy.

This document discusses the handling of IP traffic. Supporting MPLS-based traffic is a subject for future work.

1.1. Requirements Language

The key words "MUST", "MUST NOT", "REQUIRED", "SHALL", "SHALL NOT", "SHOULD", "SHOULD NOT", "RECOMMENDED", "NOT RECOMMENDED", "MAY", and "OPTIONAL" in this document are to be interpreted as described in BCP 14 [RFC2119] [RFC8174] when, and only when, they appear in all capitals, as shown here.

2. Area Proxy

In this specification, we completely abstract away the details of the Level 1 area topology within Level 2, making the entire area look like a single proxy system directly connected to all of the area's Level 2 neighbors. By only providing an abstraction of the topology, Level 2's requirement for connectivity can be satisfied without the full overhead of the area's internal topology. It then becomes the responsibility of the Level 1 area to provide the forwarding connectivity that's advertised.

For this discussion, we'll consider a single Level 1 IS-IS area to be the Inside Area and the remainder of the Level 2 area to be the Outside Area. All routers within the Inside Area speak Level 1 and Level 2 IS-IS on all of the links within the topology. We propose to implement Area Proxy by having a Level 2 Proxy Link State PDU (LSP) that represents the entire Inside Area. We will refer to this as the Proxy LSP. This is the only LSP from the area that will be flooded into the overall Level 2 LSDB.

There are four classes of routers that we need to be concerned with in this discussion:

Inside Router: A router within the Inside Area that runs Level 1 and Level 2 IS-IS. A router is recognized as an Inside Router by the existence of its LSP in the Level 1 LSDB.

Area Leader: The Area Leader is an Inside Router that is elected to represent the Level 1 area by injecting the Proxy LSP into the Level 2 LSDB. There may be multiple candidates for Area Leader, but only one is elected at a given time. Any Inside Router can be the Area Leader.

Inside Edge Router: An Inside Edge Router is an Inside Area Router that has at least one Level 2 interface outside of the Inside Area. An interface on an Inside Edge Router that is connected to an Outside Edge Router is an Area Proxy Boundary.

Outside Edge Router: An Outside Edge Router is a Level 2 router that is outside of the Inside Area that has an adjacency with an Inside Edge Router.

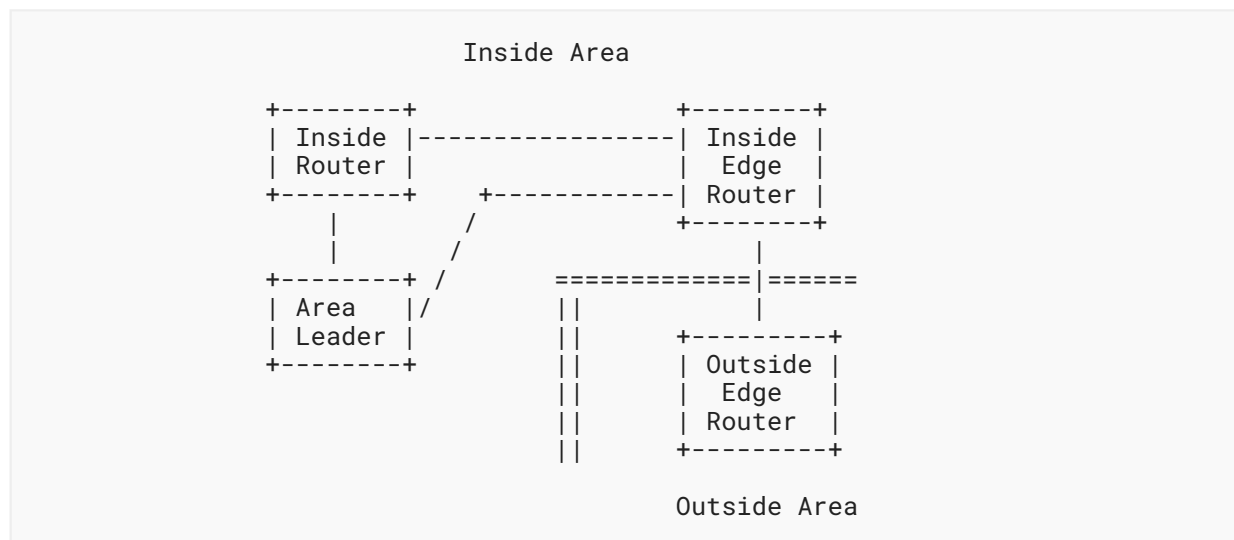


Figure 1: An Example of Router Classes

All Inside Edge Routers learn the Area Proxy System Identifier from the Area Proxy TLV advertised by the Area Leader and use that as the system identifier in their Level 2 IS-IS Hello (IIH) PDUs on all Outside interfaces. Outside Edge Routers will then advertise an adjacency to the Area Proxy System Identifier. This allows all Outside Routers to use the Proxy LSP in their SPF computations without seeing the full topology of the Inside Area.

Area Proxy functionality assumes that all circuits on Inside Routers are either Level 1-2 circuits within the Inside Area, or Level 2 circuits between Outside Edge Routers and Inside Edge Routers.

Area Proxy Boundary multi-access circuits (i.e., Ethernets in LAN mode) with multiple Inside Edge Routers on them are not supported. The Inside Edge Router on any boundary LAN **MUST NOT** flood Inside Router LSPs on this link. Boundary LANs **SHOULD NOT** be enabled for Level 1. An Inside Edge Router may be elected as the Designated Intermediate System (DIS) for a Boundary LAN. In this case, using the Area Proxy System ID as the basis for the LAN pseudonode identifier could create a collision, so the Insider Edge Router **SHOULD** compose the pseudonode

identifier using its originally configured system identifier. This choice of pseudonode identifier may confuse neighbors with an extremely strict implementation. In this case, the Inside Edge Router may be configured with priority 0, causing an Outside Router to be elected as the DIS.

2.1. Segment Routing

If the Inside Area supports Segment Routing (SR) [RFC8402], then all Inside Nodes **MUST** advertise a Segment Routing Global Block (SRGB). The first value of the SRGB advertised by all Inside Nodes **MUST** start at the same value. If the Area Leader detects SRGBs that do not start with the same value, it **MUST** log an error and not advertise an SRGB in the Proxy LSP. The range advertised for the area will be the minimum of that advertised by all Inside Nodes.

To support SR, the Area Leader will take the SRGB information found in the L1 LSDB and convey that to L2 through the Proxy LSP. Prefixes with Segment Identifier (SID) assignments will be copied to the Proxy LSP. Adjacency SIDs for Outside Edge Nodes will be copied to the Proxy LSP.

To further extend SR, it is helpful to have a segment that refers to the entire Inside Area. This allows a path to refer to an area and have any node within that area accept and forward the packet. In effect, this becomes an anycast SID that is accepted by all Inside Edge Nodes. The information about this SID is distributed in the Area SID sub-TLV as part of the Area Leader's Area Proxy TLV (Section 4.3.2). The Inside Edge Nodes **MUST** establish forwarding based on this SID. The Area Leader **SHALL** also include the Area SID in the Proxy LSP so that the remainder of L2 can use it for path construction. (Section 4.4.13).

3. Inside Router Functions

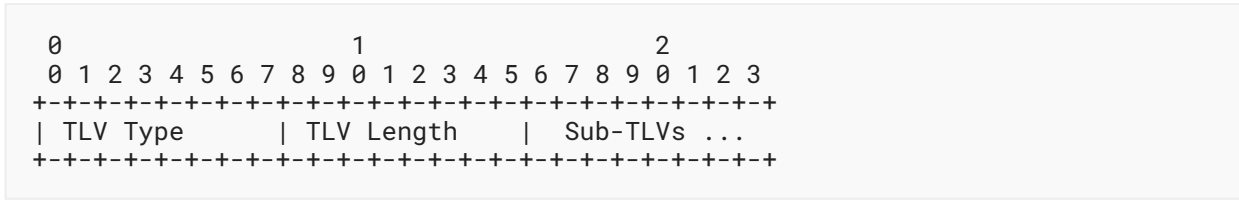
All Inside Routers run Level 1-2 IS-IS and must be explicitly instructed to enable the Area Proxy functionality. To signal their readiness to participate in Area Proxy functionality, they will advertise the Area Proxy TLV in their L2 LSP.

3.1. The Area Proxy TLV

The Area Proxy TLV serves multiple functions:

- The presence of the Area Proxy TLV in a node's LSP indicates that the node is enabled for Area Proxy.
- An LSP containing the Area Proxy TLV is also an Inside Node. All Inside Nodes, including pseudonodes, **MUST** advertise the Area Proxy TLV.
- It is a container for sub-TLVs with Area Proxy information.

A node advertises the Area Proxy TLV in fragment 0 of its L2 LSP. Nodes **MUST NOT** advertise the Area Proxy TLV in an L1 LSP. Nodes **MUST** ignore the Area Proxy TLV if it is found in an L1 LSP. The Area Proxy TLV is not used in the Proxy LSP. The format of the Area Proxy TLV is:



TLV Type: 20

TLV Length: Length of the sub-TLVs.

3.2. Level 2 SPF Computation

When Outside Routers perform a Level 2 SPF computation, they will use the Proxy LSP for computing a path transiting the Inside Area. Because the topology has been abstracted away, the cost for transiting the Inside Area will be zero.

When Inside Routers perform a Level 2 SPF computation, they **MUST** ignore the Proxy LSP. Because these systems see the Inside Area topology, the link metrics internal to the area are visible. This could lead to different and possibly inconsistent SPF results, potentially leading to forwarding loops.

To prevent this, the Inside Routers **MUST** consider the metrics of links outside of the Inside Area (inter-area metrics) separately from the metrics of the Inside Area links (intra-area metrics). Intra-area metrics **MUST** be treated as less than any inter-area metric. Thus, if two paths have different total inter-area metrics, the path with the lower inter-area metric would be preferred regardless of any intra-area metrics involved. However, if two paths have equal inter-area metrics, then the intra-area metrics would be used to compare the paths.

Point-to-point links between two Inside Routers are considered to be Inside Area links. LAN links that have a pseudonode LSP in the Level 1 LSDB are considered to be Inside Area links.

3.3. Responsibilities Concerning the Proxy LSP

The Area Leader will generate a Proxy LSP that will be flooded across the Inside Area. Inside Routers **MUST** flood the Proxy LSP and **MUST** ignore its contents. The Proxy LSP uses the Area Proxy System Identifier as its Source ID.

4. Area Leader Functions

The Area Leader has several responsibilities. First, it **MUST** inject the Area Proxy System Identifier into the Level 2 LSDB. Second, the Area Leader **MUST** generate the Proxy LSP for the Inside Area.

4.1. Area Leader Election

The Area Leader is selected using the election mechanisms and TLVs described in "Dynamic Flooding on Dense Graphs" [RFC9667].

4.2. Redundancy

If the Area Leader fails, another candidate may become Area Leader and **MUST** regenerate the Proxy LSP. The failure of the Area Leader is not visible outside of the area and appears to simply be an update of the Proxy LSP.

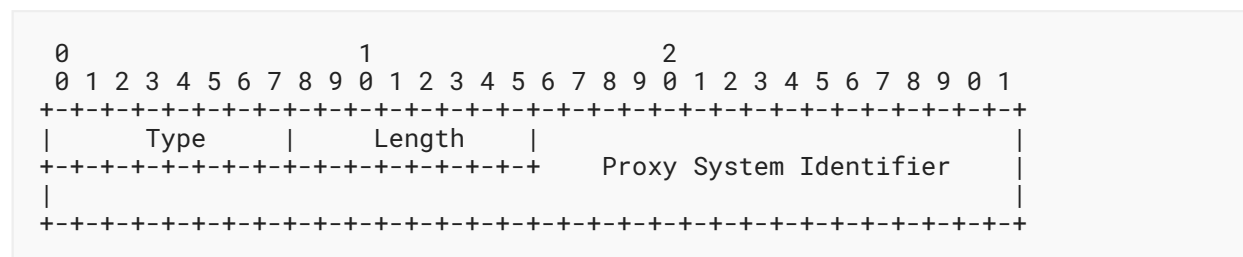
For consistency, all Area Leader candidates **SHOULD** be configured with the same Proxy System ID, Proxy Hostname, and any other information that may be inserted into the Proxy LSP.

4.3. Distributing Area Proxy Information

The Area Leader is responsible for distributing information about the area to all Inside Nodes. In particular, the Area Leader distributes the Proxy System ID and the Area SID. This is done using two sub-TLVs of the Area Proxy TLV.

4.3.1. The Area Proxy System Identifier Sub-TLV

The Area Proxy System Identifier sub-TLV **MUST** be used by the Area Leader to distribute the Area Proxy System ID. This is an additional system identifier that is used by Inside Nodes as an indication that Area Proxy is active. The format of this sub-TLV is:



Type: 1

Length: Length of a system ID (6).

Proxy System Identifier: The Area Proxy System Identifier.

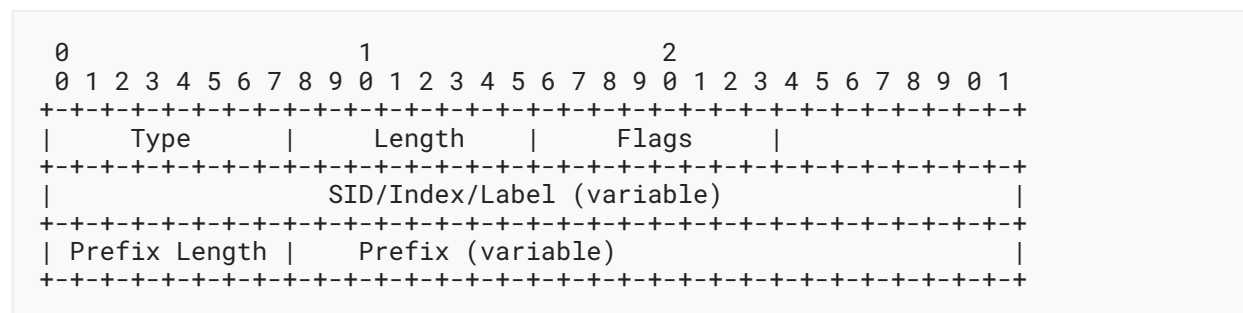
The Area Leader **MUST** advertise the Area Proxy System Identifier sub-TLV when it observes that all Inside Routers are advertising the Area Proxy TLV. Their advertisements indicate that they are individually ready to perform Area Proxy functionality. The Area Leader then advertises the Area Proxy System Identifier TLV to indicate that the Inside Area **MUST** enable Area Proxy functionality.

Other candidates for Area Leader **MAY** also advertise the Area Proxy System Identifier when they observe that all Inside Routers are advertising the Area Proxy TLV. All candidates advertising the Area Proxy System Identifier TLV **SHOULD** be advertising the same system identifier. Multiple proxy system identifiers in a single area is a misconfiguration and each unique occurrence **SHOULD** be logged. Systems should use the Proxy System ID advertised by the Area Leader.

The Area Leader and other candidates for Area Leader **MAY** withdraw the Area Proxy System Identifier when one or more Inside Routers are not advertising the Area Proxy TLV. This will disable Area Proxy functionality. However, before withdrawing the Area Proxy System Identifier, an implementation **SHOULD** protect against unnecessary churn from transients by delaying the withdrawal. The amount of delay is implementation dependent.

4.3.2. The Area SID Sub-TLV

The Area SID sub-TLV allows the Area Leader to advertise a prefix and SID that represent the entirety of the Inside Area to the Outside Area. This sub-TLV is learned by all of the Inside Edge Nodes who should consume this SID at forwarding time. The Area SID sub-TLV has the following format:

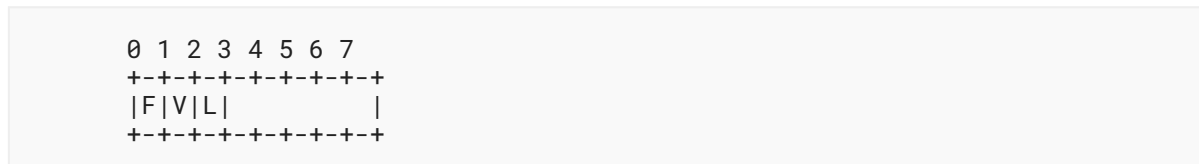


where:

Type: 2

Length: Variable (1 + SID length)

Flags: 1 octet, defined as follows.



F: Address-Family Flag. If this flag is not set, then this proxy SID is used when forwarding IPv4-encapsulated traffic. If set, then this proxy SID is used when forwarding IPv6-encapsulated traffic.

V: Value Flag. If set, then the proxy SID carries a value, as defined in [RFC8667], Section 2.1.1.1.

L: Local Flag. If set, then the value/index carried by the proxy SID has local significance, as defined in [RFC8667], Section 2.1.1.1.

Other bits: **MUST** be zero when originated and ignored when received.

SID/Index/Label: As defined in [RFC8667], Section 2.1.1.1.

Prefix Length: 1 octet

Prefix: 0-16 octets

4.4. Proxy LSP Generation

Each Inside Router generates a Level 2 LSP and the Level 2 LSPs for the Inside Edge Routers will include adjacencies to Outside Edge Routers. Unlike normal Level 2 operations, these LSPs are not advertised outside of the Inside Area and **MUST** be filtered by all Inside Edge Routers to not be flooded to Outside Routers. Only the Proxy LSP is injected into the overall Level 2 LSDB.

The Area Leader uses the Level 2 LSPs generated by the Inside Edge Routers to generate the Proxy LSP. This LSP is originated using the Area Proxy System Identifier. The Area Leader can also insert the following additional TLVs into the Proxy LSP for additional information for the Outside Area. LSPs generated by unreachable nodes **MUST NOT** be considered.

4.4.1. The Protocols Supported TLV

The Area Leader **SHOULD** insert a Protocols Supported TLV (129) [[RFC1195](#)] into the Proxy LSP. The values included in the TLV **SHOULD** be the protocols supported by the Inside Area.

4.4.2. The Area Address TLV

The Area Leader **SHOULD** insert an Area Addresses TLV (1) [[ISO10589](#)] into the Proxy LSP.

4.4.3. The Dynamic Hostname TLV

It is **RECOMMENDED** that the Area Leader insert the Dynamic Hostname TLV (137) [[RFC5301](#)] into the Proxy LSP. The contents of the hostname may be specified by configuration. The presence of the hostname helps to simplify network debugging.

4.4.4. The IS Neighbors TLV

The Area Leader can insert the IS Neighbors TLV (2) [[ISO10589](#)] into the Proxy LSP for Outside Edge Routers. The Area Leader learns of the Outside Edge Routers by examining the LSPs generated by the Inside Edge Routers copying any IS Neighbors TLVs referring to Outside Edge Routers into the Proxy LSP. Since the Outside Edge Routers advertise an adjacency to the Area Proxy System Identifier, this will result in a bidirectional adjacency.

An entry for a neighbor in both the IS Neighbors TLV and the Extended IS Neighbors TLV would be functionally redundant, so the Area Leader **SHOULD NOT** do this. The Area Leader **MAY** omit either the IS Neighbors TLV or the Extended IS Neighbors TLV, but it **MUST** include at least one of them.

4.4.5. The Extended IS Neighbors TLV

The Area Leader can insert the Extended IS Reachability TLV (22) [[RFC5305](#)] into the Proxy LSP. The Area Leader **SHOULD** copy each Extended IS Reachability TLV advertised by an Inside Edge Router about an Outside Edge Router into the Proxy LSP.

If the Inside Area supports Segment Routing, and Segment Routing selects a SID where the L-Flag is not set, then the Area Lead **SHOULD** include an Adjacency Segment Identifier sub-TLV (31) [RFC8667] using the selected SID.

If the inside area supports SRv6, the Area Leader **SHOULD** copy the "SRv6 End.X SID" and "SRv6 LAN End.X SID" sub-TLVs of the Extended IS Reachability TLVs advertised by Inside Edge Routers about Outside Edge Routers.

If the inside area supports Traffic Engineering (TE), the Area Leader **SHOULD** copy TE-related sub-TLVs ([RFC5305], Section 3) to each Extended IS Reachability TLV in the Proxy LSP.

4.4.6. The MT Intermediate Systems TLV

If the Inside Area supports Multi-Topology (MT), then the Area Leader **SHOULD** copy each Outside Edge Router advertisement that is advertised by an Inside Edge Router in an MT Intermediate Systems TLV into the Proxy LSP.

4.4.7. Reachability TLVs

The Area Leader **SHOULD** insert additional TLVs describing any routing prefixes that should be advertised on behalf of the area. These prefixes may be learned from the Level 1 LSDB, Level 2 LSDB, or redistributed from another routing protocol. This applies to all of the various types of TLVs used for prefix advertisement:

- IP Internal Reachability Information TLV (128) [RFC1195]
- IP External Reachability Information TLV (130) [RFC1195]
- Extended IP Reachability TLV (135) [RFC5305]
- IPv6 Reachability TLV (236) [RFC5308]
- Multi-Topology Reachable IPv4 Prefixes TLV (235) [RFC5120]
- Multi-Topology Reachable IPv6 Prefixes TLV (237) [RFC5120]

For TLVs in the Level 1 LSDB, for a given TLV type and prefix, the Area Leader **SHOULD** select the TLV with the lowest metric and copy that TLV into the Proxy LSP.

When examining the Level 2 LSDB for this function, the Area Leader **SHOULD** only consider TLVs advertised by Inside Routers. Further, for prefixes that represent Boundary links, the Area Leader **SHOULD** copy all TLVs that have unique sub-TLV contents.

If the Inside Area supports SR and the selected TLV includes a Prefix Segment Identifier sub-TLV (3) [RFC8667], then the sub-TLV **SHOULD** be copied as well. The P-Flag **SHOULD** be set in the copy of the sub-TLV to indicate that penultimate hop popping should not be performed for this prefix. The E-Flag **SHOULD** be reset in the copy of the sub-TLV to indicate that an explicit NULL is not required. The R-Flag **SHOULD** simply be copied.

4.4.8. The Router Capability TLV

The Area Leader **MAY** insert the Router Capability TLV (242) [RFC7981] into the Proxy LSP. If SR is supported by the inside area, as indicated by the presence of an SRGB being advertised by all Inside Nodes, then the Area Leader **SHOULD** advertise an SR-Capabilities sub-TLV (2) [RFC8667]

with an SRGB. The first value of the SRGB is the same as the first value advertised by all Inside Nodes. The range advertised for the area will be the minimum of all ranges advertised by Inside Nodes. The Area Leader **SHOULD** use its Router ID in the Router Capability TLV.

If SRv6 Capability sub-TLV [RFC7981] is advertised by all Inside Routers, the Area Leader should insert an SRv6 Capability sub-TLV in the Router Capability TLV. Each flag in the SRv6 Capability sub-TLV should be set if the flag is set by all Inside Routers.

If the Node Maximum SID Depth (MSD) sub-TLV [RFC8491] is advertised by all Inside Routers, the Area Leader should advertise the intersection of the advertised MSD types and the smallest supported MSD values for each type.

4.4.9. The Multi-Topology TLV

If the Inside Area supports multi-topology, then the Area Leader **SHOULD** insert the Multi-Topology TLV (229) [RFC5120], including the topologies supported by the Inside Nodes.

If any Inside Node is advertising the O (Overload) bit for a given topology, then the Area Leader **MUST** advertise the O bit for that topology. If any Inside Node is advertising the A (Attach) bit for a given topology, then the Area Leader **MUST** advertise the A bit for that topology.

4.4.10. The SID/Label Binding and the Multi-Topology SID/Label Binding TLV

If an Inside Node advertises the SID/Label Binding or Multi-Topology SID/Label Binding TLV [RFC8667], then the Area Leader **MAY** copy the TLV to the Proxy LSP.

4.4.11. The SRv6 Locator TLV

If the inside area supports SRv6, the Area Leader **SHOULD** copy all SRv6 locator TLVs [RFC9352] advertised by Inside Routers to the Proxy LSP.

4.4.12. Traffic Engineering Information

If the inside area supports TE, the Area Leader **SHOULD** advertise a TE Router ID TLV (134) [RFC5305] in the Proxy LSP. It **SHOULD** copy the Shared Risk Link Group (SRLS) TLVs (138) [RFC5307] advertised by Inside Edge Routers about links to Outside Edge Routers.

If the inside area supports IPv6 TE, the Area Leader **SHOULD** advertise an IPv6 TE Router ID TLV (140) [RFC6119] in the Proxy LSP. It **SHOULD** also copy the IPv6 SRLG TLVs (139) [RFC6119] advertised by Inside Edge Routers about links to Outside Edge Routers.

4.4.13. The Area SID

When SR is enabled, it may be useful to advertise an Area SID that will direct traffic to any of the Inside Edge Routers. The information for the Area SID is distributed to all Inside Edge Routers using the Area SID sub-TLV (Section 4.3.2) by the Area Leader.

The Area Leader **SHOULD** advertise the Area SID information in the Proxy LSP as a Node SID as defined in [RFC8667], Section 2.1. The advertisement in the Proxy LSP informs the Outside Area that packets directed to the SID will be forwarded to one of the Inside Edge Nodes and the Area SID will be consumed.

Other uses of the Area SID and Area SID prefix are outside the scope of this document. Documents that define other use cases for the Area SID **MUST** specify whether the SID value should be the same or different from that used in support of Area Proxy.

5. Inside Edge Router Functions

The Inside Edge Router has two additional and important functions. First, it **MUST** generate IIHs that appear to have come from the Area Proxy System Identifier. Second, it **MUST** filter the L2 LSPs, Partial Sequence Number PDUs (PSNPs), and Complete Sequence Number PDUs (CSNPs) that are being advertised to Outside Routers.

5.1. Generating L2 IIHs to Outside Routers

The Inside Edge Router has one or more Level 2 interfaces to the Outside Routers. These may be identified by explicit configuration or by the fact that they are not also Level 1 circuits. On these Level 2 interfaces, the Inside Edge Router **MUST NOT** send an IIH until it has learned the Area Proxy System ID from the Area Leader. Then, once it has learned the Area Proxy System ID, it **MUST** generate its IIHs on the circuit using the Proxy System ID as the source of the IIH.

Using the Proxy System ID causes the Outside Router to advertise an adjacency to the Proxy System ID, not to the Inside Edge Router, which supports the proxy function. The normal system ID of the Inside Edge Router **MUST NOT** be used as it will cause unnecessary adjacencies to form.

5.2. Filtering LSP Information

For the area proxy abstraction to be effective the L2 LSPs generated by the Inside Routers **MUST** be restricted to the Inside Area. The Inside Routers know which system IDs are members of the Inside Area based on the advertisement of the Area Proxy TLV. To prevent unwanted LSP information from escaping the Inside Area, the Inside Edge Router **MUST** perform filtering of LSP flooding, CSNPs, and PSNPs. Specifically:

- A Level 2 LSP with a source system identifier that is found in the Level 1 LSDB **MUST NOT** be flooded to an Outside Router.
- A Level 2 LSP that contains the Area Proxy TLV **MUST NOT** be flooded to an Outside Router.
- A Level 2 CSNP sent to an Outside Router **MUST NOT** contain any information about an LSP with a system identifier found in the Level 1 LSDB. If an Inside Edge Router filters a CSNP and there is no remaining content, then the CSNP **MUST NOT** be sent. The source address of the CSNP **MUST** be the Area Proxy System ID.
- A Level 2 PSNP sent to an Outside Router **MUST NOT** contain any information about an LSP with a system identifier found in the Level 1 LSDB. If an Inside Edge Router filters a PSNP and there is no remaining content, then the PSNP **MUST NOT** be sent. The source address of the PSNP **MUST** be the Area Proxy System ID.

6. IANA Considerations

IANA has assigned code point 20 from the "IS-IS TLV Codepoints" registry for the Area Proxy TLV. The registry fields are IIH:n, LSP:y, SNP:n, and Purge:n.

In association with this, IANA has created a "IS-IS Sub-TLVs for the Area Proxy TLV" registry. Temporary registrations may be made via early allocation [RFC7120].

The registration procedure is Expert Review [RFC8126]. The values are from 0-255, and the fields are Value, Name, and Reference. The initial assignments are as follows.

Value	Name	Reference
1	Area Proxy System Identifier	RFC 9666
2	Area SID	RFC 9666

Table 1

7. Security Considerations

This document introduces no new security issues. Security of routing within a domain is already addressed as part of the routing protocols themselves. This document proposes no changes to those security architectures. Security for IS-IS is provided by "IS-IS Cryptographic Authentication" [RFC5304] and "IS-IS Generic Cryptographic Authentication" [RFC5310].

8. References

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